

PRELIMINARY DATASHEET

## DATASHEET

**PRODUCT :** 32M (x16) Flash Memory

**MODEL No :** **LH28F320BJHG-PBTLZ2**

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# LH28F320BJHG-PBTLZ2

## 32M-BIT ( 2Mbit ×16 )

### Boot Block Flash MEMORY

- **Low Voltage Operation**
  - $V_{CC}=V_{CCW}=2.7V-3.6V$  Single Voltage
- **OTP(One Time Program) Block**
  - 3963 word + 4 word Program only array
- **16bit I/O Interface**
- **High-Performance Read Access Time**
  - 90ns( $V_{CC}=2.7V-3.6V$ )
- **Operating Temperature**
  - $-40^{\circ}C$  to  $+85^{\circ}C$
- **Low Power Management**
  - Typ. 4 $\mu$ A ( $V_{CC}=3.0V$ ) Standby Current
  - Automatic Power Savings Mode Decreases  $I_{CCR}$  in Static Mode
  - Typ. 120 $\mu$ A ( $V_{CC}=3.0V$ ,  $T_A=+25^{\circ}C$ ,  $f=32kHz$ ) Read Current
- **Optimized Array Blocking Architecture**
  - Two 4K-word Boot Blocks
  - Six 4K-word Parameter Blocks
  - Sixty-three 32K-word Main Blocks
  - Bottom Boot Location
- **Extended Cycling Capability**
  - Minimum 100,000 Block Erase Cycles
- **Enhanced Automated Suspend Options**
  - Word Write Suspend to Read
  - Block Erase Suspend to Word Write
  - Block Erase Suspend to Read
- **Enhanced Data Protection Features**
  - Absolute Protection with  $V_{CCW} \leq V_{CCWLK}$
  - Block Erase, Full Chip Erase, Word Write and Lock-Bit Configuration Lockout during Power Transitions
  - Block Locking with Command and WP#
  - Permanent Locking
- **Automated Block Erase, Full Chip Erase, Word Write and Lock-Bit Configuration**
  - Command User Interface (CUI)
  - Status Register (SR)
- **SRAM-Compatible Write Interface**
- **Chip-Size Packaging**
  - 0.75mm pitch 48-Ball CSP
- **ETOX<sup>TM</sup>\* Nonvolatile Flash Technology**
- **CMOS Process (P-type silicon substrate)**
- **Not designed or rated as radiation hardened**

SHARP's LH28F320BJHG-PBTLZ2 Flash memory is a high-density, low-cost, nonvolatile, read/write storage solution for a wide range of applications.

LH28F320BJHG-PBTLZ2 can operate at  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=2.7V-3.6V$  or  $11.7V-12.3V$ . Its low voltage operation capability realize battery life and suits for cellular phone application.

Its Boot, Parameter and Main-blocked architecture, low voltage and extended cycling provide for highly flexible component suitable for portable terminals and personal computers. Its enhanced suspend capabilities provide for an ideal solution for code + data storage applications.

For secure code storage applications, such as networking, where code is either directly executed out of flash or downloaded to DRAM, the LH28F320BJHG-PBTLZ2 offers four levels of protection: absolute protection with  $V_{CCW} \leq V_{CCWLK}$ , selective hardware block locking or flexible software block locking. These alternatives give designers ultimate control of their code security needs.

The LH28F320BJHG-PBTLZ2 is manufactured on SHARP's 0.25 $\mu$ m ETOX<sup>TM</sup>\* process technology. It come in chip-size package: the 0.75mm pitch 48-ball CSP, ideal for board constrained applications.

\*ETOX is a trademark of Intel Corporation.

## 1 INTRODUCTION

This datasheet contains LH28F320BJHG-PBTLZ2 specifications. Section 1 provides a flash memory overview. Sections 2, 3, 4 and 5 describe the memory organization and functionality. Section 6 covers electrical specifications.

### 1.1 Features

Key enhancements of LH28F320BJHG-PBTLZ2 boot block Flash memory are:

- Single low voltage operation
- Low power consumption
- Enhanced Suspend Capabilities
- Boot Block Architecture

Please note following:

- $V_{CCWLK}$  has been lowered to 1.0V to support 2.7V-3.6V block erase, full chip erase, word write and lock-bit configuration operations. The  $V_{CCW}$  voltage transitions to GND is recommended for designs that switch  $V_{CCW}$  off during read operation.

### 1.2 Product Overview

The LH28F320BJHG-PBTLZ2 is a high-performance 32M-bit Boot Block Flash memory organized as 2M-word of 16 bits. The 2M-word of data is arranged in two 4K-word boot blocks, six 4K-word parameter blocks and sixty-three 32K-word main blocks which are individually erasable, lockable and unlockable in-system. The memory map is shown in Figure 3.

The dedicated  $V_{CCW}$  pin gives complete data protection when  $V_{CCW} \leq V_{CCWLK}$ .

A Command User Interface (CUI) serves as the interface between the system processor and internal operation of the device. A valid command sequence written to the CUI initiates device automation. An internal Write State Machine (WSM) automatically executes the algorithms and timings necessary for block erase, full chip erase, word write and lock-bit configuration operations.

A block erase operation erases one of the device's 32K-word blocks typically within 1.2s ( $3V V_{CC}$ ,  $3V V_{CCW}$ ), 4K-word blocks typically within 0.6s ( $3V V_{CC}$ ,  $3V V_{CCW}$ ) independent of other blocks. Each block can be independently erased minimum 100,000 times. Block erase suspend mode allows system software to suspend block erase to read or write data from any other block.

Writing memory data is performed in word increments of the device's 32K-word blocks typically within 33 $\mu$ s ( $3V V_{CC}$ ,  $3V V_{CCW}$ ), 4K-word blocks typically within 36 $\mu$ s ( $3V V_{CC}$ ,  $3V V_{CCW}$ ). Word write suspend mode enables the system to read data or execute code from any other flash memory array location.

Individual block locking uses a combination of bits, seventy-one block lock-bits, a permanent lock-bit and WP# pin, to lock and unlock blocks. Block lock-bits gate block erase, full chip erase and word write operations, while the permanent lock-bit gates block lock-bit modification and locked block alternation. Lock-bit configuration operations (Set Block Lock-Bit, Set Permanent Lock-Bit and Clear Block Lock-Bits commands) set and cleared lock-bits.

The status register indicates when the WSM's block erase, full chip erase, word write or lock-bit configuration operation is finished.

The access time is 90ns ( $t_{AVQV}$ ) over the operating temperature range (-40°C to +85°C) and  $V_{CC}$  supply voltage range of 2.7V-3.6V.

The Automatic Power Savings (APS) feature substantially reduces active current when the device is in static mode (addresses not switching). In APS mode, the typical  $I_{CCR}$  current is 4 $\mu$ A (CMOS) at 3.0V  $V_{CC}$ .

When CE# and RP# pins are at  $V_{CC}$ , the  $I_{CC}$  CMOS standby mode is enabled. When the RP# pin is at GND, reset mode is enabled which minimizes power consumption and provides write protection. A reset time ( $t_{PHQV}$ ) is required from RP# switching high until outputs are valid. Likewise, the device has a wake time ( $t_{PHEL}$ ) from RP#-high until writes to the CUI are recognized. With RP# at GND, the WSM is reset and the status register is cleared.

Please do not execute reprogramming "0" for the bit which has already been programmed "0". Overwrite operation may generate unerasable bit. In case of reprogramming "0" to the data which has been programmed "1".

- Program "0" for the bit in which you want to change data from "1" to "0".
- Program "1" for the bit which has already been programmed "0".

For example, changing data from "10111101" to "10111100" requires "11111110" programming.

## 1.3 Product Description

### 1.3.1 Package Pinout

LH28F320BJHG-PBTLZ2 Boot Block Flash memory is available in 0.75mm pitch 48-ball CSP package (see Figure 2).

### 1.3.2 Block Organization

This product features an asymmetrically-blocked architecture providing system memory integration. Each erase block can be erased independently of the others up to 100,000 times. For the address locations of the blocks, see the memory map in Figure 3.

**Boot Blocks:** The boot block is intended to replace a dedicated boot PROM in a microprocessor or microcontroller-based system. This boot block 4K words (4,096 words) features hardware controllable write-protection to protect the crucial microprocessor boot code from accidental modification. The protection of the boot block is controlled using a combination of the  $V_{CCW}$ , RP#, WP# pins and block lock-bit.

**Parameter Blocks:** The boot block architecture includes parameter blocks to facilitate storage of frequently update small parameters that would normally require an EEPROM. By using software techniques, the word-rewrite functionality of EEPROMs can be emulated. Each boot block component contains six parameter blocks of 4K words (4,096 words) each. The protection of the parameter block is controlled using a combination of the  $V_{CCW}$ , RP# and block lock-bit.

**Main Blocks:** The remainder is divided into main blocks for data or code storage. Each 32M-bit device contains sixty-three 32K words (32,768 words) blocks. The protection of the main block is controlled using a combination of the  $V_{CCW}$ , RP# and block lock-bit.

