## OneNAND™ Specification

Density	Part No.	Vcc(core & IO)	Temperature	PKG
512Mb	KFG1216Q2A	1.8V(1.7V~1.95V)	Extended	63FBGA(LF)/63FBGA
	KFG1216D2A	2.65V(2.4V~2.9V)	Extended	63FBGA(LF)/63FBGA
	KFG1216U2A	3.3V(2.7V~3.6V)	Industrial	63FBGA(LF)/63FBGA

Version: Ver. 1.0 Date: May 17th, 2005



### 1.0 INTRODUCTION

This specification contains information about the Samsung Electronics Company OneNAND<sup>TM</sup>, Flash memory product family. Section 1.0 includes a general overview, revision history, and product ordering information.

Section 2.0 describes the OneNAND device. Section 3.0 provides information about device operation. Electrical specifications and timing waveforms are in Sections 4.0 though 6.0. Section 7.0 provides additional application and technical notes pertaining to use of the OneNAND. Package dimensions are found in Section 8.0

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# 1.1 Revision History

## **Document Title**

### **OneNAND**

## **Revision History**

Revision No.	<u>History</u>	<u>Draft Date</u>	<u>Remark</u>
0.0	Initial issue.	Nov. 4, 2004	Preliminary
0.1	1. Corrected Errata 2. Revised cache read flow chart 3. Revised standby current 4. Revised spare area description 5. Added CE don't care state for Asynch Write, Load, Program, and Block Erase timing diagram	Dec. 7, 2004	Preliminary
0.2	Added Copy Back Operation with Random Data Input     Changed tBA from 11ns to 11.5ns     Pended Active Erase Current	Dec. 24, 2004	Preliminary
0.3	<ol> <li>Corrected the errata</li> <li>Revised typical value of ISB from 50uA to 10uA</li> <li>Revised maximum value of ISB from 100uA to 50uA</li> <li>Revised maximum value of tCE, tAA and tACC from 70ns to 76ns</li> <li>Revised Vcc-IO description</li> <li>Revised Spare Area description</li> <li>Added extra information on Controller Status Register</li> <li>Added commands related to Interrupt Status Register bits</li> <li>Revised Write Protection Status on Chapter 3.4.3</li> <li>Revised Copy-Back Program Operation description</li> <li>Added extra information on Multi-Block Erase Operation</li> <li>Disabled FBA restriction in OTP operation</li> <li>Revised Cache Read Flow Chart</li> <li>Revised Reset Parameter descriptions</li> <li>Added RDY information on Warm Reset Timing diagram</li> <li>Added information on Data Protection Timing During Power Down</li> <li>Revised Interrupt pin rise and falling slope graph</li> <li>Added restriction on address register setting on Dual Operations</li> <li>Added restriction on address register setting on Cache Read Operation</li> </ol>	Jan. 13, 2005	Preliminary
0.4	1. Corrected the errata 2. Updated DC parameters to RMS values 3. Added Speed Information on Product Number 4. Revised tOEZ description 5. Revised OTP register setting restriction 6. Added Boot Sequence Infrormation on Technical Notes 7. Added Cint Information	Feb. 25, 2005	Preliminary



Final

May. 17, 2005

### **Revision History**

Revision No. History Draft Date Remark

1.0 1. Corrected the errata

2. Added Data Protection flow chart.

3. Removed Cache Read Operation.

4. Added additional information on command register.

5. Revised Interrupt status register information.

6. Added INT pin schematic.

7. Changed tPGM1 to 205 from 320us, tPGM2 to 220 from 350us.

8. Revised AC/DC parameters

9. Revised ECC Bypass Description.

10. Revised Reset Parameters and Timing Diagrams.

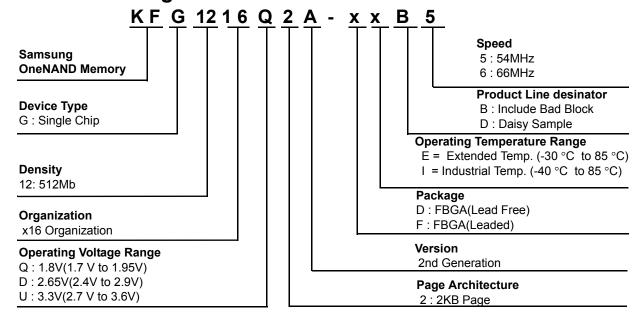
## 1.2 Flash Product Type Selector

Samsung offers a variety of Flash solutions including NAND Flash, OneNAND<sup>TM</sup> and NOR Flash. Samsung offers Flash products both component and a variety of card formats including RS-MMC, MMC, CompactFlash, and SmartMedia.

To determine which Samsung Flash product solution is best for your application, refer the product selector chart.

Application Requires	Samsung Flash Products				
Application Requires	NAND	OneNAND	NOR		
Fast Random Read			•		
Fast Sequential Read	•	•			
Fast Write/Program	•	•			
Multi Block Erase		• (Max 64 Blocks)	•		
Erase Suspend/Resume		•	•		
Copyback	• (EDC)	• (ECC)			
Lock/Unlock/Lock-Tight		•	•		
ECC	External (Hardware/Software)	Internal	X		
Scalability	•	•			

## 1.3 Ordering Information





### 1.4 Architectural Benefits

OneNAND is a highly integrated non-volatile memory solution based around a NAND Flash memory array.