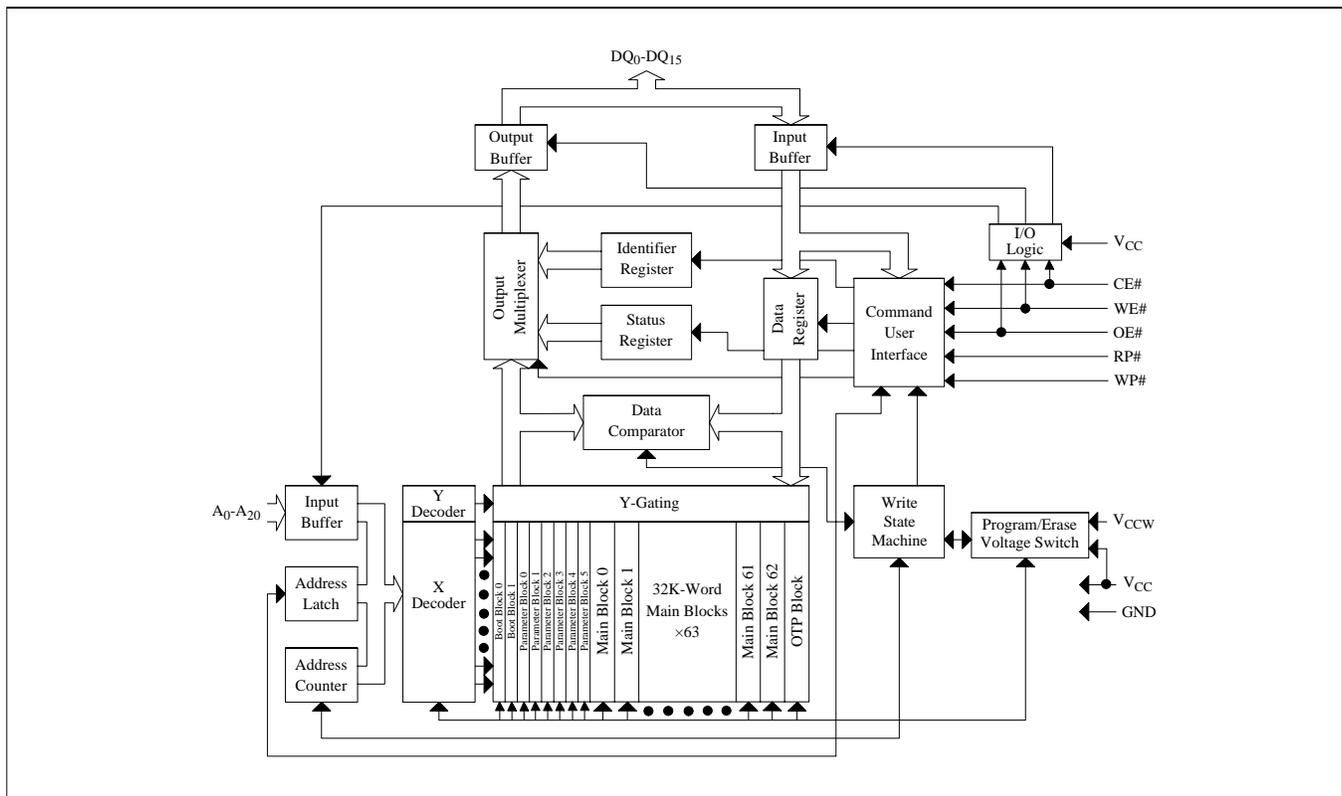


Figure 1. Block Diagram

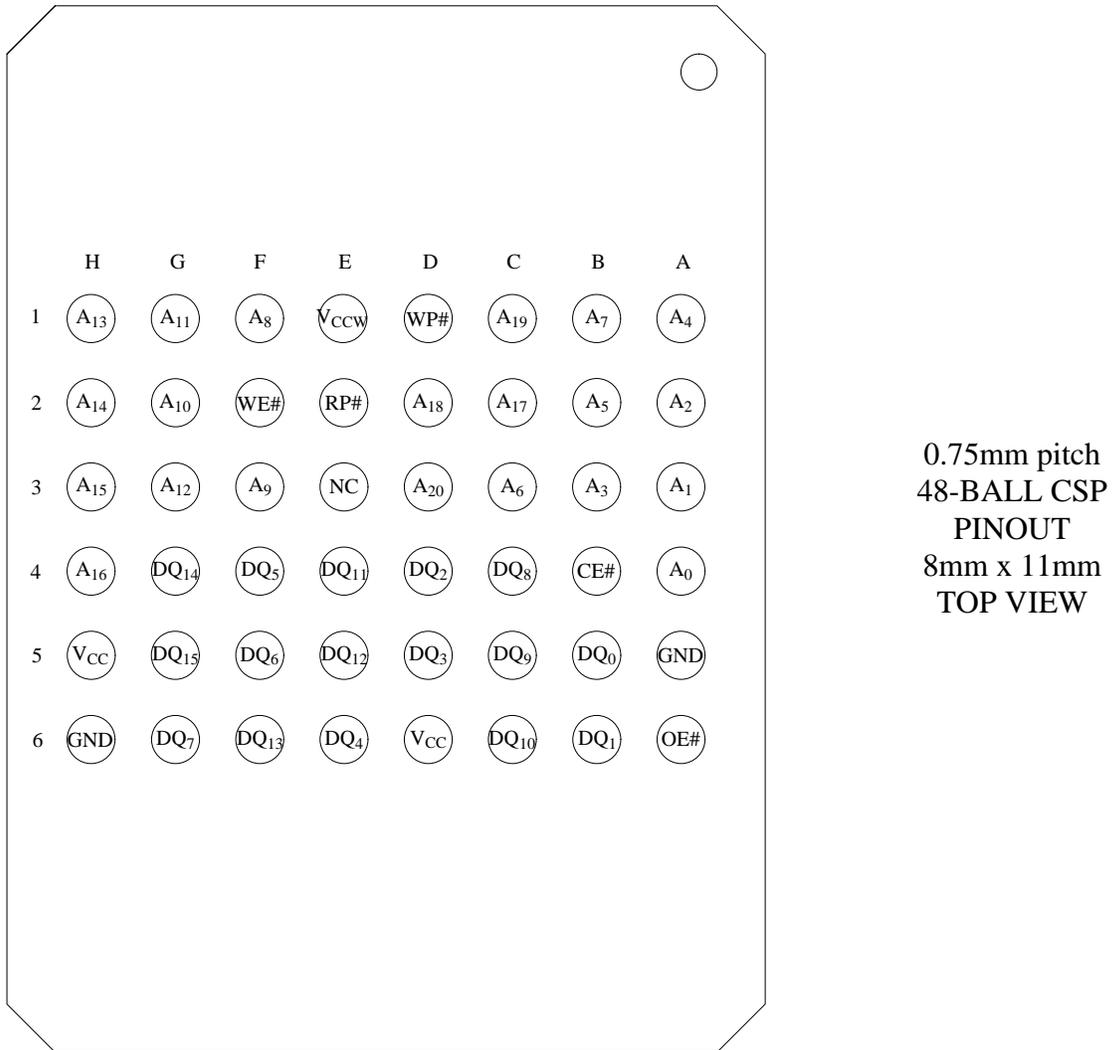


Figure 2. 0.75mm pitch CSP 48-Ball Pinout

Table 1. Pin Descriptions

Symbol	Type	Name and Function
A <sub>0</sub> -A <sub>20</sub>	INPUT	ADDRESS INPUTS: Inputs for addresses during read and write operations. Addresses are internally latched during a write cycle. A <sub>15</sub> -A <sub>20</sub> : Main Block Address. A <sub>12</sub> -A <sub>20</sub> : Boot and Parameter Block Address.
DQ <sub>0</sub> -DQ <sub>15</sub>	INPUT/ OUTPUT	DATA INPUT/OUTPUTS: Inputs data and commands during CUI write cycles; outputs data during memory array, status register and identifier code read cycles. Data pins float to high-impedance when the chip is deselected or outputs are disabled. Data is internally latched during a write cycle.
CE#	INPUT	CHIP ENABLE: Activates the device's control logic, input buffers, decoders and sense amplifiers. CE#-high deselects the device and reduces power consumption to standby levels.
RP#	INPUT	RESET: Resets the device internal automation. RP#-high enables normal operation. When driven low, RP# inhibits write operations which provides data protection during power transitions. Exit from reset mode sets the device to read array mode. RP# must be V <sub>IL</sub> during power-up.
OE#	INPUT	OUTPUT ENABLE: Gates the device's outputs during a read cycle.
WE#	INPUT	WRITE ENABLE: Controls writes to the CUI and array blocks. Addresses and data are latched on the rising edge of the WE# pulse.
WP#	INPUT	WRITE PROTECT: When WP# is V <sub>IL</sub> , boot blocks cannot be written or erased. When WP# is V <sub>IH</sub> , locked boot blocks can not be written or erased. WP# is not affected parameter and main blocks.
V <sub>CCW</sub>	SUPPLY	BLOCK ERASE, FULL CHIP ERASE, WORD WRITE OR LOCK-BIT CONFIGURATION POWER SUPPLY: For erasing array blocks, writing words or configuring lock-bits. With V <sub>CCW</sub> ≤ V <sub>CCWLK</sub> , memory contents cannot be altered. Block erase, full chip erase, word write and lock-bit configuration with an invalid V <sub>CCW</sub> (see 6.2.3 DC Characteristics) produce spurious results and should not be attempted. Applying 12V±0.3V to V <sub>CCW</sub> during erase/write can only be done for a maximum of 1000 cycles on each block. V <sub>CCW</sub> may be connected to 12V±0.3V for a total of 80 hours maximum.
V <sub>CC</sub>	SUPPLY	DEVICE POWER SUPPLY: Do not float any power pins. With V <sub>CC</sub> ≤ V <sub>LKO</sub> , all write attempts to the flash memory are inhibited. Device operations at invalid V <sub>CC</sub> voltage (see 6.2.3 DC Characteristics) produce spurious results and should not be attempted.
GND	SUPPLY	GROUND: Do not float any ground pins.
NC		NO CONNECT: Lead is not internal connected; it may be driven or floated.

## Bottom Boot

[A<sub>20</sub>-A<sub>0</sub>]

0FFFFF	32K-word Main Block	30
0F8000		
0F7FFF	32K-word Main Block	29
0F0000		
0EFFFF	32K-word Main Block	28
0E8000		
0E7FFF	32K-word Main Block	27
0E0000		
0DFFFF	32K-word Main Block	26
0D8000		
0D7FFF	32K-word Main Block	25
0D0000		
0CFFFF	32K-word Main Block	24
0C8000		
0C7FFF	32K-word Main Block	23
0C0000		
0BFFFF	32K-word Main Block	22
0B8000		
0B7FFF	32K-word Main Block	21
0B0000		
0AFFFF	32K-word Main Block	20
0A8000		
0A7FFF	32K-word Main Block	19
0A0000		
09FFFF	32K-word Main Block	18
098000		
097FFF	32K-word Main Block	17
090000		
08FFFF	32K-word Main Block	16
088000		
087FFF	32K-word Main Block	15
080000		
07FFFF	32K-word Main Block	14
078000		
077FFF	32K-word Main Block	13
070000		
06FFFF	32K-word Main Block	12
068000		
067FFF	32K-word Main Block	11
060000		
05FFFF	32K-word Main Block	10
058000		
057FFF	32K-word Main Block	9
050000		
04FFFF	32K-word Main Block	8
048000		
047FFF	32K-word Main Block	7
040000		
03FFFF	32K-word Main Block	6
038000		
037FFF	32K-word Main Block	5
030000		
02FFFF	32K-word Main Block	4
028000		
027FFF	32K-word Main Block	3
020000		
01FFFF	32K-word Main Block	2
018000		
017FFF	32K-word Main Block	1
010000		
00FFFF	32K-word Main Block	0
008000		
007FFF	4K-word Parameter Block	5
007000		
006FFF	4K-word Parameter Block	4
006000		
005FFF	4K-word Parameter Block	3
005000		
004FFF	4K-word Parameter Block	2
004000		
003FFF	4K-word Parameter Block	1
003000		
002FFF	4K-word Parameter Block	0
002000		
001FFF	4K-word Boot Block	1
001000		
000FFF	4K-word Boot Block	0
000000		

[A<sub>20</sub>-A<sub>0</sub>]

1FFFFF	32K-word Main Block	62
1F8000		
1F7FFF	32K-word Main Block	61
1F0000		
1EFFFF	32K-word Main Block	60
1E8000		
1E7FFF	32K-word Main Block	59
1E0000		
1DFFFF	32K-word Main Block	58
1D8000		
1D7FFF	32K-word Main Block	57
1D0000		
1CFFFF	32K-word Main Block	56
1C8000		
1C7FFF	32K-word Main Block	55
1C0000		
1BFFFF	32K-word Main Block	54
1B8000		
1B7FFF	32K-word Main Block	53
1B0000		
1AFFFF	32K-word Main Block	52
1A8000		
1A7FFF	32K-word Main Block	51
1A0000		
19FFFF	32K-word Main Block	50
198000		
197FFF	32K-word Main Block	49
190000		
18FFFF	32K-word Main Block	48
188000		
187FFF	32K-word Main Block	47
180000		
17FFFF	32K-word Main Block	46
178000		
177FFF	32K-word Main Block	45
170000		
16FFFF	32K-word Main Block	44
168000		
167FFF	32K-word Main Block	43
160000		
15FFFF	32K-word Main Block	42
158000		
157FFF	32K-word Main Block	41
150000		
14FFFF	32K-word Main Block	40
148000		
147FFF	32K-word Main Block	39
140000		
13FFFF	32K-word Main Block	38
138000		
137FFF	32K-word Main Block	37
130000		
12FFFF	32K-word Main Block	36
128000		
127FFF	32K-word Main Block	35
120000		
11FFFF	32K-word Main Block	34
118000		
117FFF	32K-word Main Block	33
110000		
10FFFF	32K-word Main Block	32
108000		
107FFF	32K-word Main Block	31
100000		

Figure 3. Memory Map

## 2 PRINCIPLES OF OPERATION

The LH28F320BJHG-PBTLZ2 flash memory includes an on-chip WSM to manage block erase, full chip erase, word write and lock-bit configuration functions. It allows for: fixed power supplies during block erase, full chip erase, word write and lock-bit configuration, and minimal processor overhead with RAM-like interface timings.

After initial device power-up or return from reset mode (see section 3 Bus Operations), the device defaults to read array mode. Manipulation of external memory control pins allow array read, standby and output disable operations.

Status register and identifier codes can be accessed through the CUI independent of the  $V_{CCW}$  voltage. High voltage on  $V_{CCW}$  enables successful block erase, full chip erase, word write and lock-bit configurations. All functions associated with altering memory contents—block erase, full chip erase, word write, lock-bit configuration, status and identifier codes—are accessed via the CUI and verified through the status register.

Commands are written using standard microprocessor write timings. The CUI contents serve as input to the WSM, which controls the block erase, full chip erase, word write and lock-bit configuration. The internal algorithms are regulated by the WSM, including pulse repetition, internal verification and margining of data. Addresses and data are internally latched during write cycles. Writing the appropriate command outputs array data, accesses the identifier codes or outputs status register data.

Interface software that initiates and polls progress of block erase, full chip erase, word write and lock-bit configuration can be stored in any block. This code is copied to and executed from system RAM during flash memory updates. After successful completion, reads are again possible via the Read Array command. Block erase suspend allows system software to suspend a block erase to read/write data from/to blocks other than that which is suspend. Word write suspend allows system software to suspend a word write to read data from any other flash memory array location.

### 2.1 Data Protection

When  $V_{CCW} \leq V_{CCWLK}$ , memory contents cannot be altered. The CUI, with two-step block erase, full chip erase, word write or lock-bit configuration command sequences, provides protection from unwanted operations even when high voltage is applied to  $V_{CCW}$ . All write functions are disabled when  $V_{CC}$  is below the write lockout voltage  $V_{LKO}$  or when  $RP\#$  is at  $V_{IL}$ . The device's block locking capability provides additional protection from inadvertent code or data alteration by gating block erase, full chip erase and word write operations. Refer to Table 5 for write protection alternatives.

### 3 BUS OPERATION

The local CPU reads and writes flash memory in-system. All bus cycles to or from the flash memory conform to standard microprocessor bus cycles.

#### 3.1 Read

Information can be read from any block, identifier codes or status register independent of the  $V_{CCW}$  voltage. RP# can be at  $V_{IH}$ .

The first task is to write the appropriate read mode command (Read Array, Read Identifier Codes or Read Status Register) to the CUI. Upon initial device power-up or after exit from reset mode, the device automatically resets to read array mode. Five control pins dictate the data flow in and out of the component: CE#, OE#, WE#, RP# and WP#. CE# and OE# must be driven active to obtain data at the outputs. CE# is the device selection control, and when active enables the selected memory device. OE# is the data output (DQ<sub>0</sub>-DQ<sub>15</sub>) control and when active drives the selected memory data onto the I/O bus. WE# must be at  $V_{IH}$ , RP# must be at  $V_{IH}$ , and WP# must be at  $V_{IL}$  or  $V_{IH}$ . Figure 16 illustrates read cycle.

#### 3.2 Output Disable

With OE# at a logic-high level ( $V_{IH}$ ), the device outputs are disabled. Output pins (DQ<sub>0</sub>-DQ<sub>15</sub>) are placed in a high-impedance state.

#### 3.3 Standby

CE# at a logic-high level ( $V_{IH}$ ) places the device in standby mode which substantially reduces device power consumption. DQ<sub>0</sub>-DQ<sub>15</sub> outputs are placed in a high-impedance state independent of OE#. If deselected during block erase, full chip erase, word write or lock-bit configuration, the device continues functioning, and consuming active power until the operation completes.

#### 3.4 Reset

RP# at  $V_{IL}$  initiates the reset mode.

In read modes, RP#-low deselected the memory, places output drivers in a high-impedance state and turns off all internal circuits. RP# must be held low for a minimum of 100ns. Time  $t_{PHQV}$  is required after return from reset mode until initial memory access outputs are valid. After this wake-up interval, normal operation is restored. The CUI is reset to read array mode and status register is set to 80H.

During block erase, full chip erase, word write or lock-bit configuration modes, RP#-low will abort the operation. SR.7 remains "0" until the reset operation is complete. Memory contents being altered are no longer valid; the data may be partially erased or written. Time  $t_{PHWL}$  is required after RP# goes to logic-high ( $V_{IH}$ ) before another command can be written.

As with any automated device, it is important to assert RP# during system reset. When the system comes out of reset, it expects to read from the flash memory. Automated flash memories provide status information when accessed during block erase, full chip erase, word write or lock-bit configuration modes. If a CPU reset occurs with no flash memory reset, proper CPU initialization may not occur because the flash memory may be providing status information instead of array data. SHARP's flash memories allow proper CPU initialization following a system reset through the use of the RP# input. In this application, RP# is controlled by the same RESET# signal that resets the system CPU.

#### 3.5 Read Identifier Codes

The read identifier codes operation outputs the manufacturer code, device code, block lock configuration codes for each block and the permanent lock configuration code (see Figure 4). Using the manufacturer and device codes, the system CPU can automatically match the device with its proper algorithms. The block lock and permanent lock configuration codes identify locked and unlocked blocks and permanent lock-bit setting.

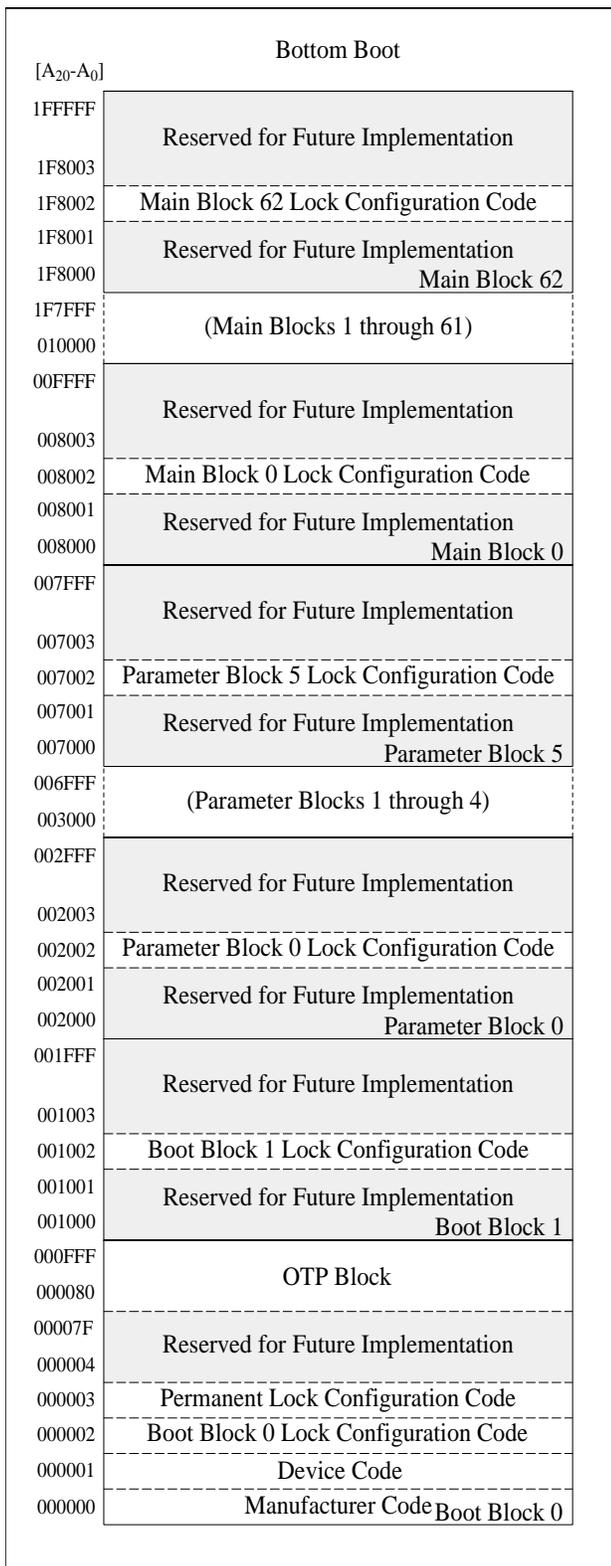


Figure 4. Device Identifier Code Memory Map

### 3.6 OTP(One Time Program) Block

The OTP block is a special block that can not be erased. The block is divided into two parts. One is a factory program area where a unique number can be written according to customer requirements in SHARP factory. This factory program area is "READ ONLY" (Already locked). The other is a customer program area that can be used by customers. This customer program area can be locked. After locking, this customer program area is protected permanently.

The OTP block is read in Configuration Read Mode by writing Read Identifier Codes command(90H). To return to Read Array Mode, write Read Array command(FFH).

The OTP block is programmed by writing OTP Program command(C0H). First write OTP Program command and then write data with address to the device (See Figure 5). If OTP program is failed, SR.4(WORD WRITE AND SET LOCK-BIT STATUS) bit is set to "1". And if this OTP block is locked, SR.1(DEVICE PROTECT STATUS) bit is set to "1" too.

The OTP block is also locked by writing OTP Program command(C0H). First write OTP Program command and then write data "FFFDH" with address "80H" to the device. Address "80H" of OTP block is OTP lock information. Bit 0 of address "80H" means factory program area lock status("1" is "NOT LOCKED", "0" is "LOCKED"). Bit 1 of address "80H" means customer program area lock status. The OTP lock information can not be cleared, after once it is set.

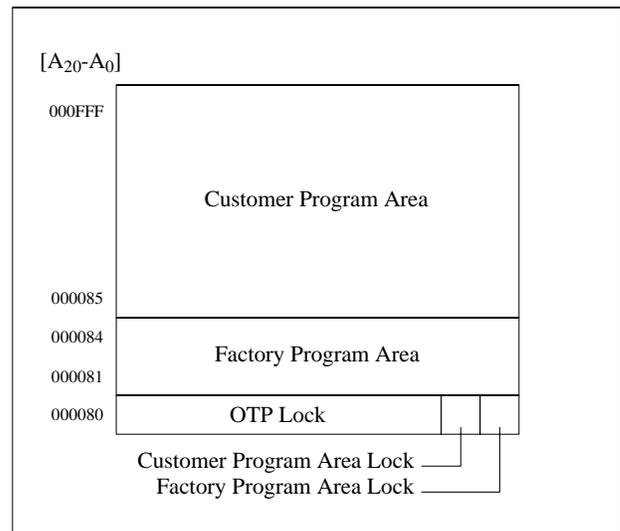


Figure 5. OTP Block Address Map

### 3.7 Write

Writing commands to the CUI enable reading of device data and identifier codes. They also control inspection and clearing of the status register. When  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ , the CUI additionally controls block erase, full chip erase, word write and lock-bit configuration.

The Block Erase command requires appropriate command data and an address within the block to be erased. The Full Chip Erase command requires appropriate command data and an address within the device. The Word Write command requires the command and address of the location to be written. Set Permanent and Block Lock-Bit commands require the command and address within the device (Permanent Lock) or block within the device (Block Lock) to be locked. The Clear Block Lock-Bits command requires the command and address within the device.

The CUI does not occupy an addressable memory location. It is written when WE# and CE# are active. The address and data needed to execute a command are latched on the rising edge of WE# or CE# (whichever goes high first). Standard microprocessor write timings are used. Figures 17 and 18 illustrate WE# and CE# controlled write operations.

## 4 COMMAND DEFINITIONS

When the  $V_{CCW}$  voltage  $\leq V_{CCWLK}$ , read operations from the status register, identifier codes, or blocks are enabled. Placing  $V_{CCWH1/2}$  on  $V_{CCW}$  enables successful block erase, full chip erase, word write and lock-bit configuration operations.

Device operations are selected by writing specific commands into the CUI. Table 3 defines these commands.

Table 2. Bus Operations<sup>(1,2)</sup>

Mode	Notes	RP#	CE#	OE#	WE#	Address	$V_{CCW}$	DQ <sub>0-15</sub>
Read	7	$V_{IH}$	$V_{IL}$	$V_{IL}$	$V_{IH}$	X	X	$D_{OUT}$
Output Disable		$V_{IH}$	$V_{IL}$	$V_{IH}$	$V_{IH}$	X	X	High Z
Standby		$V_{IH}$	$V_{IH}$	X	X	X	X	High Z
Reset	3	$V_{IL}$	X	X	X	X	X	High Z
Read Identifier Codes	7	$V_{IH}$	$V_{IL}$	$V_{IL}$	$V_{IH}$	See Figure 4, 5	X	Note 4
Write	5,6,7	$V_{IH}$	$V_{IL}$	$V_{IH}$	$V_{IL}$	X	X	$D_{IN}$

#### NOTES:

1. Refer to DC Characteristics. When  $V_{CCW} \leq V_{CCWLK}$ , memory contents can be read, but not altered.
2. X can be  $V_{IL}$  or  $V_{IH}$  for control pins and addresses, and  $V_{CCWLK}$  or  $V_{CCWH1/2}$  for  $V_{CCW}$ . See DC Characteristics for  $V_{CCWLK}$  voltages.
3. RP# at  $GND \pm 0.2V$  ensures the lowest power consumption.
4. See Section 4.2 for read identifier code data.
5. Command writes involving block erase, full chip erase, word write or lock-bit configuration are reliably executed when  $V_{CCW}=V_{CCWH1/2}$  and  $V_{CC}=2.7V-3.6V$ .
6. Refer to Table 3 for valid  $D_{IN}$  during a write operation.
7. Never hold OE# low and WE# low at the same timing.

Table 3. Command Definitions<sup>(10)</sup>

Command	Bus Cycles Req'd.	Notes	First Bus Cycle			Second Bus Cycle		
			Oper <sup>(1)</sup>	Addr <sup>(2)</sup>	Data <sup>(3)</sup>	Oper <sup>(1)</sup>	Addr <sup>(2)</sup>	Data <sup>(3)</sup>
Read Array/Reset	1		Write	X	FFH			
Read Identifier Codes	≥2	4	Write	X	90H	Read	IA	ID
Read Status Register	2		Write	X	70H	Read	X	SRD
Clear Status Register	1		Write	X	50H			
Block Erase	2	5	Write	X	20H	Write	BA	D0H
Full Chip Erase	2		Write	X	30H	Write	X	D0H
Word Write	2	5,6	Write	X	40H or 10H	Write	WA	WD
Block Erase and Word Write Suspend	1	5	Write	X	B0H			
Block Erase and Word Write Resume	1	5	Write	X	D0H			
Set Block Lock-Bit	2	8	Write	X	60H	Write	BA	01H
Clear Block Lock-Bits	2	7,8	Write	X	60H	Write	X	D0H
Set Permanent Lock-Bit	2	9	Write	X	60H	Write	X	F1H
OTP Program	2		Write	X	C0H	Write	OA	OD

## NOTES:

- BUS operations are defined in Table 2.
- X=Any valid address within the device.  
IA=Identifier Code Address: see Figure 4.  
BA=Address within the block being erased.  
WA=Address of memory location to be written.  
OA=Address of OTP block to be written : see Figure 5.
- ID=Data read from identifier codes.  
SRD=Data read from status register. See Table 6 for a description of the status register bits.  
WD=Data to be written at location WA. Data is latched on the rising edge of WE# or CE# (whichever goes high first).  
OD=Data to be written at location OA. Data is latched on the rising edge of WE# or CE# (whichever goes high first).
- Following the Read Identifier Codes command, read operations access manufacturer, device, block lock configuration and permanent lock configuration codes. See Section 4.2 for read identifier code data.
- If WP# is V<sub>IL</sub>, boot blocks are locked without block lock-bits state. If WP# is V<sub>IH</sub>, boot blocks are locked by block lock-bits. The parameter and main blocks are locked by block lock-bits without WP# state.
- Either 40H or 10H are recognized by the WSM as the word write setup.
- The clear block lock-bits operation simultaneously clears all block lock-bits.
- If the permanent lock-bit is set, Set Block Lock-Bit and Clear Block Lock-Bits commands can not be done.
- Once the permanent lock-bit is set, permanent lock-bit reset is unable.
- Commands other than those shown above are reserved by SHARP for future device implementations and should not be used.

## 4.1 Read Array Command

Upon initial device power-up and after exit from reset mode, the device defaults to read array mode. This operation is also initiated by writing the Read Array command. The device remains enabled for reads until another command is written. Once the internal WSM has started a block erase, full chip erase, word write or lock-bit configuration the device will not recognize the Read Array command until the WSM completes its operation unless the WSM is suspended via an Erase Suspend or Word Write Suspend command. The Read Array command functions independently of the  $V_{CCW}$  voltage and RP# can be  $V_{IH}$ .

## 4.2 Read Identifier Codes Command

The identifier code operation is initiated by writing the Read Identifier Codes command. Following the command write, read cycles from addresses shown in Figure 4 retrieve the manufacturer, device, block lock configuration and permanent lock configuration codes (see Table 4 for identifier code values). To terminate the operation, write another valid command. Like the Read Array command, the Read Identifier Codes command functions independently of the  $V_{CCW}$  voltage and RP# can be  $V_{IH}$ . Following the Read Identifier Codes command, the following information can be read:

Table 4. Identifier Codes

Code	Address [A <sub>20</sub> -A <sub>0</sub> ]	Data [DQ <sub>15</sub> -DQ <sub>0</sub> ]
Manufacture Code	00000H	00B0H
Device Code	00001H	00E3H
Block Lock Configuration	BA <sup>(1)</sup> +2	
•Block is Unlocked		DQ <sub>0</sub> =0
•Block is Locked		DQ <sub>0</sub> =1
•Reserved for Future Use		DQ <sub>1-7</sub>
Permanent Lock Configuration	00003H	
•Device is Unlocked		DQ <sub>0</sub> =0
•Device is Locked		DQ <sub>0</sub> =1
•Reserved for Future Use		DQ <sub>1-7</sub>

### NOTE:

1. BA selects the specific block lock configuration code to be read. See Figure 4 for the device identifier code memory map.

## 4.3 Read Status Register Command

The status register may be read to determine when a block erase, full chip erase, word write or lock-bit configuration is complete and whether the operation completed successfully. It may be read at any time by writing the Read Status Register command. After writing this command, all subsequent read operations output data from the status register until another valid command is written. The status register contents are latched on the falling edge of OE# or CE#, whichever occurs. OE# or CE# must toggle to  $V_{IH}$  before further reads to update the status register latch. The Read Status Register command functions independently of the  $V_{CCW}$  voltage. RP# can be  $V_{IH}$ .

## 4.4 Clear Status Register Command

Status register bits SR.5, SR.4, SR.3 or SR.1 are set to "1"s by the WSM and can only be reset by the Clear Status Register command. These bits indicate various failure conditions (see Table 6). By allowing system software to reset these bits, several operations (such as cumulatively erasing multiple blocks or writing several words in sequence) may be performed. The status register may be polled to determine if an error occurred during the sequence.

To clear the status register, the Clear Status Register command (50H) is written. It functions independently of the applied  $V_{CCW}$  Voltage. RP# can be  $V_{IH}$ . This command is not functional during block erase or word write suspend modes.

## 4.5 Block Erase Command

Erase is executed one block at a time and initiated by a two-cycle command. A block erase setup is first written, followed by an block erase confirm. This command sequence requires appropriate sequencing and an address within the block to be erased (erase changes all block data to FFFFH). Block preconditioning, erase, and verify are handled internally by the WSM (invisible to the system). After the two-cycle block erase sequence is written, the device automatically outputs status register data when read (see Figure 6). The CPU can detect block erase completion by analyzing the status register bit SR.7.

When the block erase is complete, status register bit SR.5 should be checked. If a block erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read status register mode until a new command is issued.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Block Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable block erasure can only occur when  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ . In the absence of this high voltage, block contents are protected against erasure. If block erase is attempted while  $V_{CCW}\leq V_{CCWLK}$ , SR.3 and SR.5 will be set to "1". Successful block erase requires for boot blocks that WP# is  $V_{IH}$  and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If block erase is attempted when the excepting above conditions, SR.1 and SR.5 will be set to "1".

## 4.6 Full Chip Erase Command

This command followed by a confirm command erases all of the unlocked blocks. A full chip erase setup (30H) is first written, followed by a full chip erase confirm (D0H). After a confirm command is written, device erases the all unlocked blocks block by block. This command sequence requires appropriate sequencing. Block preconditioning, erase and verify are handled internally by the WSM (invisible to the system). After the two-cycle full chip erase sequence is written, the device automatically outputs status register data when read (see Figure 7). The CPU can detect full chip erase completion by analyzing the output data of the status register bit SR.7.

When the full chip erase is complete, status register bit SR.5 should be checked. If erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read

status register mode until a new command is issued. If error is detected on a block during full chip erase operation, WSM stops erasing. Full chip erase operation start from lower address block, finish the higher address block. Full chip erase can not be suspended.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Full Chip Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable full chip erasure can only occur when  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ . In the absence of this high voltage, block contents are protected against erasure. If full chip erase is attempted while  $V_{CCW}\leq V_{CCWLK}$ , SR.3 and SR.5 will be set to "1". Successful full chip erase requires for boot blocks that WP# is  $V_{IH}$  and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If all blocks are locked, SR.1 and SR.5 will be set to "1".

## 4.7 Word Write Command

Word write is executed by a two-cycle command sequence. Word write setup (standard 40H or alternate 10H) is written, followed by a second write that specifies the address and data (latched on the rising edge of WE#). The WSM then takes over, controlling the word write and write verify algorithms internally. After the word write sequence is written, the device automatically outputs status register data when read (see Figure 8). The CPU can detect the completion of the word write event by analyzing the status register bit SR.7.

When word write is complete, status register bit SR.4 should be checked. If word write error is detected, the status register should be cleared. The internal WSM verify only detects errors for "1"s that do not successfully write to "0"s. The CUI remains in read status register mode until it receives another command.

Reliable word writes can only occur when  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ . In the absence of this high voltage, memory contents are protected against word writes. If word write is attempted while  $V_{CCW}\leq V_{CCWLK}$ , status register bits SR.3 and SR.4 will be set to "1". Successful word write requires for boot blocks that WP# is  $V_{IH}$  and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If word write is attempted when the excepting above conditions, SR.1 and SR.4 will be set to "1".

#### 4.8 Block Erase Suspend Command

The Block Erase Suspend command allows block-erase interruption to read or word write data in another block of memory. Once the block erase process starts, writing the Block Erase Suspend command requests that the WSM suspend the block erase sequence at a predetermined point in the algorithm. The device outputs status register data when read after the Block Erase Suspend command is written. Polling status register bits SR.7 and SR.6 can determine when the block erase operation has been suspended (both will be set to "1"). Specification  $t_{WHR12}$  defines the block erase suspend latency.

When Block Erase Suspend command write to the CUI, if block erase was finished, the device places read array mode. Therefore, after Block Erase Suspend command write to the CUI, Read Status Register command (70H) has to write to CUI, then status register bit SR.6 should be checked for places the device in suspend mode.

At this point, a Read Array command can be written to read data from blocks other than that which is suspended. A Word Write command sequence can also be issued during erase suspend to program data in other blocks. Using the Word Write Suspend command (see Section 4.9), a word write operation can also be suspended. During a word write operation with block erase suspended, status register bit SR.7 will return to "0". However, SR.6 will remain "1" to indicate block erase suspend status.

The only other valid commands while block erase is suspended are Read Status Register and Block Erase Resume. After a Block Erase Resume command is written to the flash memory, the WSM will continue the block erase process. Status register bits SR.6 and SR.7 will automatically clear. After the Erase Resume command is written, the device automatically outputs status register data when read (see Figure 9).  $V_{CCW}$  must remain at  $V_{CCWH1/2}$  (the same  $V_{CCW}$  level used for block erase) while block erase is suspended. RP# must also remain at  $V_{IH}$ . WP# must also remain at  $V_{IL}$  or  $V_{IH}$  (the same WP# level used for block erase). Block erase cannot resume until word write operations initiated during block erase suspend have completed.

If the time between writing the Block Erase Resume command and writing the Block Erase Suspend command is shorter than  $t_{ERES}$  and both commands are written repeatedly, a longer time is required than standard block erase until the completion of the operation.

#### 4.9 Word Write Suspend Command

The Word Write Suspend command allows word write interruption to read data in other flash memory locations. Once the word write process starts, writing the Word Write Suspend command requests that the WSM suspend the Word write sequence at a predetermined point in the algorithm. The device continues to output status register data when read after the Word Write Suspend command is written. Polling status register bits SR.7 and SR.2 can determine when the word write operation has been suspended (both will be set to "1"). Specification  $t_{WHR11}$  defines the word write suspend latency.

When Word Write Suspend command write to the CUI, if word write was finished, the device places read array mode. Therefore, after Word Write Suspend command write to the CUI, Read Status Register command (70H) has to write to CUI, then status register bit SR.2 should be checked for places the device in suspend mode.