The chip integrates system features including:

- · A BootRAM and bootloader
- Two independent bi-directional 2KB DataRAM buffers
- A High-Speed x16 Host Interface
- On-chip Error Correction
- On-chip NOR interface controller

This on-chip integration enables system designers to reduce external system logic and use high-density NAND Flash in applications that would otherwise have to use more NOR components.

OneNAND takes advantage of the higher performance NAND program time, low power, and high density and combines it with the synchronous read performance of NOR. The NOR Flash host interface makes OneNAND an ideal solution for applications like G3 Smart Phones, Camera Phones, and mobile applications that have large, advanced multimedia applications and operating systems, but lack a NAND controller.

When integrated into a Samsung Multi-Chip-Package with Samsung Mobile DDR SDRAM, designers can complete a high-performance, small footprint solution.

The attached datasheets are prepared and approved by SAMSUNG Electronics. SAMSUNG Electronics CO., LTD. reserve the right to change the specifications. SAMSUNG Electronics will evaluate and reply to your requests and questions about device. If you have any questions, please contact the SAMSUNG branch office near you.



#### 1.5 **Product Features**

#### **Device Architecture**

· Design Technology: 90nm

· Supply Voltage: 1.8V (1.7V ~ 1.95V), 2.65V (2.4 ~ 2.9V), 3.3V (2.7 ~ 3.6V)

· Host Interface: 16 bit

• 5KB Internal BufferRAM: 1KB BootRAM, 4KB DataRAM

 SLC NAND Array: (2K+64)B Page Size, (128K+4K)B Block Size

#### **Device Performance**

· Host Interface Type: Synchronous Burst Read

- Up to 54MHz clock frequency

- Linear Burst 4-, 8-, 16, 32-words with wrap around

- Continuous 1K word Sequential Burst

Asynchronous Random Read

- 76ns access time

Asynchronous Random Write Latency 3(up to 40MHz), 4, 5, 6, and 7

• Programmable Burst Read Latency Multiple Sector Read: Up to 4 sectors using Sector Count Register

Cold/Warm/Hot/NAND Flash Resets

up to 64 blocks

· Multi Block Erase: · Low Power Dissipation: Typical Power,

> - Standby current : 10uA@1.8Vdevice,

35uA@2.65V/3.3V device

- Synchronous Burst Read current: 12mA@1.8Vdevice

22mA@2.65V/3.3V device

30mA@all device - Load current : - Program current : 25mA@1.8v device 28mA@2.65V/3.3V device

- Erase current : 20mA@1.8V device

23mA@2.65V/3.3V device 20mA@1.8V device - Multi Block Erase current :

23mA@2.65V/3.3V device

#### **System Hardware**

• Voltage detector generating internal reset signal from Vcc

Hardware reset input (/RP)

• Multiple Reset Modes:

· Data Protection Modes - Write Protection for BootRAM

- Write Protection for NAND Flash Array - Write Protection during power-up - Write Protection during power-down

• User-controlled One Time Programmable(OTP) area

Internal 2bit EDC / 1bit ECC

· Internal Bootloader supports Booting Solution in system

· Handshaking Feature - INT pin indicates Ready / Busy

- Polling the interrupt register status bit

· Detailed chip information - by ID register

63ball, 9.5mm x 12mm x max 1.0mmt, 0.8mm ball pitch FBGA **Packaging** 



### 1.6 General Overview

OneNAND<sup>TM</sup>, is a monolithic integrated circuit with a NAND Flash array using a NOR Flash interface. This device includes control logic, a NAND Flash array, and 5KB of internal BufferRAM. The BufferRAM reserves 1KB for boot code buffering (BootRAM) and 4KB for data buffering (DataRAM), split between 2 independent buffers. It has a x16 Host Interface and a random access time speed of ~76ns.

The device operates up to a maximum host-driven clock frequency of 54MHz for synchronous reads at Vcc(or Vccq. Refer to chapter 4.2) with minimum 4-clock latency. Below 40MHz it is accessible with minimum 3-clock latency. Appropriate wait cycles are determined by programmable read latency.

OneNAND provides for multiple sector read operations by assigning the number of sectors to be read in the sector counter register. The device includes one block-sized OTP (One Time Programmable) area that can be used to increase system security or to provide identification capabilities.



### 2.0 DEVICE DESCRIPTION

### 2.1 Detailed Product Description

The OneNAND is an advanced generation, high-performance NAND-based Flash memory.

It integrates on-chip a single-level-cell (SLC) NAND Flash Array memory with two independent data buffers, boot RAM buffer, a page buffer for the Flash array, and a one-time-programmable block.

The combination of these memory areas enable high-speed pipelining of reads from host, BufferRAM, Page Buffer, and NAND Flash Array memory.

Clock speeds up to 54MHz with a x16 wide I/O yields a 68MByte/second bandwidth.

The OneNAND also includes a Boot RAM and boot loader. This enables the device to efficiently load boot code at device startup from the NAND Array without the need for off-chip boot device.

One block of the NAND Array is set aside as an OTP memory area. This area, available to the user, can be configured and locked with secured user information.