At this point, a Read Array command can be written to read data from locations other than that which is suspended. The only other valid commands while word write is suspended are Read Status Register and Word Write Resume. After Word Write Resume command is written to the flash memory, the WSM will continue the word write process. Status register bits SR.2 and SR.7 will automatically clear. After the Word Write Resume command is written, the device automatically outputs status register data when read (see Figure 10).  $V_{CCW}$  must remain at  $V_{CCWH1/2}$  (the same  $V_{CCW}$  level used for word write) while in word write suspend mode. RP# must also remain at  $V_{IH}$ . WP# must also remain at  $V_{IL}$  or  $V_{IH}$  (the same WP# level used for word write).

If the time between writing the Word Write Resume command and writing the Word Write Suspend command is short and both commands are written repeatedly, a longer time is required than standard word write until the completion of the operation.

#### 4.10 Set Block and Permanent Lock-Bit Commands

A flexible block locking and unlocking scheme is enabled via a combination of block lock-bits, a permanent lock-bit and WP# pin. The block lock-bits and WP# pin gates program and erase operations while the permanent lock-bit gates block-lock bit modification. With the permanent lock-bit not set, individual block lock-bits can be set using the Set Block Lock-Bit command. The Set Permanent Lock-Bit command, sets the permanent lock-bit. After the permanent lock-bit is set, block lock-bits and locked block contents cannot be altered. See Table 5 for a summary of hardware and software write protection options.

Set block lock-bit and permanent lock-bit are executed by a two-cycle command sequence. The set block or permanent lock-bit setup along with appropriate block or device address is written followed by either the set block lock-bit confirm (and an address within the block to be locked) or the set permanent lock-bit confirm (and any device address). The WSM then controls the set lock-bit algorithm. After the sequence is written, the device automatically outputs status register data when read (see Figure 11). The CPU can detect the completion of the set lock-bit event by analyzing the status register bit SR.7.

When the set lock-bit operation is complete, status register bit SR.4 should be checked. If an error is detected, the status register should be cleared. The CUI will remain in read status register mode until a new command is issued.

This two-step sequence of set-up followed by execution ensures that lock-bits are not accidentally set. An invalid Set Block or Permanent Lock-Bit command will result in status register bits SR.4 and SR.5 being set to "1". Also, reliable operations occur only when  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ . In the absence of this high voltage, lock-bit contents are protected against alteration.

A successful set block lock-bit operation requires that the permanent lock-bit be cleared. If it is attempted with the permanent lock-bit set, SR.1 and SR.4 will be set to "1" and the operation will fail.

#### 4.11 Clear Block Lock-Bits Command

All set block lock-bits are cleared in parallel via the Clear Block Lock-Bits command. With the permanent lock-bit not set, block lock-bits can be cleared using only the Clear Block Lock-Bits command. If the permanent lock-bit is set, block lock-bits cannot be cleared. See Table 5 for a summary of hardware and software write protection options.

Clear block lock-bits operation is executed by a two-cycle command sequence. A clear block lock-bits setup is first written. After the command is written, the device automatically outputs status register data when read (see Figure 12). The CPU can detect completion of the clear block lock-bits event by analyzing the status register bit SR.7.

When the operation is complete, status register bit SR.5 should be checked. If a clear block lock-bit error is detected, the status register should be cleared. The CUI will remain in read status register mode until another command is issued.

This two-step sequence of set-up followed by execution ensures that block lock-bits are not accidentally cleared. An invalid Clear Block Lock-Bits command sequence will result in status register bits SR.4 and SR.5 being set to "1". Also, a reliable clear block lock-bits operation can only occur when  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ . If a clear block lock-bits operation is attempted while  $V_{CCW} \leq V_{CCWLK}$ , SR.3 and SR.5 will be set to "1". In the absence of this high voltage, the block lock-bits content are protected against alteration. A successful clear block lock-bits operation requires that the permanent lock-bit is not set. If it is attempted with the permanent lock-bit set, SR.1 and SR.5 will be set to "1" and the operation will fail.

If a clear block lock-bits operation is aborted due to  $V_{CCW}$  or  $V_{CC}$  transitioning out of valid range or RP# active transition, block lock-bit values are left in an undetermined state. A repeat of clear block lock-bits is required to initialize block lock-bit contents to known values. Once the permanent lock-bit is set, it cannot be cleared.

#### 4.12 OTP Program Command

OTP program is executed by a two-cycle command sequence. OTP program command(C0H) is written, followed by a second write cycle that specifies the address and data (latched on the rising edge of WE#). The WSM then takes over, controlling the OTP program and program verify algorithms internally. After the OTP program command sequence is completed, the device automatically outputs status register data when read (see Figure 13). The CPU can detect the completion of the OTP program by analyzing the status register bit SR.7.

When OTP program is completed, status register bit SR.4 should be checked. If OTP program error is detected, the status register should be cleared. The internal WSM verify only detects errors for "1"s that do not successfully program to "0"s. The CUI remains in read status register mode until it receives other commands.

Reliable OTP program can be executed only when  $V_{CC}=2.7V-3.6V$  and  $V_{CCW}=V_{CCWH1/2}$ . In the absence of this voltage, memory contents are protected against OTP programs. If OTP program is attempted while  $V_{CCW} \leq V_{CCWLK}$ , status register bits SR.3 and SR.4 is set to "1". If OTP write is attempted when the OTP Lock-bit is set, SR.1 and SR.4 is set to "1".

#### 4.13 Block Locking by the WP#

This Boot Block Flash memory architecture features two hardware-lockable boot blocks so that the kernel code for the system can be kept secure while other blocks are programmed or erased as necessary.

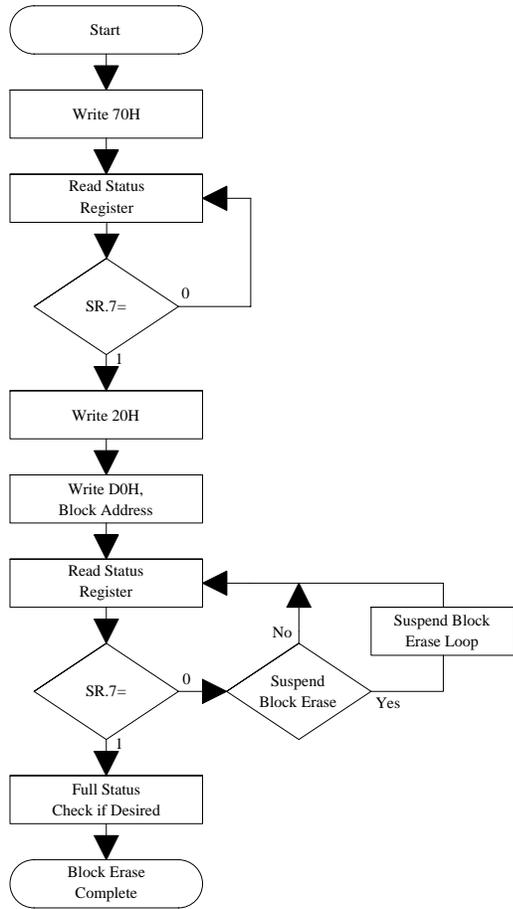
The lockable two boot blocks are locked when  $WP\#=V_{IL}$ ; any program or erase operation to a locked block will result in an error, which will be reflected in the status register. For top configuration, the top two boot blocks are lockable. For the bottom configuration, the bottom two boot blocks are lockable. If  $WP\#$  is  $V_{IH}$  and block lock-bit is not set, boot block can be programmed or erased normally (Unless  $V_{CCW}$  is below  $V_{CCWLK}$ ).  $WP\#$  is valid only two boot blocks, other blocks are not affected.

Table 5. Write Protection Alternatives

Operation	V <sub>CCWLK</sub>	RP#	Permanent Lock-Bit	Block Lock-bit	WP#	Effect
Block Erase or Word Write	≤V <sub>CCWLK</sub>	X	X	X	X	All Blocks Locked.
	>V <sub>CCWLK</sub>	V <sub>IL</sub>	X	X	X	All Blocks Locked.
		V <sub>IH</sub>	X	0	V <sub>IL</sub>	2 Boot Blocks Locked.
					V <sub>IH</sub>	Block Erase and Word Write Enabled.
		1	V <sub>IL</sub>	Block Erase and Word Write Disabled.		
V <sub>IH</sub>	Block Erase and Word Write Disabled.					
Full Chip Erase	≤V <sub>CCWLK</sub>	X	X	X	X	All Blocks Locked.
	>V <sub>CCWLK</sub>	V <sub>IL</sub>	X	X	X	All Blocks Locked.
		V <sub>IH</sub>	X	X	V <sub>IL</sub>	All Unlocked Blocks are Erased. 2 Boot Blocks and Locked Blocks are NOT Erased.
					V <sub>IH</sub>	All Unlocked Blocks are Erased, Locked Blocks are NOT Erased.
Set Block Lock-Bit	≤V <sub>CCWLK</sub>	X	X	X	X	Set Block Lock-Bit Disabled.
	>V <sub>CCWLK</sub>	V <sub>IL</sub>	X	X	X	Set Block Lock-Bit Disabled.
		V <sub>IH</sub>	0	X	X	Set Block Lock-Bit Enabled.
			1	X	X	Set Block Lock-Bit Disabled.
Clear Block Lock-Bits	≤V <sub>CCWLK</sub>	X	X	X	X	Clear Block Lock-Bits Disabled.
	>V <sub>CCWLK</sub>	V <sub>IL</sub>	X	X	X	Clear Block Lock-Bits Disabled.
		V <sub>IH</sub>	0	X	X	Clear Block Lock-Bits Enabled.
			1	X	X	Clear Block Lock-Bits Disabled.
Set Permanent Lock-Bit	≤V <sub>CCWLK</sub>	X	X	X	X	Set Permanent Lock-Bit Disabled.
	>V <sub>CCWLK</sub>	V <sub>IL</sub>	X	X	X	Set Permanent Lock-Bit Disabled.
		V <sub>IH</sub>	X	X	X	Set Permanent Lock-Bit Enabled.

Table 6. Status Register Definition

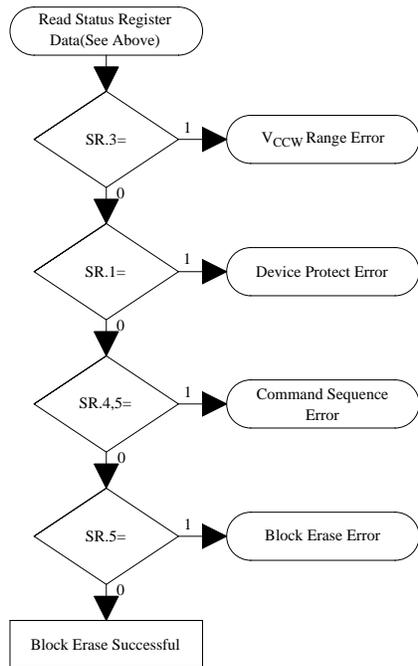
WSMS	BESS	ECBLBS	WWSLBS	VCCWS	WWSS	DPS	R
7	6	5	4	3	2	1	0
<p>SR.7 = WRITE STATE MACHINE STATUS (WSMS) 1 = Ready 0 = Busy</p> <p>SR.6 = BLOCK ERASE SUSPEND STATUS (BESS) 1 = Block Erase Suspended 0 = Block Erase in Progress/Completed</p> <p>SR.5 = ERASE AND CLEAR BLOCK LOCK-BITS STATUS (ECBLBS) 1 = Error in Block Erase, Full Chip Erase or Clear Block Lock-Bits 0 = Successful Block Erase, Full Chip Erase or Clear Block Lock-Bits</p> <p>SR.4 = WORD WRITE AND SET LOCK-BIT STATUS (WWSLBS) 1 = Error in Word Write or Set Block/Permanent Lock-Bit 0 = Successful Word Write or Set Block/Permanent Lock-Bit</p> <p>SR.3 = <math>V_{CCW}</math> STATUS (VCCWS) 1 = <math>V_{CCW}</math> Low Detect, Operation Abort 0 = <math>V_{CCW}</math> OK</p> <p>SR.2 = WORD WRITE SUSPEND STATUS (WWSS) 1 = Word Write Suspended 0 = Word Write in Progress/Completed</p> <p>SR.1 = DEVICE PROTECT STATUS (DPS) 1 = Block Lock-Bit, Permanent Lock-Bit and/or WP# Lock Detected, Operation Abort 0 = Unlock</p> <p>SR.0 = RESERVED FOR FUTURE ENHANCEMENTS (R)</p>				<p>NOTES:</p> <p>Check SR.7 to determine block erase, full chip erase, word write or lock-bit configuration completion. SR.6-0 are invalid while SR.7="0".</p> <p>If both SR.5 and SR.4 are "1"s after a block erase, full chip erase or lock-bit configuration attempt, an improper command sequence was entered.</p> <p>SR.3 does not provide a continuous indication of <math>V_{CCW}</math> level. The WSM interrogates and indicates the <math>V_{CCW}</math> level only after Block Erase, Full Chip Erase, Word Write or Lock-Bit Configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when <math>V_{CCW} \neq V_{CCWH1/2}</math>.</p> <p>SR.1 does not provide a continuous indication of permanent and block lock-bit and WP# values. The WSM interrogates the permanent lock-bit, block lock-bit and WP# only after Block Erase, Full Chip Erase, Word Write or Lock-Bit Configuration command sequences. It informs the system, depending on the attempted operation, if the block lock-bit is set, permanent lock-bit is set and/or WP# is <math>V_{IL}</math>. Reading the block lock and permanent lock configuration codes after writing the Read Identifier Codes command indicates permanent and block lock-bit status.</p> <p>SR.0 is reserved for future use and should be masked out when polling the status register.</p>			



Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Erase Setup	Data=20H Addr=X
Write	Erase Confirm	Data=D0H Addr=Within Block to be Erased
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Repeat for subsequent block erasures.  
Full status check can be done after each block erase or after a sequence of block erasures.  
Write FFH after the last operation to place device in read array mode.

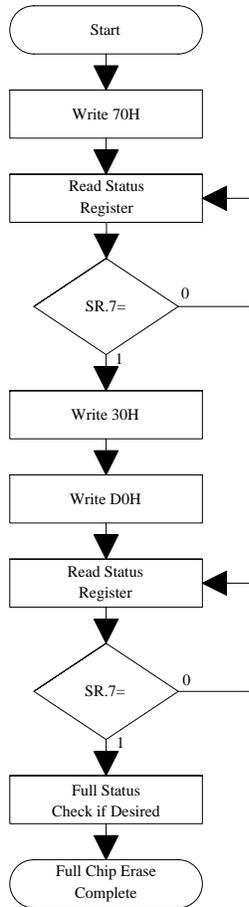
FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V <sub>CCW</sub> Error Detect
Standby		Check SR.1 1=Device Protect Detect
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.5 1=Block Erase Error

SR.5, SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register Command in cases where multiple blocks are erased before full status is checked.  
If error is detected, clear the Status Register before attempting retry or other error recovery.

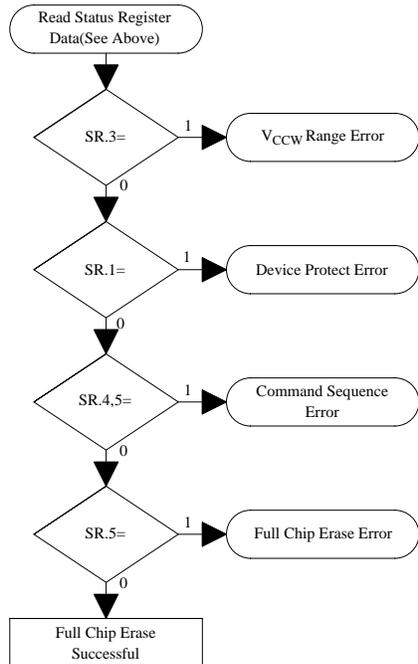
Figure 6. Automated Block Erase Flowchart



Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Full Chip Erase Setup	Data=30H Addr=X
Write	Full Chip Erase Confirm	Data=D0H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Full status check can be done after each full chip erase.  
Write FFH after the last operation to place device in read array mode.

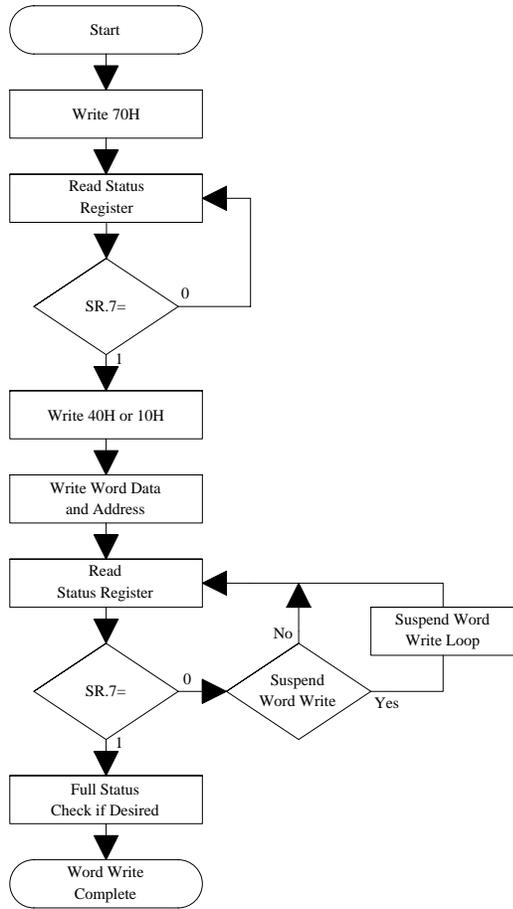
FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V <sub>CCW</sub> Error Detect
Standby		Check SR.1 1=Device Protect Detect (All Blocks are locked)
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.5 1=Full Chip Erase Error

SR.5, SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register Command in cases where multiple blocks are erased before full status is checked.  
If error is detected, clear the Status Register before attempting retry or other error recovery.

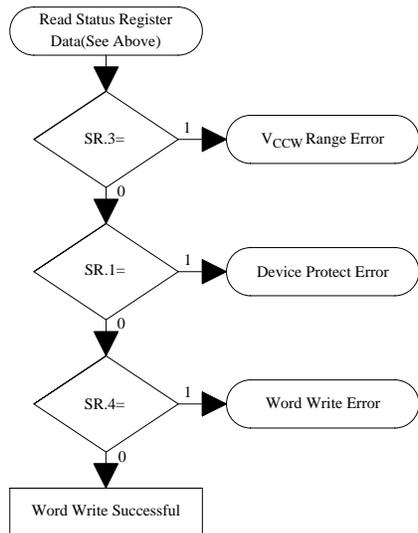
Figure 7. Automated Full Chip Erase Flowchart



Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Setup Word Write	Data=40H or 10H Addr=X
Write	Word Write	Data=Data to Be Written Addr=Location to Be Written
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Repeat for subsequent word writes.  
SR full status check can be done after each word write, or after a sequence of word writes.  
Write FFH after the last word write operation to place device in read array mode.

FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V <sub>CCW</sub> Error Detect
Standby		Check SR.1 1=Device Protect Detect
Standby		Check SR.4 1=Data Write Error

SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register command in cases where multiple locations are written before full status is checked.  
If error is detected, clear the Status Register before attempting retry or other error recovery.

Figure 8. Automated Word Write Flowchart

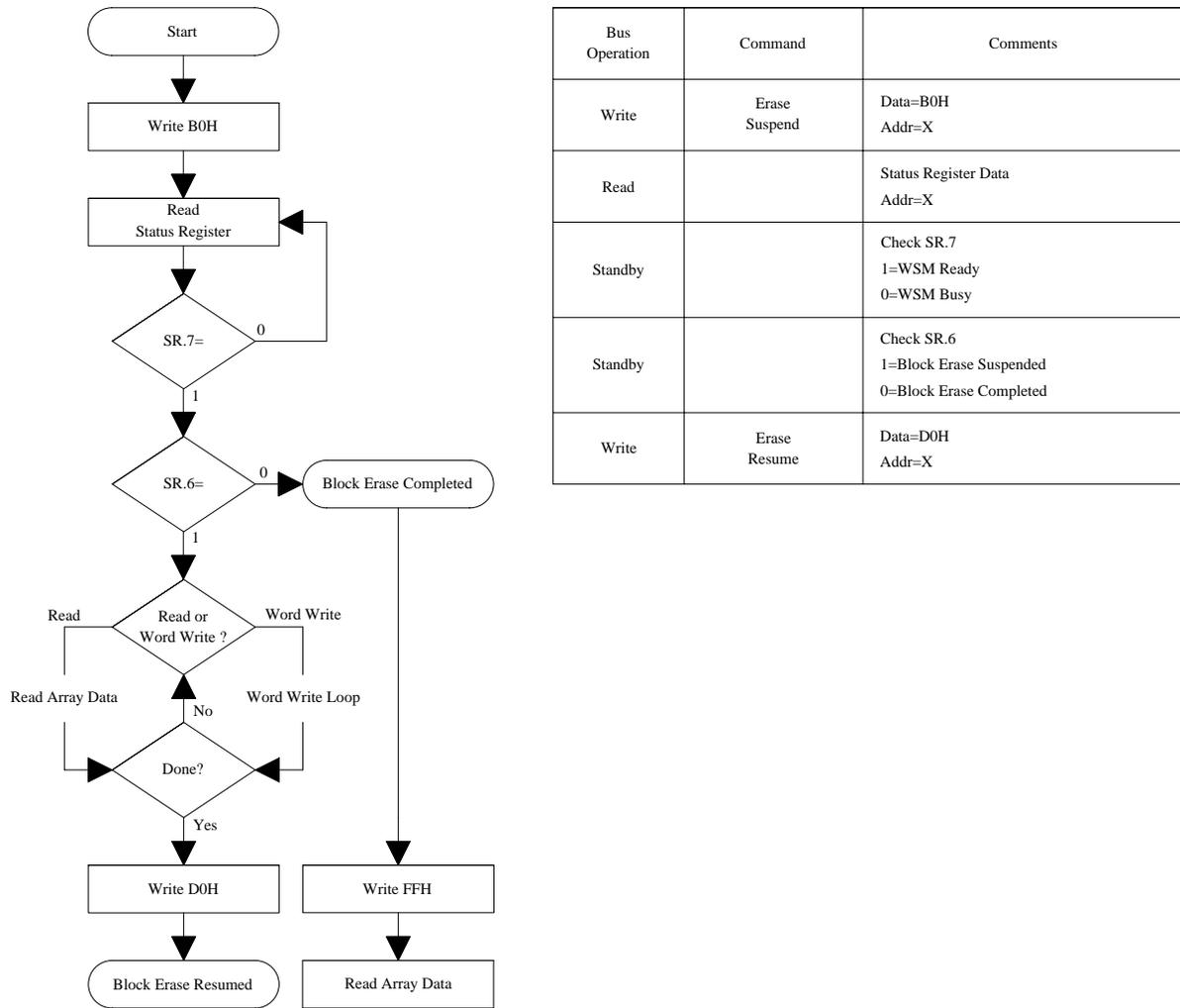
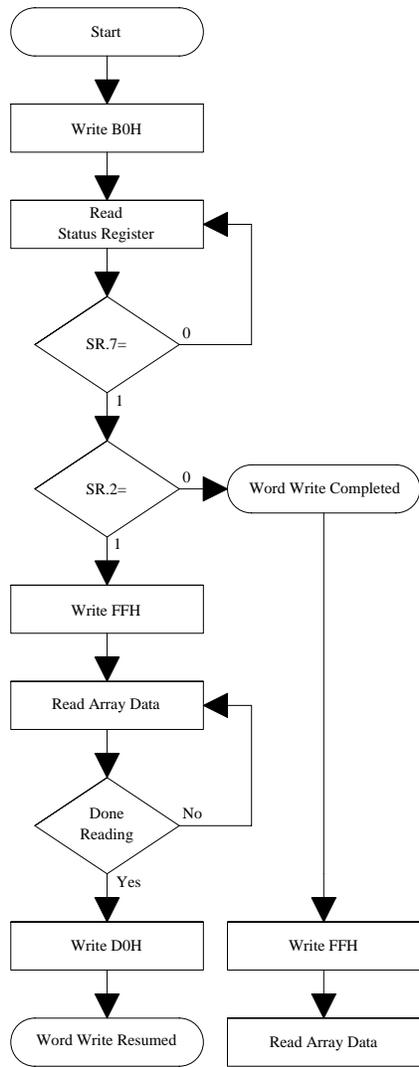
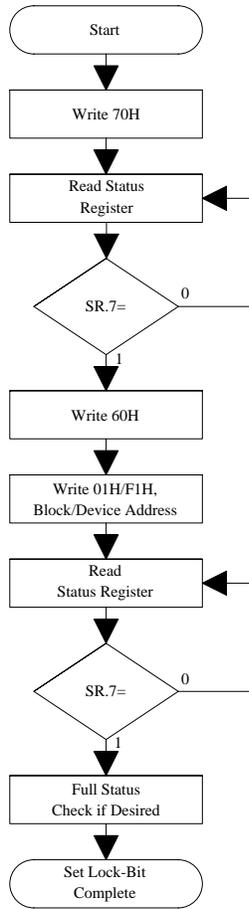


Figure 9. Block Erase Suspend/Resume Flowchart



Bus Operation	Command	Comments
Write	Word Write Suspend	Data=B0H Addr=X
Read		Status Register Data Addr=X
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Standby		Check SR.2 1=Word Write Suspended 0=Word Write Completed
Write	Read Array	Data=FFH Addr=X
Read		Read Array locations other than that being written.
Write	Word Write Resume	Data=D0H Addr=X

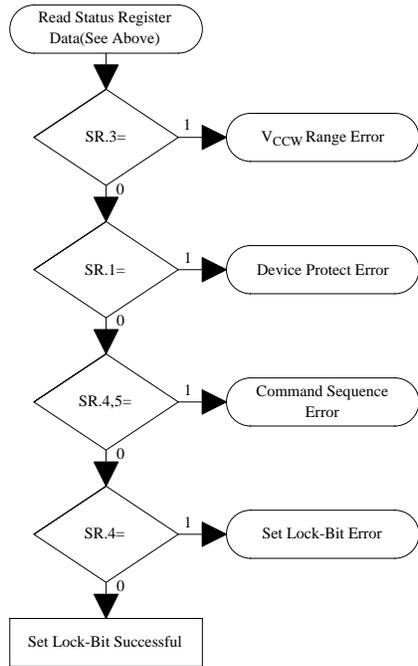
Figure 10. Word Write Suspend/Resume Flowchart



Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Set Block/Permanent Lock-Bit Setup	Data=60H Addr=X
Write	Set Block or Permanent Lock-Bit Confirm	Data=01H(Block), F1H(Permanent) Addr=Block Address(Block), Device Address(Permanent)
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Repeat for subsequent lock-bit set operations.  
Full status check can be done after each lock-bit set operation or after a sequence of lock-bit set operations.  
Write FFH after the last lock-bit set operation to place device in read array mode.

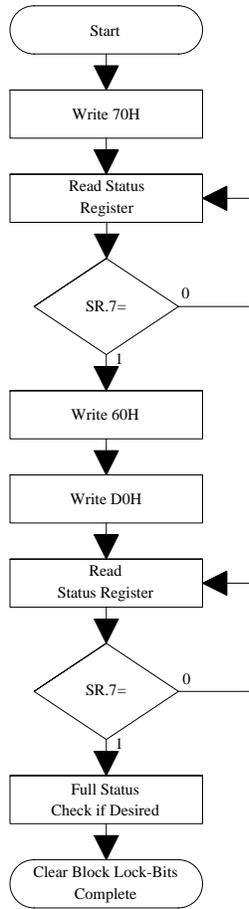
FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V <sub>CCW</sub> Error Detect
Standby		Check SR.1 1=Device Protect Detect Permanent Lock-Bit is Set (Set Block Lock-Bit Operation)
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.4 1=Set Lock-Bit Error

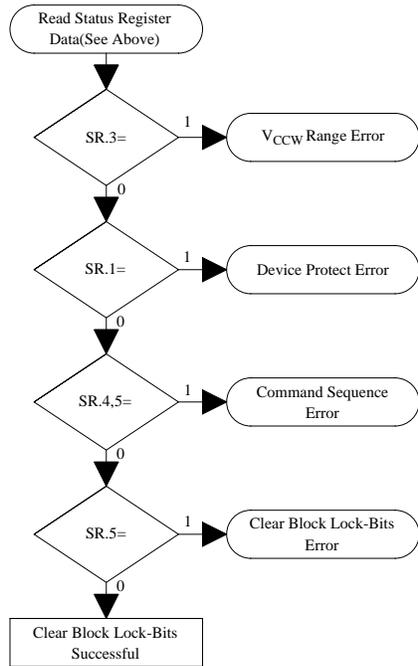
SR.5, SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register command in cases where multiple lock-bits are set before full status is checked.  
If error is detected, clear the Status Register before attempting retry or other error recovery.

Figure 11. Set Block and Permanent Lock-Bit Flowchart



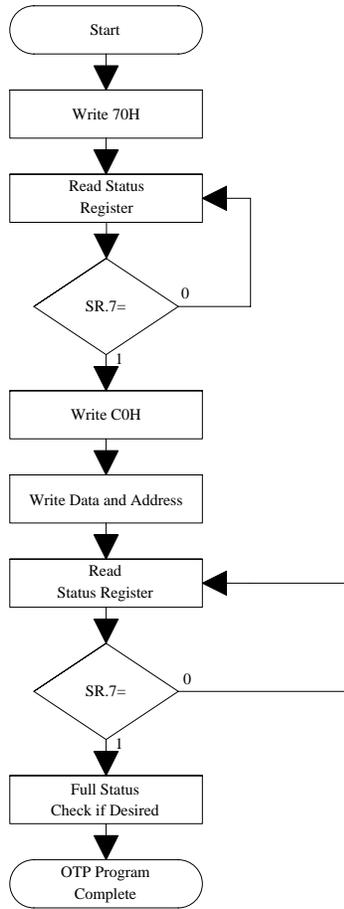
Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Clear Block Lock-Bits Setup	Data=60H Addr=X
Write	Clear Block Lock-Bits Confirm	Data=D0H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write FFH after the Clear Block Lock-Bits operation to place device in read array mode.		

FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V <sub>CCW</sub> Error Detect
Standby		Check SR.1 1=Device Protect Detect Permanent Lock-Bit is Set
Standby		Check SR.4,5 Both 1=Command Sequence Error
Standby		Check SR.5 1=Clear Block Lock-Bits Error
SR.5, SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register command. If error is detected, clear the Status Register before attempting retry or other error recovery.		

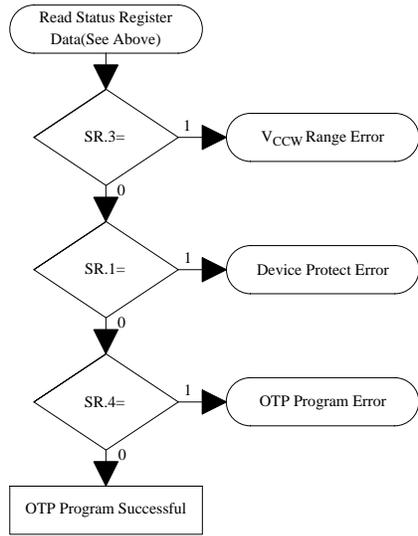
Figure 12. Clear Block Lock-Bits Flowchart



Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Setup OTP Program	Data=C0H Addr=X
Write	OTP Program	Data=Data to Be Written Addr=Location to Be Written
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Repeat for subsequent OTP programs.  
 SR full status check can be done after each OTP program, or after a sequence of OTP programs.  
 Write FFH after the last OTP program operation to place device in read array mode.

FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 1=V <sub>CCW</sub> Error Detect
Standby		Check SR.1 1=Device Protect Detect
Standby		Check SR.4 1=Data Write Error

SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register command in cases where multiple locations are written before full status is checked.  
 If error is detected, clear the Status Register before attempting retry or other error recovery.

Figure 13. Automated OTP Program Flowchart

## 5 DESIGN CONSIDERATIONS

### 5.1 Three-Line Output Control

The device will often be used in large memory arrays. SHARP provides three control inputs to accommodate multiple memory connections. Three-line control provides for:

- Lowest possible memory power dissipation.
- Complete assurance that data bus contention will not occur.