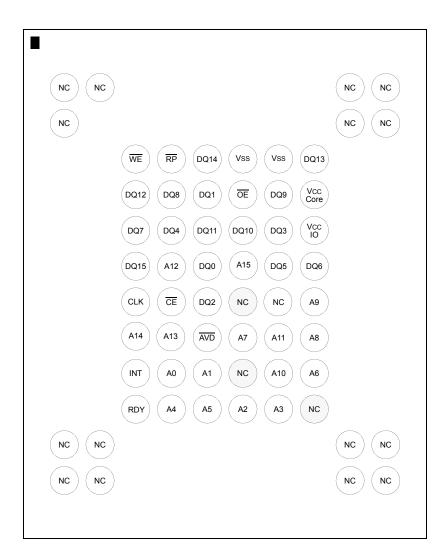
On-chip controller interfaces enable the device to operate in systems without NAND Host controllers.



## 2.2 **Definitions**

B (capital letter)	Byte, 8bits
W (capital letter)	Word, 16bits
b (lower-case letter)	Bit
ECC	Error Correction Code
Calculated ECC	ECC that has been calculated during a load or program access
Written ECC	ECC that has been stored as data in the NAND Flash array or in the BufferRAM
BufferRAM	On-chip internal buffer consisting of BootRAM and DataRAM
BootRAM	A 1KB portion of the BufferRAM reserved for Boot Code buffering
DataRAM	A 4KB portion of the BufferRAM reserved for Data buffering
Sector	Part of a Page; 512B for the main data area and 16B for the spare area.  It is also the minimum Load/Program/Copy-Back Program unit during a 1~4 sector operation is available.
Data unit	Possible data unit to be read from memory to BufferRAM or to be programmed to memory 528B of which 512B is in main area and 16B in spare area 1056B of which 1024B is in main area and 32B in spare area 1584B of which 1536B is in main area and 48B in spare area 2112B of which 2048B is in main area and 64B in spare area

# 2.3 Pin Configuration



# (TOP VIEW, Balls Facing Down) 63ball FBGA OneNAND Chip

63ball, 9.5mm x 12mm x max 1.0mmt , 0.8mm ball pitch FBGA

# 2.4 Pin Description

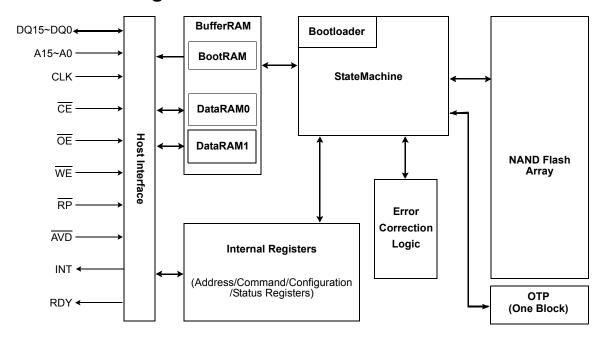
Pin Name	Type	Nameand Description
Host Interface		
A15~A0	I	Address Inputs - Inputs for addresses during read and write operation, which are for addressing BufferRAM & Register.
DQ15~DQ0	I/O	Data Inputs/Outputs - Inputs data during program and commands during all operations, outputs data for memory array/ register read cycles. Data pins float to high-impedance when the chip is deselected or outputs are disabled.
INT	0	Interrupt Notifies the Host when a command is completed. It is open drain output with internal resistor(~50kohms). After power-up, it is at hi-z condition. Once IOBE is set to 1, it does not float to hi-z condition even when the chip is deselected or when outputs are disabled.
RDY	0	
CLK	I	Clock CLK synchronizes the device to the system bus frequency in synchronous read mode. The first rising edge of CLK in conjunction with AVD low latches address input.
WE	I	Write Enable WE controls writes to the bufferRAM and registers. Datas are latched on the WE pulse's rising edge
ĀVD	I	Address Valid Detect Indicates valid address presence on address inputs. During asynchronous read operation, all addresses are latched on AVD's rising edge, and during synchronous read operation, all addresses are latched on CLK's rising edge while AVD is held low for one clock cycle.  > Low: for asynchronous mode, indicates valid address; for burst mode, causes starting address to be latched on rising edge on CLK  > High: device ignores address inputs
RP	I	Reset Pin When low, RP resets internal operation of OneNAND. RP status is don't care during power-up and bootloading.
CE	I	Chip Enable CE-low activates internal control logic, and CE-high deselects the device, places it in standby state, and places DQ in Hi-Z
ŌĒ	I	Output Enable  OE-low enables the device's output data buffers during a read cycle.
Power Supply		
Vcc-Core / Vcc		Power for MuxOneNAND Core This is the power supply for OneNAND Core.
Vcc-IO / Vccq		Power for MuxOneNAND I/O This is the power supply for OneNAND I/O Vcc-IO / Vccq is internally separated from Vcc-Core / Vcc.
Vss		Ground for OneNAND
etc.		
DNU		Do Not Use Leave it disconnected. These pins are used for testing.
NC		No Connection Lead is not internally connected.

#### NOTE

Do not leave power supply(Vcc-Core/Vcc-IO, VSS) disconnected.



### 2.5 Block Diagram



### 2.6 Memory Array Organization

The OneNAND architecture integrates several memory areas on a single chip.

## 2.6.1 Internal (NAND Array) Memory Organization

The on-chip internal memory is a single-level-cell (SLC) NAND array used for data storage and code. The internal memory is divided into a main area and a spare area.

#### Main Area

The main area is the primary memory array. This main area is divided into Blocks of 64 Pages. Within a Block, each Page is 2KB and is comprised of 4 Sectors. Within a Page, each Sector is 512B and is comprised of 256 Words.