To use these control inputs efficiently, an address decoder should enable CE# while OE# should be connected to all memory devices and the system's READ# control line. This assures that only selected memory devices have active outputs while deselected memory devices are in standby mode. RP# should be connected to the system POWERGOOD signal to prevent unintended writes during system power transitions. POWERGOOD should also toggle during system reset.

### 5.2 Power Supply Decoupling

Flash memory power switching characteristics require careful device decoupling. System designers are interested in three supply current issues; standby current levels, active current levels and transient peaks produced by falling and rising edges of CE# and OE#. Transient current magnitudes depend on the device outputs' capacitive and inductive loading. Two-line control and proper decoupling capacitor selection will suppress transient voltage peaks. Each device should have a 0.1 $\mu$ F ceramic capacitor connected between its V<sub>CC</sub> and GND and between its V<sub>CCW</sub> and GND. These high-frequency, low inductance capacitors should be placed as close as possible to package leads. Additionally, for every eight devices, a 4.7 $\mu$ F electrolytic capacitor should be placed at the array's power supply connection between V<sub>CC</sub> and GND. The bulk capacitor will overcome voltage slumps caused by PC board trace inductance.

### 5.3 V<sub>CCW</sub> Trace on Printed Circuit Boards

Updating flash memories that reside in the target system requires that the printed circuit board designer pay attention to the V<sub>CCW</sub> Power supply trace. The V<sub>CCW</sub> pin supplies the memory cell current for word writing and block erasing. Use similar trace widths and layout considerations given to the V<sub>CC</sub> power bus. Adequate V<sub>CCW</sub> supply traces and decoupling will decrease V<sub>CCW</sub> voltage spikes and overshoots.

### 5.4 V<sub>CC</sub>, V<sub>CCW</sub>, RP# Transitions

Block erase, full chip erase, word write and lock-bit configuration are not guaranteed if V<sub>CCW</sub> falls outside of a valid V<sub>CCWH1/2</sub> range, V<sub>CC</sub> falls outside of a valid 2.7V-3.6V range, or RP# $\neq$ V<sub>IH</sub>. If V<sub>CCW</sub> error is detected, status register bit SR.3 is set to "1" along with SR.4 or SR.5, depending on the attempted operation. If RP# transitions to V<sub>IL</sub> during block erase, full chip erase, word write or lock-bit configuration, SR.7 will remain "0" until the reset operation is complete. Then, the operation will abort and the device will enter reset mode. The aborted operation may leave data partially altered. Therefore, the command sequence must be repeated after normal operation is restored. Device power-off or RP# transitions to V<sub>IL</sub> clear the status register.

The CUI latches commands issued by system software and is not altered by V<sub>CCW</sub> or CE# transitions or WSM actions. Its state is read array mode upon power-up, after exit from reset mode or after V<sub>CC</sub> transitions below V<sub>LKO</sub>.

## 5.5 Power-Up/Down Protection

The device is designed to offer protection against accidental block erase, full chip erase, word write or lock-bit configuration during power transitions. Upon power-up, the device is indifferent as to which power supply ( $V_{CCW}$  or  $V_{CC}$ ) powers-up first. Internal circuitry resets the CUI to read array mode at power-up.

A system designer must guard against spurious writes for  $V_{CC}$  voltages above  $V_{LKO}$  when  $V_{CCW}$  is active. Since both  $WE\#$  and  $CE\#$  must be low for a command write, driving either to  $V_{IH}$  will inhibit writes. The CUI's two-step command sequence architecture provides added level of protection against data alteration.

In-system block lock and unlock capability prevents inadvertent data alteration. The device is disabled while  $RP\#=V_{IL}$  regardless of its control inputs state.

## 5.6 Power Dissipation

When designing portable systems, designers must consider battery power consumption not only during device operation, but also for data retention during system idle time. Flash memory's nonvolatility increases usable battery life because data is retained when system power is removed.

## 5.7 Data Protection Method

Noises having a level exceeding the limit specified in the specification may be generated under specific operating conditions on some systems. Such noises, when induced onto  $WE\#$  signal or power supply, may be interpreted as false commands, causing undesired memory updating. To protect the data stored in the flash memory against unwanted overwriting, systems operating with the flash memory should have the following write protect designs, as appropriate:

### 1) Protecting data in specific block

When a lock bit is set, the corresponding block (includes the 2 boot blocks) is protected against overwriting. By setting a  $WP\#$  to low, only the 2 boot blocks can be protected against overwriting. By using this feature, the flash memory space can be divided into the program section (locked section) and data section (unlocked section). The permanent lock bit can be used to prevent false block bit setting. For further information on setting/resetting lock-bit, refer to the specification. (See chapter 4.10 and 4.11.)

### 2) Data protection through $V_{CCW}$

When the level of  $V_{CCW}$  is lower than  $V_{CCWLK}$  (lockout voltage), write operation on the flash memory is disabled. All blocks are locked and the data in the blocks are completely write protected. For the lockout voltage, refer to the specification. (See chapter 6.2.3.)

### 3) Data protection through $RP\#$

When the  $RP\#$  is kept low during read mode, the flash memory will be reset mode, then write protecting all blocks. When the  $RP\#$  is kept low during power up and power down sequence such as voltage transition, write operation on the flash memory is disabled, write protecting all blocks. For the details of  $RP\#$  control, refer to the specification. (See chapter 5.5 and 6.2.7.)

## 6 ELECTRICAL SPECIFICATIONS

### 6.1 Absolute Maximum Ratings\*

Operating Temperature	
During Read, Block Erase, Full Chip Erase, Word Write and Lock-Bit Configuration .....	-40°C to +85°C <sup>(1)</sup>
Storage Temperature	
During under Bias .....	-40°C to +85°C
During non Bias .....	-65°C to +125°C
Voltage On Any Pin (except V <sub>CC</sub> and V <sub>CCW</sub> ) .....	-0.5V to V <sub>CC</sub> +0.5V <sup>(2)</sup>
V <sub>CC</sub> Supply Voltage.....	-0.2V to +4.6V <sup>(2)</sup>
V <sub>CCW</sub> Supply Voltage.....	-0.2V to +13.0V <sup>(2,3)</sup>
Output Short Circuit Current.....	100mA <sup>(4)</sup>

*\*WARNING: Stressing the device beyond the "Absolute Maximum Ratings" may cause permanent damage. These are stress ratings only. Operation beyond the "Operating Conditions" is not recommended and extended exposure beyond the "Operating Conditions" may affect device reliability.*

#### NOTES:

1. Operating temperature is for extended temperature product defined by this specification.
2. All specified voltages are with respect to GND. Minimum DC voltage is -0.5V on input/output pins and -0.2V on V<sub>CC</sub> and V<sub>CCW</sub> pins. During transitions, this level may undershoot to -2.0V for periods <20ns. Maximum DC voltage on input/output pins are V<sub>CC</sub>+0.5V which, during transitions, may overshoot to V<sub>CC</sub>+2.0V for periods <20ns.
3. Maximum DC voltage on V<sub>CCW</sub> may overshoot to +13.0V for periods <20ns. Applying 12V±0.3V to V<sub>CCW</sub> during erase/write can only be done for a maximum of 1000 cycles on each block. V<sub>CCW</sub> may be connected to 12V±0.3V for a total of 80 hours maximum.
4. Output shorted for no more than one second. No more than one output shorted at a time.

### 6.2 Operating Conditions

Temperature and V<sub>CC</sub> Operating Conditions

Symbol	Parameter	Min.	Max.	Unit	Test Condition
T <sub>A</sub>	Operating Temperature	-40	+85	°C	Ambient Temperature
V <sub>CC</sub>	V <sub>CC</sub> Supply Voltage (2.7V-3.6V)	2.7	3.6	V	

#### 6.2.1 CAPACITANCE<sup>(1)</sup>

T<sub>A</sub>=+25°C, f=1MHz

Symbol	Parameter	Typ.	Max.	Unit	Condition
C <sub>IN</sub>	Input Capacitance	7	10	pF	V <sub>IN</sub> =0.0V
C <sub>OUT</sub>	Output Capacitance	9	12	pF	V <sub>OUT</sub> =0.0V

#### NOTE:

1. Sampled, not 100% tested.

## 6.2.2 AC INPUT/OUTPUT TEST CONDITIONS

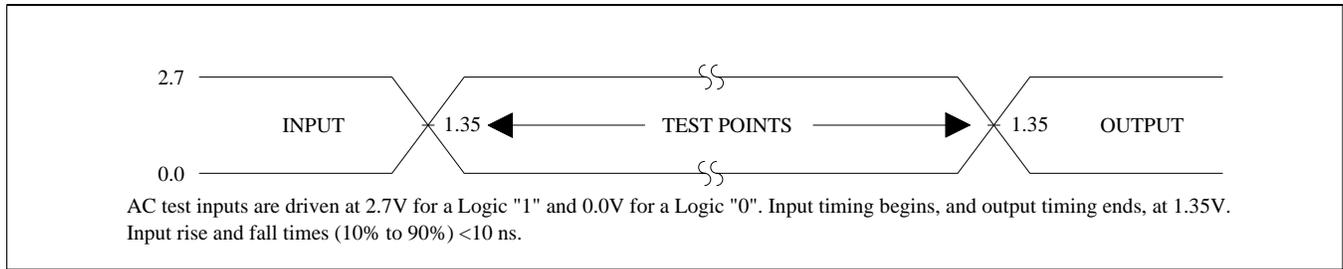
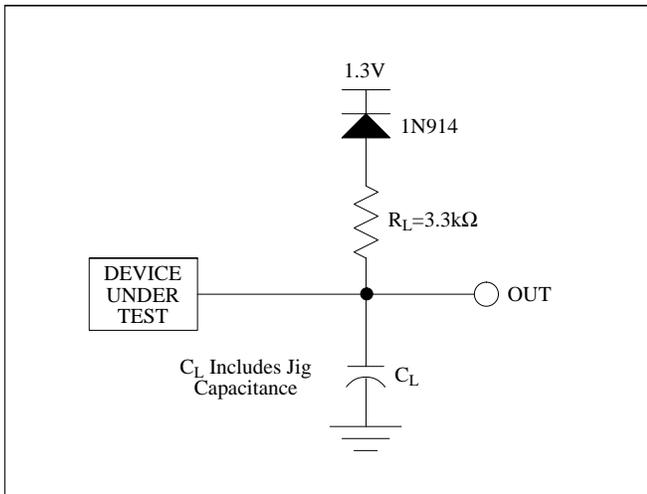
Figure 14. Transient Input/Output Reference Waveform for  $V_{CC}=2.7V-3.6V$ 

Figure 15. Transient Equivalent Testing Load Circuit

Test Configuration Capacitance Loading Value

Test Configuration	$C_L$ (pF)
$V_{CC}=2.7V-3.6V$	50

## 6.2.3 DC CHARACTERISTICS

DC Characteristics

Sym.	Parameter	Notes	V <sub>CC</sub> =2.7V-3.6V		Unit	Test Conditions
			Typ.	Max.		
I <sub>LI</sub>	Input Load Current	1		±0.5	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. V <sub>IN</sub> =V <sub>CC</sub> or GND
I <sub>LO</sub>	Output Leakage Current	1		±0.5	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. V <sub>OUT</sub> =V <sub>CC</sub> or GND
I <sub>CCS</sub>	V <sub>CC</sub> Standby Current	1	4	20	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. CE#=RP#=V <sub>CC</sub> ±0.2V
I <sub>CCAS</sub>	V <sub>CC</sub> Auto Power-Save Current	1,4	4	20	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. CE#=GND±0.2V
I <sub>CCD</sub>	V <sub>CC</sub> Reset Power-Down Current	1	4	20	μA	RP#=GND±0.2V
I <sub>CCR</sub>	V <sub>CC</sub> Read Current	1	15	30	mA	V <sub>CC</sub> =V <sub>CC</sub> Max., CE#=GND f=5MHz, I <sub>OUT</sub> =0mA
I <sub>CCW</sub>	V <sub>CC</sub> Word Write or Set Lock-Bit Current	1,5	5	17	mA	V <sub>CCW</sub> =2.7V-3.6V
			5	12	mA	V <sub>CCW</sub> =11.7V-12.3V
I <sub>CCE</sub>	V <sub>CC</sub> Block Erase, Full Chip Erase or Clear Block Lock-Bits Current	1,5	4	17	mA	V <sub>CCW</sub> =2.7V-3.6V
			4	12	mA	V <sub>CCW</sub> =11.7V-12.3V
I <sub>CCWS</sub> I <sub>CCES</sub>	V <sub>CC</sub> Word Write or Block Erase Suspend Current	1,2	1	6	mA	CE#=V <sub>IH</sub>
I <sub>CCWS</sub> I <sub>CCWR</sub>	V <sub>CCW</sub> Standby or Read Current	1	±2	±15	μA	V <sub>CCW</sub> ≤V <sub>CC</sub>
			10	200	μA	V <sub>CCW</sub> >V <sub>CC</sub>
I <sub>CCWAS</sub>	V <sub>CCW</sub> Auto Power-Save Current	1,4	0.1	5	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. CE#=GND±0.2V
I <sub>CCWD</sub>	V <sub>CCW</sub> Reset Power-Down Current	1	0.1	5	μA	RP#=GND±0.2V
I <sub>CCWW</sub>	V <sub>CCW</sub> Word Write or Set Lock-Bit Current	1,5	12	40	mA	V <sub>CCW</sub> =2.7V-3.6V
				30	mA	V <sub>CCW</sub> =11.7V-12.3V
I <sub>CCWE</sub>	V <sub>CCW</sub> Block Erase, Full Chip Erase or Clear Block Lock-Bits Current	1,5	8	25	mA	V <sub>CCW</sub> =2.7V-3.6V
				20	mA	V <sub>CCW</sub> =11.7V-12.3V
I <sub>CCWWS</sub> I <sub>CCWES</sub>	V <sub>CCW</sub> Word Write or Block Erase Suspend Current	1	10	200	μA	V <sub>CCW</sub> =V <sub>CCWH1/2</sub>

## DC Characteristics (Continued)

Sym.	Parameter	Notes	$V_{CC}=2.7V-3.6V$		Unit	Test Conditions
			Min.	Max.		
$V_{IL}$	Input Low Voltage	5	-0.5	0.4	V	
$V_{IH}$	Input High Voltage	5	$V_{CC}-0.4$	$V_{CC}+0.5$	V	
$V_{OL}$	Output Low Voltage	5		0.4	V	$V_{CC}=V_{CC} \text{ Min.}$ $I_{OL}=2.0\text{mA}$
$V_{OH}$	Output High Voltage	5	$V_{CC}-0.4$		V	$V_{CC}=V_{CC} \text{ Min.}$ $I_{OH}=-100\mu\text{A}$
$V_{CCWLK}$	$V_{CCW}$ Lockout during Normal Operations	3,5		1.0	V	
$V_{CCWH1}$	$V_{CCW}$ during Block Erase, Full Chip Erase, Word Write or Lock-Bit Configuration Operations		2.7	3.6	V	
$V_{CCWH2}$	$V_{CCW}$ during Block Erase, Full Chip Erase, Word Write or Lock-Bit Configuration Operations	6	11.7	12.3	V	
$V_{LKO}$	$V_{CC}$ Lockout Voltage		2.0		V	

## NOTES:

- All currents are in RMS unless otherwise noted. Typical values at nominal  $V_{CC}$  voltage and  $T_A=+25^\circ\text{C}$ .
- $I_{CCWS}$  and  $I_{CCES}$  are specified with the device de-selected. If read or word written while in erase suspend mode, the device's current draw is the sum of  $I_{CCWS}$  or  $I_{CCES}$  and  $I_{CCR}$  or  $I_{CCW}$ , respectively.
- Block erases, full chip erase, word writes and lock-bit configurations are inhibited when  $V_{CCW} \leq V_{CCWLK}$ , and not guaranteed in the range between  $V_{CCWLK}(\text{max.})$  and  $V_{CCWH1}(\text{min.})$ , between  $V_{CCWH1}(\text{max.})$  and  $V_{CCWH2}(\text{min.})$  and above  $V_{CCWH2}(\text{max.})$ .
- The Automatic Power Savings (APS) feature is placed automatically power save mode that addresses not switching more than 300ns while read mode.
- Sampled, not 100% tested.
- Applying  $12V \pm 0.3V$  to  $V_{CCW}$  during erase/write can only be done for a maximum of 1000 cycles on each block.  $V_{CCW}$  may be connected to  $12V \pm 0.3V$  for a total of 80 hours maximum.

6.2.4 AC CHARACTERISTICS - READ-ONLY OPERATIONS<sup>(1)</sup>
 $V_{CC}=2.7V-3.6V, T_A=-40^{\circ}C \text{ to } +85^{\circ}C$ 

Sym.	Parameter	Notes	Min.	Max.	Unit
$t_{AVAV}$	Read Cycle Time		90		ns
$t_{AVQV}$	Address to Output Delay			90	ns
$t_{ELQV}$	CE# to Output Delay	2		90	ns
$t_{PHQV}$	RP# High to Output Delay			600	ns
$t_{GLQV}$	OE# to Output Delay	2		40	ns
$t_{ELOX}$	CE# to Output in Low Z	3	0		ns
$t_{EHQZ}$	CE# High to Output in High Z	3		40	ns
$t_{GLQX}$	OE# to Output in Low Z	3	0		ns
$t_{GHQZ}$	OE# High to Output in High Z	3		15	ns
$t_{OH}$	Output Hold from Address, CE# or OE# Change, Whichever Occurs First	3	0		ns

## NOTES:

1. See AC Input/Output Reference Waveform for maximum allowable input slew rate.
2. OE# may be delayed up to  $t_{ELQV}-t_{GLQV}$  after the falling edge of CE# without impact on  $t_{ELQV}$ .
3. Sampled, not 100% tested.

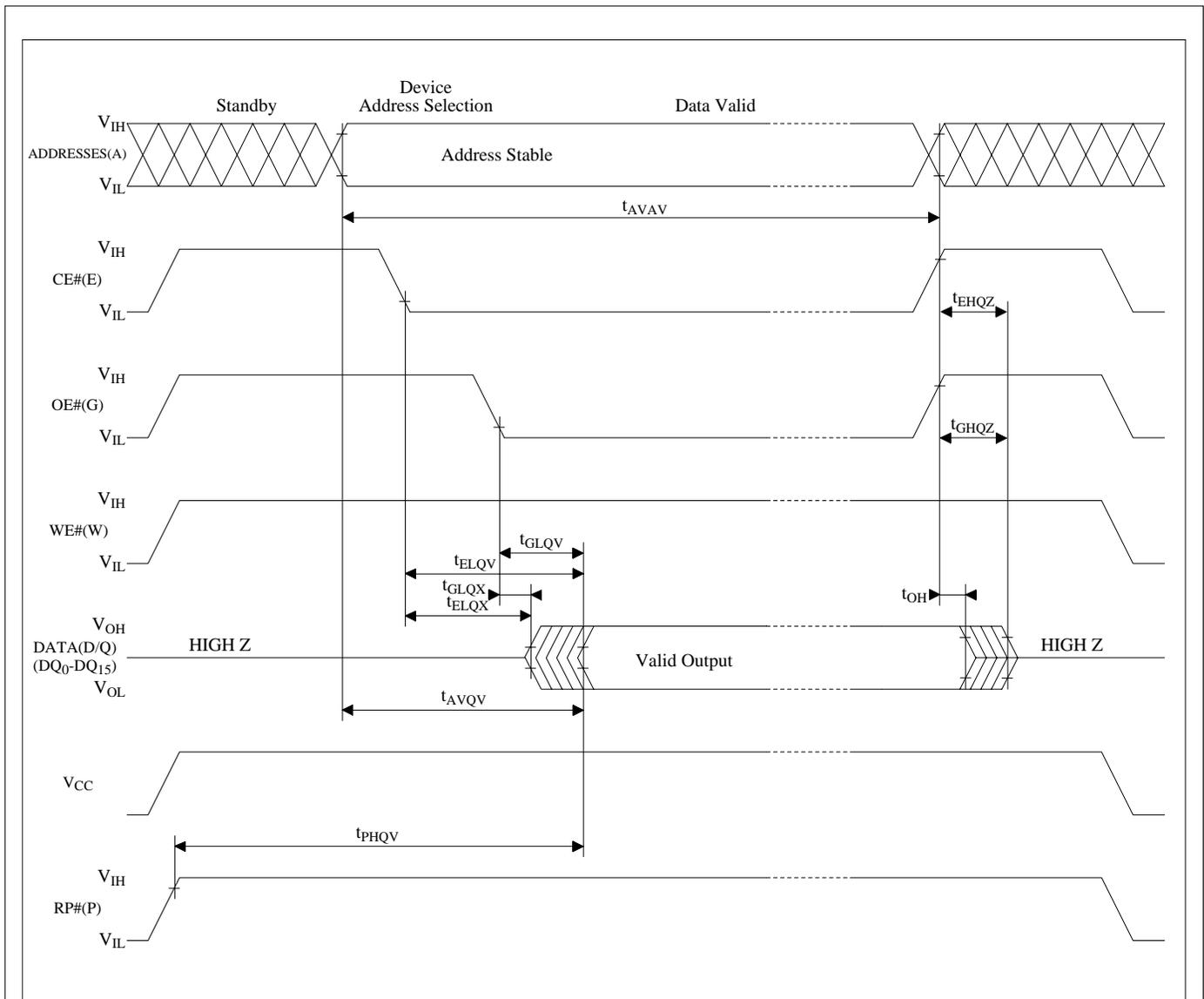


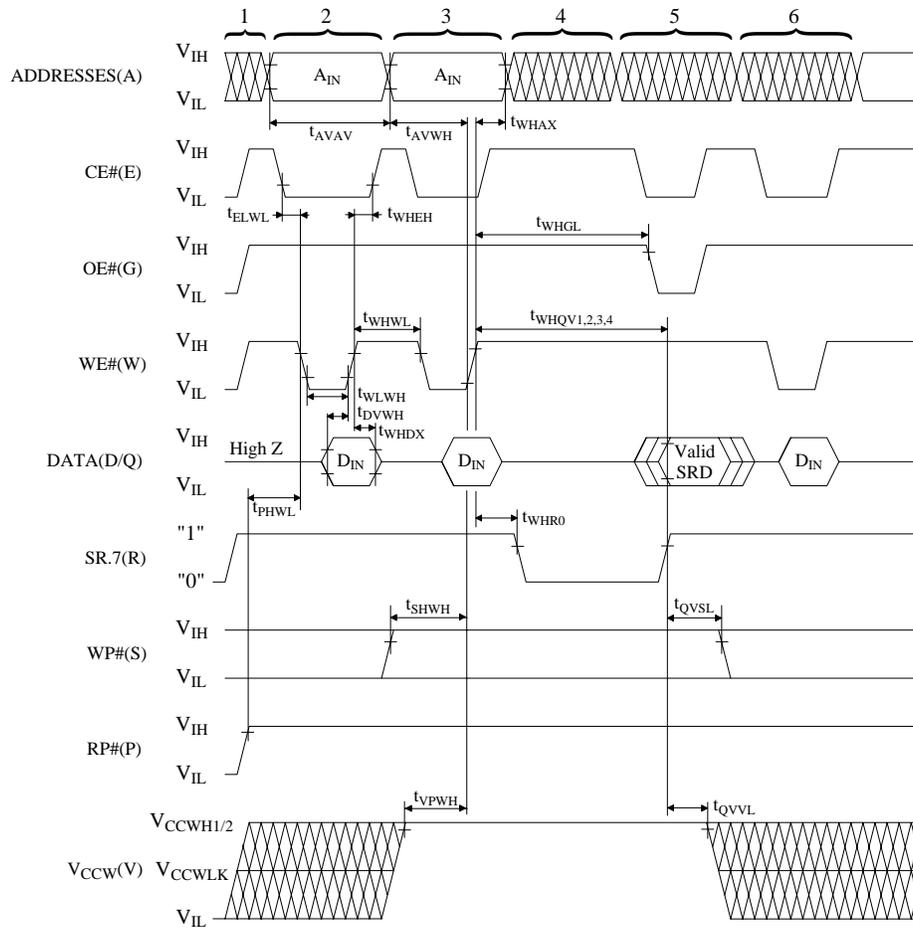
Figure 16. AC Waveform for Read Operations

6.2.5 AC CHARACTERISTICS - WRITE OPERATIONS<sup>(1)</sup>
 $V_{CC}=2.7V-3.6V, T_A=-40^{\circ}C \text{ to } +85^{\circ}C$ 

Sym.	Parameter	Notes	Min.	Max.	Unit
$t_{AVAV}$	Write Cycle Time		90		ns
$t_{PHWL}$	RP# High Recovery to WE# Going Low	2	1		$\mu s$
$t_{ELWL}$	CE# Setup to WE# Going Low		10		ns
$t_{WLWH}$	WE# Pulse Width		50		ns
$t_{SHWH}$	WP# $V_{IH}$ Setup to WE# Going High	2	100		ns
$t_{VPWH}$	$V_{CCW}$ Setup to WE# Going High	2	100		ns
$t_{AVWH}$	Address Setup to WE# Going High	3	50		ns
$t_{DVWH}$	Data Setup to WE# Going High	3	50		ns
$t_{WHDX}$	Data Hold from WE# High		0		ns
$t_{WHAX}$	Address Hold from WE# High		0		ns
$t_{WHEH}$	CE# Hold from WE# High		10		ns
$t_{WHWL}$	WE# Pulse Width High		30		ns
$t_{WHR0}$	WE# High to SR.7 Going "0"			100	ns
$t_{WHGL}$	Write Recovery before Read		0		ns
$t_{QVVL}$	$V_{CCW}$ Hold from Valid SRD	2,4	0		ns
$t_{QVSL}$	WP# $V_{IH}$ Hold from Valid SRD	2,4	0		ns

## NOTES:

1. Read timing characteristics during block erase, full chip erase, word write and lock-bit configuration operations are the same as during read-only operations. Refer to AC Characteristics for read-only operations.
2. Sampled, not 100% tested.
3. Refer to Table 3 for valid  $A_{IN}$  and  $D_{IN}$  for block erase, full chip erase, word write or lock-bit configuration.
4.  $V_{CCW}$  should be held at  $V_{CCWH1/2}$  until determination of block erase, full chip erase, word write or lock-bit configuration success (SR.1/3/4/5=0).



## NOTES:

1.  $V_{CC}$  power-up and standby.
2. Write each setup command.
3. Write each confirm command or valid address and data.
4. Automated erase or program delay.
5. Read status register data.
6. Write Read Array command.

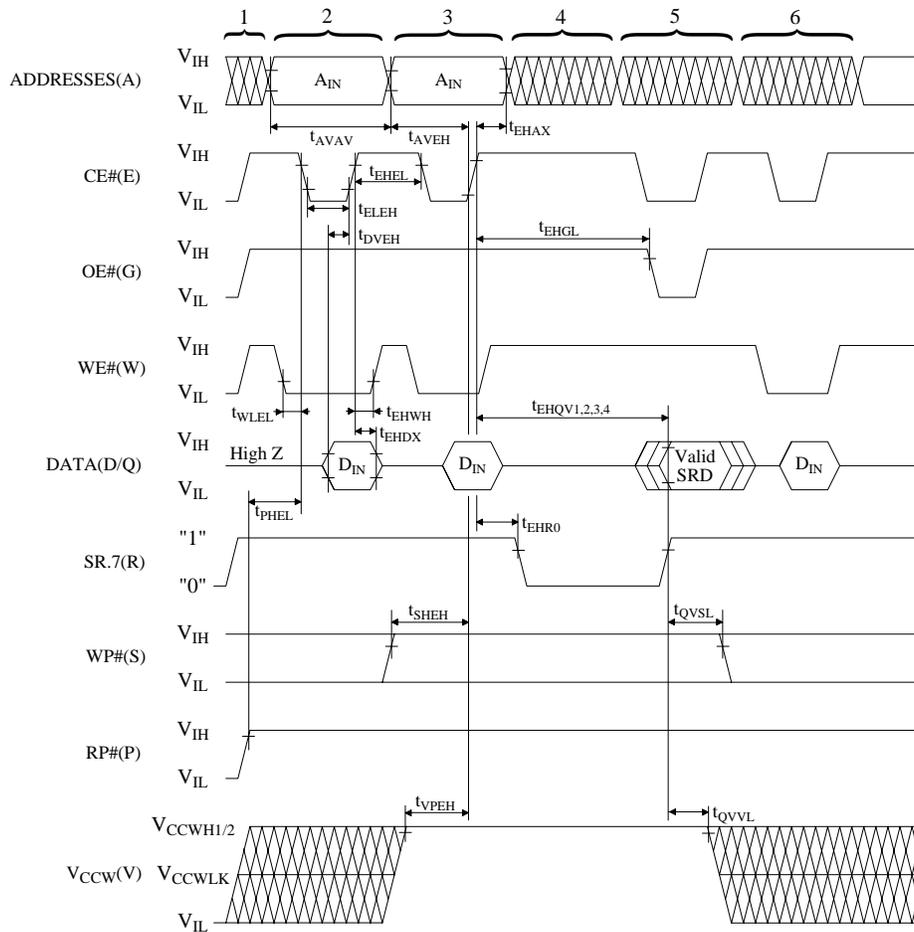
Figure 17. AC Waveform for WE#-Controlled Write Operations

6.2.6 ALTERNATIVE CE#-CONTROLLED WRITES<sup>(1)</sup>
 $V_{CC}=2.7V-3.6V, T_A=-40^{\circ}C \text{ to } +85^{\circ}C$ 

Sym.	Parameter	Notes	Min.	Max.	Unit
$t_{AVAV}$	Write Cycle Time		90		ns
$t_{PHEL}$	RP# High Recovery to CE# Going Low	2	1		$\mu s$
$t_{WLEL}$	WE# Setup to CE# Going Low		0		ns
$t_{ELEH}$	CE# Pulse Width		65		ns
$t_{SHEH}$	WP# $V_{IH}$ Setup to CE# Going High	2	100		ns
$t_{VPEH}$	$V_{CCW}$ Setup to CE# Going High	2	100		ns
$t_{AVEH}$	Address Setup to CE# Going High	3	50		ns
$t_{DVEH}$	Data Setup to CE# Going High	3	50		ns
$t_{EHDX}$	Data Hold from CE# High		0		ns
$t_{EHAX}$	Address Hold from CE# High		0		ns
$t_{EHWH}$	WE# Hold from CE# High		0		ns
$t_{EHEL}$	CE# Pulse Width High		25		ns
$t_{EHR0}$	CE# High to SR.7 Going "0"			100	ns
$t_{EHGL}$	Write Recovery before Read		0		ns
$t_{QVVL}$	$V_{CCW}$ Hold from Valid SRD	2,4	0		ns
$t_{QVSL}$	WP# $V_{IH}$ Hold from Valid SRD	2,4	0		ns

## NOTES:

1. In systems where CE# defines the write pulse width (within a longer WE# timing waveform), all setup, hold, and inactive WE# times should be measured relative to the CE# waveform.
2. Sampled, not 100% tested.
3. Refer to Table 3 for valid  $A_{IN}$  and  $D_{IN}$  for block erase, full chip erase, word write or lock-bit configuration.
4.  $V_{CCW}$  should be held at  $V_{CCWH1/2}$  until determination of block erase, full chip erase, word write or lock-bit configuration success (SR.1/3/4/5=0).



## NOTES:

1. V<sub>CC</sub> power-up and standby.
2. Write each setup command.
3. Write each confirm command or valid address and data.
4. Automated erase or program delay.
5. Read status register data.
6. Write Read Array command.

Figure 18. AC Waveform for CE#-Controlled Write Operations

## 6.2.7 RESET OPERATIONS

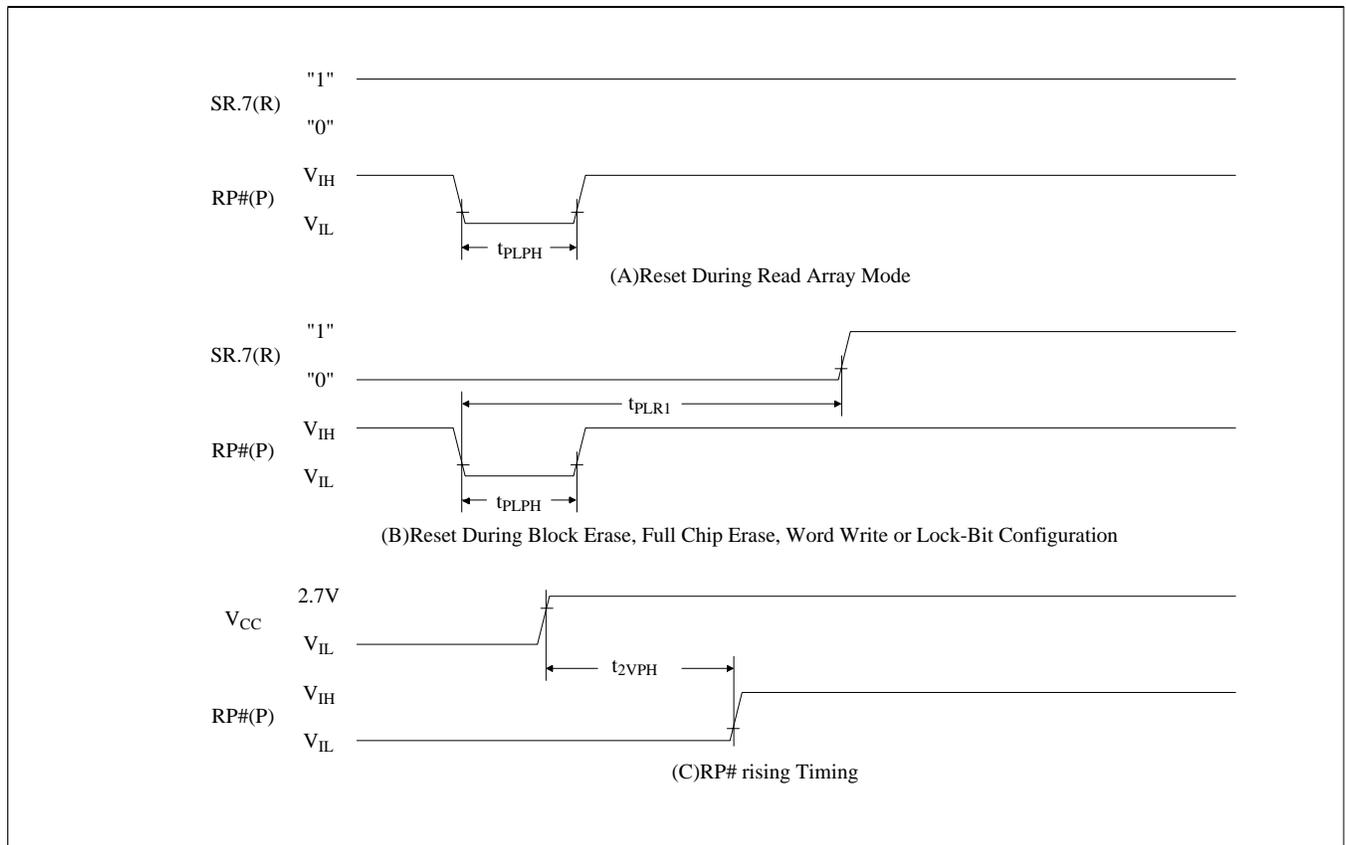


Figure 19. AC Waveform for Reset Operation

## Reset AC Specifications

Sym.	Parameter	Notes	Min.	Max.	Unit
$t_{PLPH}$	RP# Pulse Low Time	2	100		ns
$t_{PLR1}$	RP# Low to Reset during Block Erase, Full Chip Erase, Word Write or Lock-Bit Configuration	1,2		30	$\mu$ s
$t_{2VPH}$	$V_{CC}$ 2.7V to RP# High	2,3	100		ns

## NOTES:

1. If RP# is asserted while a block erase, full chip erase, word write or lock-bit configuration operation is not executing, the reset will complete within 100ns.
2. A reset time,  $t_{PHQV}$ , is required from the later of SR.7 going "1" or RP# going high until outputs are valid. Refer to AC Characteristics - Read-Only Operations for  $t_{PHQV}$ .
3. When the device power-up, holding RP# low minimum 100ns is required after  $V_{CC}$  has been in predefined range and also has been in stable there.

### 6.2.8 BLOCK ERASE, FULL CHIP ERASE, WORD WRITE AND LOCK-BIT CONFIGURATION PERFORMANCE<sup>(3)</sup>

$V_{CC}=2.7V-3.6V$ ,  $T_A=-40^{\circ}C$  to  $+85^{\circ}C$

Sym.	Parameter		Notes	$V_{CCW}=2.7V-3.6V$			$V_{CCW}=11.7V-12.3V$			Unit
				Min.	Typ. <sup>(1)</sup>	Max.	Min.	Typ. <sup>(1)</sup>	Max.	
$t_{WHQV1}$	Word Write Time	32K word Block	2		33	200		20		$\mu s$
$t_{EHQV1}$		4K word Block	2		36	200		27		$\mu s$
	Block Write Time	32K word Block	2		1.1	4		0.66		s
		4K word Block	2		0.15	0.5		0.12		s
$t_{WHQV2}$	Block Erase Time	32K word Block	2		1.2	6		0.9		s
$t_{EHQV2}$		4K word Block	2		0.6	5		0.5		s
	Full Chip Erase Time		2		84	420		64		s
$t_{WHQV3}$ $t_{EHQV3}$	Set Lock-Bit Time		2		56	200		42		$\mu s$
$t_{WHQV4}$ $t_{EHQV4}$	Clear Block Lock-Bits Time		2		1	5		0.69		s
$t_{WHR11}$ $t_{EHR11}$	Word Write Suspend Latency Time to Read		4		6	15		6	15	$\mu s$
$t_{WHR12}$ $t_{EHR12}$	Block Erase Suspend Latency Time to Read		4		16	30		16	30	$\mu s$
$t_{ERES}$	Latency Time from Block Erase Resume Command to Block Erase Suspend Command		5	600			600			$\mu s$

#### NOTES:

1. Typical values measured at  $T_A=+25^{\circ}C$  and  $V_{CC}=3.0V$ ,  $V_{CCW}=3.0V$  or  $12.0V$ . Assumes corresponding lock-bits are not set. Subject to change based on device characterization.
2. Excludes system-level overhead.
3. Sampled but not 100% tested.
4. A latency time is required from issuing suspend command(WE# or CE# going high) until SR.7 going "1".
5. If the time between writing the Block Erase Resume command and writing the Block Erase Suspend command is shorter than  $t_{ERES}$  and both commands are written repeatedly, a longer time is required than standard block erase until the completion of the operation.

## A-1 RECOMMENDED OPERATING CONDITIONS

### A-1.1 At Device Power-Up

AC timing illustrated in Figure A-1 is recommended for the supply voltages and the control signals at device power-up. If the timing in the figure is ignored, the device may not operate correctly.

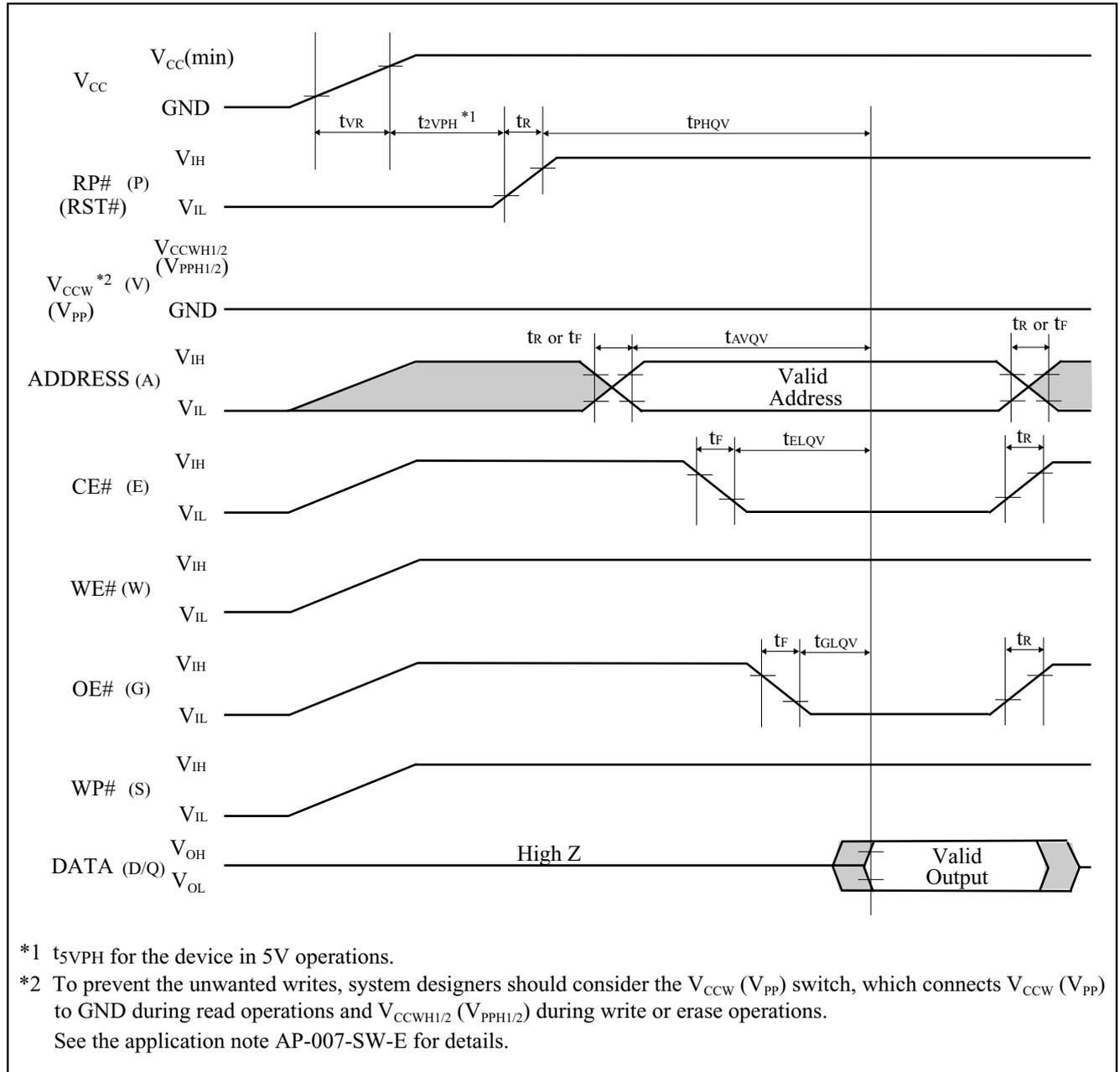


Figure A-1. AC Timing at Device Power-Up

For the AC specifications  $t_{VR}$ ,  $t_R$ ,  $t_F$  in the figure, refer to the next page. See the “ELECTRICAL SPECIFICATIONS” described in specifications for the supply voltage range, the operating temperature and the AC specifications not shown in the next page.

## A-1.1.1 Rise and Fall Time

Symbol	Parameter	Notes	Min.	Max.	Unit
$t_{VR}$	$V_{CC}$ Rise Time	1	0.5	30000	$\mu\text{s}/\text{V}$
$t_R$	Input Signal Rise Time	1, 2		1	$\mu\text{s}/\text{V}$
$t_F$	Input Signal Fall Time	1, 2		1	$\mu\text{s}/\text{V}$

## NOTES:

1. Sampled, not 100% tested.
2. This specification is applied for not only the device power-up but also the normal operations.  
 $t_R(\text{Max.})$  and  $t_F(\text{Max.})$  for RP# (RST#) are  $50\mu\text{s}/\text{V}$ .

## A-1.2 Glitch Noises

Do not input the glitch noises which are below  $V_{IH}$  (Min.) or above  $V_{IL}$  (Max.) on address, data, reset, and control signals, as shown in Figure A-2 (b). The acceptable glitch noises are illustrated in Figure A-2 (a).

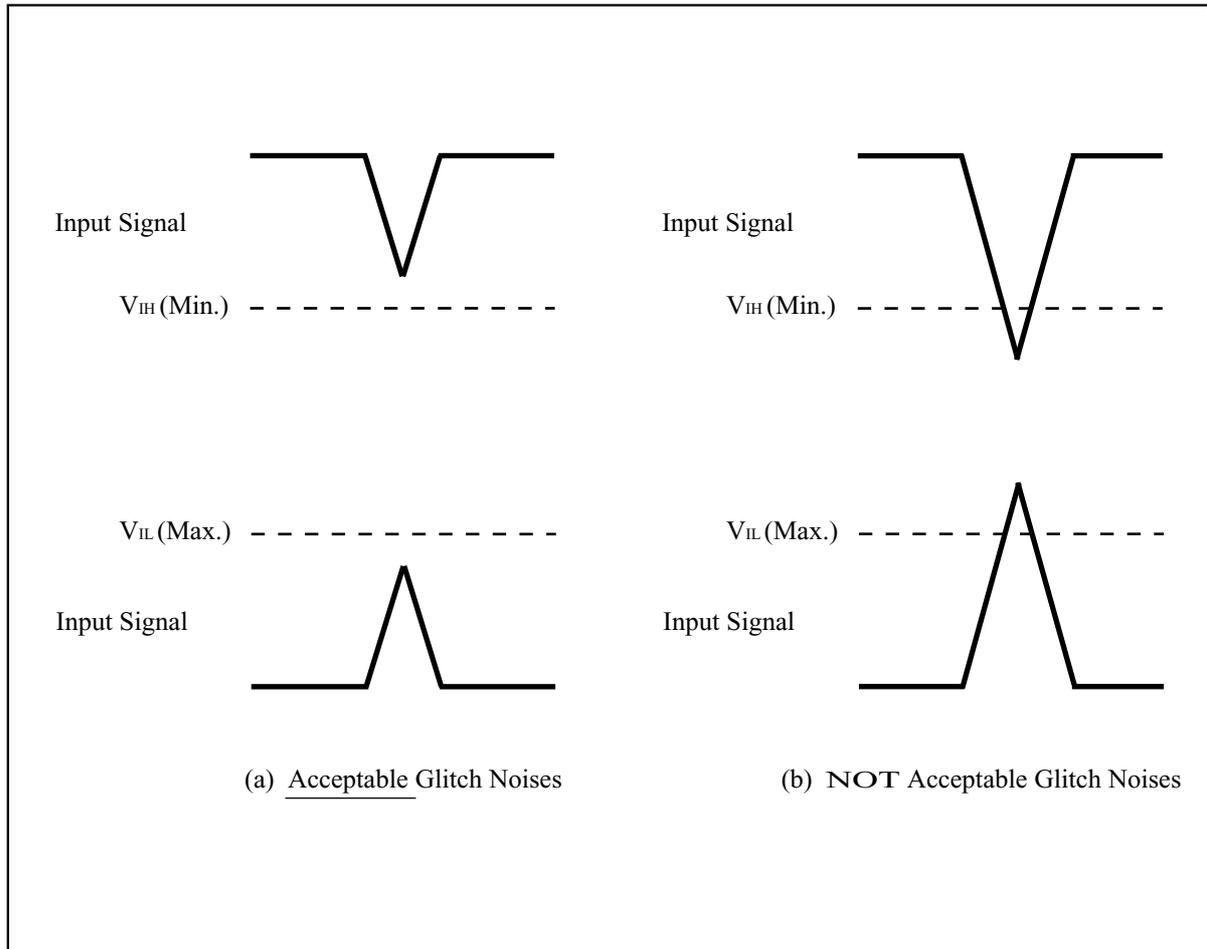


Figure A-2. Waveform for Glitch Noises

See the "DC CHARACTERISTICS" described in specifications for  $V_{IH}$  (Min.) and  $V_{IL}$  (Max.).

A-2 RELATED DOCUMENT INFORMATION<sup>(1)</sup>

Document No.	Document Name
AP-001-SD-E	Flash Memory Family Software Drivers
AP-006-PT-E	Data Protection Method of SHARP Flash Memory
AP-007-SW-E	RP#, $V_{PP}$ Electric Potential Switching Circuit

## NOTE:

1. International customers should contact their local SHARP or distribution sales office.