#### Spare Area

The spare area is used for invalid block information and ECC storage. Spare area of internal memory is associated with corresponding main area of internal memory. Within a Block, each Page has four 16B Sectors of spare area. Each spare area Sector is 8 words.



### **Internal Memory Array Information**

Area	Block	Page	Sector
Main	128KB	2KB	512B
Spare	4KB	64B	16B

#### **Internal Memory Array Organization**

Sector							
Main Area	Spare Area						
512B	16B						



 Block

 Main Area
 Spare Area

 2KB Page0
 64B Page0
 Page 0

 ...
 ...
 ...

 2KB Page63
 64B Page63
 Page 63

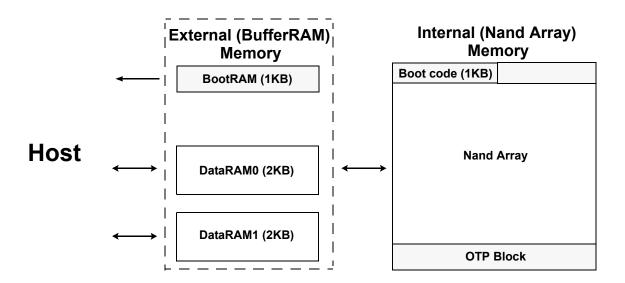
 128KB
 4KB

### 2.6.2 External (BufferRAM) Memory Organization

The on-chip external memory is comprised of 3 buffers used for Boot Code storage and data buffering.

The BootRAM is a 1KB buffer that receives Boot Code from the internal memory and makes it available to the host at start up.

There are two independent 2KB bi-directional data buffers, DataRAM0 and DataRAM1. These dual buffers enable the host to execute simultaneous Read-While load, and Write-While-program operations after Boot Up. During Boot Up, the BootRam is used by the host to initialize the main memory, and deliver boot code from NAND Flash core to host.

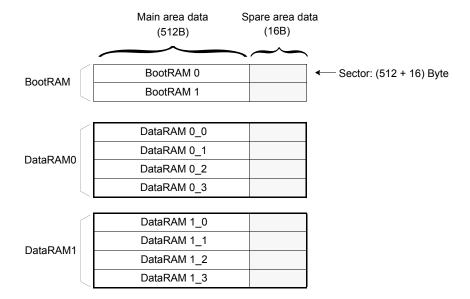


The external memory is divided into a main area and a spare area. Each buffer is the equivalent size of a Sector. The main area data is 512B. The spare area data is 16B.

### **External Memory Array Information**

Area		BootRAM	BootRAM DataRAM0	
Total Size		1KB+32B	2KB+64B	2KB+64B
Number of Sectors		2	4	4
Soctor	Main	512B	512B	512B
Sector	Spare	16B	16B	16B

### **External Memory Array Organization**



# 2.7 Memory Map

The following tables are the memory maps for the OneNAND.

### 2.7.1 Internal (NAND Array) Memory Organization

The following tables show the Internal Memory address map in word order.

Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block0	0000h	0000h~00FFh	128KB	Block32	0020h	0000h~00FFh	128KB
Block1	0001h	0000h~00FFh	128KB	Block33	0021h	0000h~00FFh	128KB
Block2	0002h	0000h~00FFh	128KB	Block34	0022h	0000h~00FFh	128KB
Block3	0003h	0000h~00FFh	128KB	Block35	0023h	0000h~00FFh	128KB
Block4	0004h	0000h~00FFh	128KB	Block36	0024h	0000h~00FFh	128KB
Block5	0005h	0000h~00FFh	128KB	Block37	0025h	0000h~00FFh	128KB
Block6	0006h	0000h~00FFh	128KB	Block38	0026h	0000h~00FFh	128KB
Block7	0007h	0000h~00FFh	128KB	Block39	0027h	0000h~00FFh	128KB
Block8	0008h	0000h~00FFh	128KB	Block40	0028h	0000h~00FFh	128KB
Block9	0009h	0000h~00FFh	128KB	Block41	0029h	0000h~00FFh	128KB
Block10	000Ah	0000h~00FFh	128KB	Block42	002Ah	0000h~00FFh	128KB
Block11	000Bh	0000h~00FFh	128KB	Block43	002Bh	0000h~00FFh	128KB
Block12	000Ch	0000h~00FFh	128KB	Block44	002Ch	0000h~00FFh	128KB
Block13	000Dh	0000h~00FFh	128KB	Block45	002Dh	0000h~00FFh	128KB
Block14	000Eh	0000h~00FFh	128KB	Block46	002Eh	0000h~00FFh	128KB
Block15	000Fh	0000h~00FFh	128KB	Block47	002Fh	0000h~00FFh	128KB
Block16	0010h	0000h~00FFh	128KB	Block48	0030h	0000h~00FFh	128KB
Block17	0011h	0000h~00FFh	128KB	Block49	0031h	0000h~00FFh	128KB
Block18	0012h	0000h~00FFh	128KB	Block50	0032h	0000h~00FFh	128KB
Block19	0013h	0000h~00FFh	128KB	Block51	0033h	0000h~00FFh	128KB
Block20	0014h	0000h~00FFh	128KB	Block52	0034h	0000h~00FFh	128KB
Block21	0015h	0000h~00FFh	128KB	Block53	0035h	0000h~00FFh	128KB
Block22	0016h	0000h~00FFh	128KB	Block54	0036h	0000h~00FFh	128KB
Block23	0017h	0000h~00FFh	128KB	Block55	0037h	0000h~00FFh	128KB
Block24	0018h	0000h~00FFh	128KB	Block56	0038h	0000h~00FFh	128KB
Block25	0019h	0000h~00FFh	128KB	Block57	0039h	0000h~00FFh	128KB
Block26	001Ah	0000h~00FFh	128KB	Block58	003Ah	0000h~00FFh	128KB
Block27	001Bh	0000h~00FFh	128KB	Block59	003Bh	0000h~00FFh	128KB
Block28	001Ch	0000h~00FFh	128KB	Block60	003Ch	0000h~00FFh	128KB
Block29	001Dh	0000h~00FFh	128KB	Block61	003Dh	0000h~00FFh	128KB
Block30	001Eh	0000h~00FFh	128KB	Block62	003Eh	0000h~00FFh	128KB
Block31	001Fh	0000h~00FFh	128KB	Block63	003Fh	0000h~00FFh	128KB

Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block64	0040h	0000h~00FFh	128KB	Block96	0060h	0000h~00FFh	128KB
Block65	0041h	0000h~00FFh	128KB	Block97	0061h	0000h~00FFh	128KB
Block66	0042h	0000h~00FFh	128KB	Block98	0062h	0000h~00FFh	128KB
Block67	0043h	0000h~00FFh	128KB	Block99	0063h	0000h~00FFh	128KB
Block68	0044h	0000h~00FFh	128KB	Block100	0064h	0000h~00FFh	128KB
Block69	0045h	0000h~00FFh	128KB	Block101	0065h	0000h~00FFh	128KB
Block70	0046h	0000h~00FFh	128KB	Block102	0066h	0000h~00FFh	128KB
Block71	0047h	0000h~00FFh	128KB	Block103	0067h	0000h~00FFh	128KB
Block72	0048h	0000h~00FFh	128KB	Block104	0068h	0000h~00FFh	128KB
Block73	0049h	0000h~00FFh	128KB	Block105	0069h	0000h~00FFh	128KB
Block74	004Ah	0000h~00FFh	128KB	Block106	006Ah	0000h~00FFh	128KB
Block75	004Bh	0000h~00FFh	128KB	Block107	006Bh	0000h~00FFh	128KB
Block76	004Ch	0000h~00FFh	128KB	Block108	006Ch	0000h~00FFh	128KB
Block77	004Dh	0000h~00FFh	128KB	Block109	006Dh	0000h~00FFh	128KB
Block78	004Eh	0000h~00FFh	128KB	Block110	006Eh	0000h~00FFh	128KB
Block79	004Fh	0000h~00FFh	128KB	Block111	006Fh	0000h~00FFh	128KB
Block80	0050h	0000h~00FFh	128KB	Block112	0070h	0000h~00FFh	128KB
Block81	0051h	0000h~00FFh	128KB	Block113	0071h	0000h~00FFh	128KB
Block82	0052h	0000h~00FFh	128KB	Block114	0072h	0000h~00FFh	128KB
Block83	0053h	0000h~00FFh	128KB	Block115	0073h	0000h~00FFh	128KB
Block84	0054h	0000h~00FFh	128KB	Block116	0074h	0000h~00FFh	128KB
Block85	0055h	0000h~00FFh	128KB	Block117	0075h	0000h~00FFh	128KB
Block86	0056h	0000h~00FFh	128KB	Block118	0076h	0000h~00FFh	128KB
Block87	0057h	0000h~00FFh	128KB	Block119	0077h	0000h~00FFh	128KB
Block88	0058h	0000h~00FFh	128KB	Block120	0078h	0000h~00FFh	128KB
Block89	0059h	0000h~00FFh	128KB	Block121	0079h	0000h~00FFh	128KB
Block90	005Ah	0000h~00FFh	128KB	Block122	007Ah	0000h~00FFh	128KB
Block91	005Bh	0000h~00FFh	128KB	Block123	007Bh	0000h~00FFh	128KB
Block92	005Ch	0000h~00FFh	128KB	Block124	007Ch	0000h~00FFh	128KB
Block93	005Dh	0000h~00FFh	128KB	Block125	007Dh	0000h~00FFh	128KB
Block94	005Eh	0000h~00FFh	128KB	Block126	007Eh	0000h~00FFh	128KB
Block95	005Fh	0000h~00FFh	128KB	Block127	007Fh	0000h~00FFh	128KB



Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block128	0080h	0000h~00FFh	128KB	Block160	00A0h	0000h~00FFh	128KB
Block129	0081h	0000h~00FFh	128KB	Block161	00A1h	0000h~00FFh	128KB
Block130	0082h	0000h~00FFh	128KB	Block162	00A2h	0000h~00FFh	128KB
Block131	0083h	0000h~00FFh	128KB	Block163	00A3h	0000h~00FFh	128KB
Block132	0084h	0000h~00FFh	128KB	Block164	00A4h	0000h~00FFh	128KB
Block133	0085h	0000h~00FFh	128KB	Block165	00A5h	0000h~00FFh	128KB
Block134	0086h	0000h~00FFh	128KB	Block166	00A6h	0000h~00FFh	128KB
Block135	0087h	0000h~00FFh	128KB	Block167	00A7h	0000h~00FFh	128KB
Block136	0088h	0000h~00FFh	128KB	Block168	00A8h	0000h~00FFh	128KB
Block137	0089h	0000h~00FFh	128KB	Block169	00A9h	0000h~00FFh	128KB
Block138	008Ah	0000h~00FFh	128KB	Block170	00AAh	0000h~00FFh	128KB
Block139	008Bh	0000h~00FFh	128KB	Block171	00ABh	0000h~00FFh	128KB
Block140	008Ch	0000h~00FFh	128KB	Block172	00ACh	0000h~00FFh	128KB
Block141	008Dh	0000h~00FFh	128KB	Block173	00ADh	0000h~00FFh	128KB
Block142	008Eh	0000h~00FFh	128KB	Block174	00AEh	0000h~00FFh	128KB
Block143	008Fh	0000h~00FFh	128KB	Block175	00AFh	0000h~00FFh	128KB
Block144	0090h	0000h~00FFh	128KB	Block176	00B0h	0000h~00FFh	128KB
Block145	0091h	0000h~00FFh	128KB	Block177	00B1h	0000h~00FFh	128KB
Block146	0092h	0000h~00FFh	128KB	Block178	00B2h	0000h~00FFh	128KB
Block147	0093h	0000h~00FFh	128KB	Block179	00B3h	0000h~00FFh	128KB
Block148	0094h	0000h~00FFh	128KB	Block180	00B4h	0000h~00FFh	128KB
Block149	0095h	0000h~00FFh	128KB	Block181	00B5h	0000h~00FFh	128KB
Block150	0096h	0000h~00FFh	128KB	Block182	00B6h	0000h~00FFh	128KB
Block151	0097h	0000h~00FFh	128KB	Block183	00B7h	0000h~00FFh	128KB
Block152	0098h	0000h~00FFh	128KB	Block184	00B8h	0000h~00FFh	128KB
Block153	0099h	0000h~00FFh	128KB	Block185	00B9h	0000h~00FFh	128KB
Block154	009Ah	0000h~00FFh	128KB	Block186	00BAh	0000h~00FFh	128KB
Block155	009Bh	0000h~00FFh	128KB	Block187	00BBh	0000h~00FFh	128KB
Block156	009Ch	0000h~00FFh	128KB	Block188	00BCh	0000h~00FFh	128KB
Block157	009Dh	0000h~00FFh	128KB	Block189	00BDh	0000h~00FFh	128KB
Block158	009Eh	0000h~00FFh	128KB	Block190	00BEh	0000h~00FFh	128KB
Block159	009Fh	0000h~00FFh	128KB	Block191	00BFh	0000h~00FFh	128KB



Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block192	00C0h	0000h~00FFh	128KB	Block224	00E0h	0000h~00FFh	128KB
Block193	00C1h	0000h~00FFh	128KB	Block225	00E1h	0000h~00FFh	128KB
Block194	00C2h	0000h~00FFh	128KB	Block226	00E2h	0000h~00FFh	128KB
Block195	00C3h	0000h~00FFh	128KB	Block227	00E3h	0000h~00FFh	128KB
Block196	00C4h	0000h~00FFh	128KB	Block228	00E4h	0000h~00FFh	128KB
Block197	00C5h	0000h~00FFh	128KB	Block229	00E5h	0000h~00FFh	128KB
Block198	00C6h	0000h~00FFh	128KB	Block230	00E6h	0000h~00FFh	128KB
Block199	00C7h	0000h~00FFh	128KB	Block231	00E7h	0000h~00FFh	128KB
Block200	00C8h	0000h~00FFh	128KB	Block232	00E8h	0000h~00FFh	128KB
Block201	00C9h	0000h~00FFh	128KB	Block233	00E9h	0000h~00FFh	128KB
Block202	00CAh	0000h~00FFh	128KB	Block234	00EAh	0000h~00FFh	128KB
Block203	00CBh	0000h~00FFh	128KB	Block235	00EBh	0000h~00FFh	128KB
Block204	00CCh	0000h~00FFh	128KB	Block236	00ECh	0000h~00FFh	128KB
Block205	00CDh	0000h~00FFh	128KB	Block237	00EDh	0000h~00FFh	128KB
Block206	00CEh	0000h~00FFh	128KB	Block238	00EEh	0000h~00FFh	128KB
Block207	00CFh	0000h~00FFh	128KB	Block239	00EFh	0000h~00FFh	128KB
Block208	00D0h	0000h~00FFh	128KB	Block240	00F0h	0000h~00FFh	128KB
Block209	00D1h	0000h~00FFh	128KB	Block241	00F1h	0000h~00FFh	128KB
Block210	00D2h	0000h~00FFh	128KB	Block242	00F2h	0000h~00FFh	128KB
Block211	00D3h	0000h~00FFh	128KB	Block243	00F3h	0000h~00FFh	128KB
Block212	00D4h	0000h~00FFh	128KB	Block244	00F4h	0000h~00FFh	128KB
Block213	00D5h	0000h~00FFh	128KB	Block245	00F5h	0000h~00FFh	128KB
Block214	00D6h	0000h~00FFh	128KB	Block246	00F6h	0000h~00FFh	128KB
Block215	00D7h	0000h~00FFh	128KB	Block247	00F7h	0000h~00FFh	128KB
Block216	00D8h	0000h~00FFh	128KB	Block248	00F8h	0000h~00FFh	128KB
Block217	00D9h	0000h~00FFh	128KB	Block249	00F9h	0000h~00FFh	128KB
Block218	00DAh	0000h~00FFh	128KB	Block250	00FAh	0000h~00FFh	128KB
Block219	00DBh	0000h~00FFh	128KB	Block251	00FBh	0000h~00FFh	128KB
Block220	00DCh	0000h~00FFh	128KB	Block252	00FCh	0000h~00FFh	128KB
Block221	00DDh	0000h~00FFh	128KB	Block253	00FDh	0000h~00FFh	128KB
Block222	00DEh	0000h~00FFh	128KB	Block254	00FEh	0000h~00FFh	128KB
Block223	00DFh	0000h~00FFh	128KB	Block255	00FFh	0000h~00FFh	128KB



Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block256	0100h	0000h~00FFh	128KB	Block288	0120h	0000h~00FFh	128KB
Block257	0101h	0000h~00FFh	128KB	Block289	0121h	0000h~00FFh	128KB
Block258	0102h	0000h~00FFh	128KB	Block290	0122h	0000h~00FFh	128KB
Block259	0103h	0000h~00FFh	128KB	Block291	0123h	0000h~00FFh	128KB
Block260	0104h	0000h~00FFh	128KB	Block292	0124h	0000h~00FFh	128KB
Block261	0105h	0000h~00FFh	128KB	Block293	0125h	0000h~00FFh	128KB
Block262	0106h	0000h~00FFh	128KB	Block294	0126h	0000h~00FFh	128KB
Block263	0107h	0000h~00FFh	128KB	Block295	0127h	0000h~00FFh	128KB
Block264	0108h	0000h~00FFh	128KB	Block296	0128h	0000h~00FFh	128KB
Block265	0109h	0000h~00FFh	128KB	Block297	0129h	0000h~00FFh	128KB
Block266	010Ah	0000h~00FFh	128KB	Block298	012Ah	0000h~00FFh	128KB
Block267	010Bh	0000h~00FFh	128KB	Block299	012Bh	0000h~00FFh	128KB
Block268	010Ch	0000h~00FFh	128KB	Block300	012Ch	0000h~00FFh	128KB
Block269	010Dh	0000h~00FFh	128KB	Block301 012Dh		0000h~00FFh	128KB
Block270	010Eh	0000h~00FFh	128KB	Block302	012Eh	0000h~00FFh	128KB
Block271	010Fh	0000h~00FFh	128KB	Block303	012Fh	0000h~00FFh	128KB
Block272	0110h	0000h~00FFh	128KB	Block304	0130h	0000h~00FFh	128KB
Block273	0111h	0000h~00FFh	128KB	Block305	0131h	0000h~00FFh	128KB
Block274	0112h	0000h~00FFh	128KB	Block306	0132h	0000h~00FFh	128KB
Block275	0113h	0000h~00FFh	128KB	Block307	0133h	0000h~00FFh	128KB
Block276	0114h	0000h~00FFh	128KB	Block308	0134h	0000h~00FFh	128KB
Block277	0115h	0000h~00FFh	128KB	Block309	0135h	0000h~00FFh	128KB
Block278	0116h	0000h~00FFh	128KB	Block310	0136h	0000h~00FFh	128KB
Block279	0117h	0000h~00FFh	128KB	Block311	0137h	0000h~00FFh	128KB
Block280	0118h	0000h~00FFh	128KB	Block312	0138h	0000h~00FFh	128KB
Block281	0119h	0000h~00FFh	128KB	Block313	0139h	0000h~00FFh	128KB
Block282	011Ah	0000h~00FFh	128KB	Block314	013Ah	0000h~00FFh	128KB
Block283	011Bh	0000h~00FFh	128KB	Block315	013Bh	0000h~00FFh	128KB
Block284	011Ch	0000h~00FFh	128KB	Block316	013Ch	0000h~00FFh	128KB
Block285	011Dh	0000h~00FFh	128KB	Block317	013Dh	0000h~00FFh	128KB
Block286	011Eh	0000h~00FFh	128KB	Block318	013Eh	0000h~00FFh	128KB
Block287	011Fh	0000h~00FFh	128KB	Block319	013Fh	0000h~00FFh	128KB



Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block320	0140h	0000h~00FFh	128KB	Block352	0160h	0000h~00FFh	128KB
Block321	0141h	0000h~00FFh	128KB	Block353	0161h	0000h~00FFh	128KB
Block322	0142h	0000h~00FFh	128KB	Block354	0162h	0000h~00FFh	128KB
Block323	0143h	0000h~00FFh	128KB	Block355	0163h	0000h~00FFh	128KB
Block324	0144h	0000h~00FFh	128KB	Block356	0164h	0000h~00FFh	128KB
Block325	0145h	0000h~00FFh	128KB	Block357	0165h	0000h~00FFh	128KB
Block326	0146h	0000h~00FFh	128KB	Block358	0166h	0000h~00FFh	128KB
Block327	0147h	0000h~00FFh	128KB	Block359	0167h	0000h~00FFh	128KB
Block328	0148h	0000h~00FFh	128KB	Block360	0168h	0000h~00FFh	128KB
Block329	0149h	0000h~00FFh	128KB	Block361	0169h	0000h~00FFh	128KB
Block330	014Ah	0000h~00FFh	128KB	Block362	016Ah	0000h~00FFh	128KB
Block331	014Bh	0000h~00FFh	128KB	Block363	016Bh	0000h~00FFh	128KB
Block332	014Ch	0000h~00FFh	128KB	Block364	016Ch	0000h~00FFh	128KB
Block333	014Dh	0000h~00FFh	128KB	Block365	016Dh	0000h~00FFh	128KB
Block334	014Eh	0000h~00FFh	128KB	Block366 016Eh		0000h~00FFh	128KB
Block335	014Fh	0000h~00FFh	128KB	Block367	016Fh	0000h~00FFh	128KB
Block336	0150h	0000h~00FFh	128KB	Block368	0170h	0000h~00FFh	128KB
Block337	0151h	0000h~00FFh	128KB	Block369	0171h	0000h~00FFh	128KB
Block338	0152h	0000h~00FFh	128KB	Block370	0172h	0000h~00FFh	128KB
Block339	0153h	0000h~00FFh	128KB	Block371	0173h	0000h~00FFh	128KB
Block340	0154h	0000h~00FFh	128KB	Block372	0174h	0000h~00FFh	128KB
Block341	0155h	0000h~00FFh	128KB	Block373	0175h	0000h~00FFh	128KB
Block342	0156h	0000h~00FFh	128KB	Block374	0176h	0000h~00FFh	128KB
Block343	0157h	0000h~00FFh	128KB	Block375	0177h	0000h~00FFh	128KB
Block344	0158h	0000h~00FFh	128KB	Block376	0178h	0000h~00FFh	128KB
Block345	0159h	0000h~00FFh	128KB	Block377	0179h	0000h~00FFh	128KB
Block346	015Ah	0000h~00FFh	128KB	Block378	017Ah	0000h~00FFh	128KB
Block347	015Bh	0000h~00FFh	128KB	Block379	017Bh	0000h~00FFh	128KB
Block348	015Ch	0000h~00FFh	128KB	Block380	017Ch	0000h~00FFh	128KB
Block349	015Dh	0000h~00FFh	128KB	Block381	017Dh	0000h~00FFh	128KB
Block350	015Eh	0000h~00FFh	128KB	Block382	017Eh	0000h~00FFh	128KB
Block351	015Fh	0000h~00FFh	128KB	Block383	017Fh	0000h~00FFh	128KB



Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block384	0180h	0000h~00FFh	128KB	Block416	01A0h	0000h~00FFh	128KB
Block385	0181h	0000h~00FFh	128KB	Block417	01A1h	0000h~00FFh	128KB
Block386	0182h	0000h~00FFh	128KB	Block418	01A2h	0000h~00FFh	128KB
Block387	0183h	0000h~00FFh	128KB	Block419	01A3h	0000h~00FFh	128KB
Block388	0184h	0000h~00FFh	128KB	Block420	01A4h	0000h~00FFh	128KB
Block389	0185h	0000h~00FFh	128KB	Block421	01A5h	0000h~00FFh	128KB
Block390	0186h	0000h~00FFh	128KB	Block422	01A6h	0000h~00FFh	128KB
Block391	0187h	0000h~00FFh	128KB	Block423	01A7h	0000h~00FFh	128KB
Block392	0188h	0000h~00FFh	128KB	Block424	01A8h	0000h~00FFh	128KB
Block393	0189h	0000h~00FFh	128KB	Block425	01A9h	0000h~00FFh	128KB
Block394	018Ah	0000h~00FFh	128KB	Block426	01AAh	0000h~00FFh	128KB
Block395	018Bh	0000h~00FFh	128KB	Block427	01ABh	0000h~00FFh	128KB
Block396	018Ch	0000h~00FFh	128KB	Block428	01ACh	0000h~00FFh	128KB
Block397	018Dh	0000h~00FFh	128KB	Block429	01ADh	0000h~00FFh	128KB
Block398	018Eh	0000h~00FFh	128KB	Block430	01AEh	0000h~00FFh	128KB
Block399	018Fh	0000h~00FFh	128KB	Block431	01AFh	0000h~00FFh	128KB
Block400	0190h	0000h~00FFh	128KB	Block432	01B0h	0000h~00FFh	128KB
Block401	0191h	0000h~00FFh	128KB	Block433	01B1h	0000h~00FFh	128KB
Block402	0192h	0000h~00FFh	128KB	Block434	01B2h	0000h~00FFh	128KB
Block403	0193h	0000h~00FFh	128KB	Block435	01B3h	0000h~00FFh	128KB
Block404	0194h	0000h~00FFh	128KB	Block436	01B4h	0000h~00FFh	128KB
Block405	0195h	0000h~00FFh	128KB	Block437	01B5h	0000h~00FFh	128KB
Block406	0196h	0000h~00FFh	128KB	Block438	01B6h	0000h~00FFh	128KB
Block407	0197h	0000h~00FFh	128KB	Block439	01B7h	0000h~00FFh	128KB
Block408	0198h	0000h~00FFh	128KB	Block440	01B8h	0000h~00FFh	128KB
Block409	0199h	0000h~00FFh	128KB	Block441	01B9h	0000h~00FFh	128KB
Block410	019Ah	0000h~00FFh	128KB	Block442	01BAh	0000h~00FFh	128KB
Block411	019Bh	0000h~00FFh	128KB	Block443	01BBh	0000h~00FFh	128KB
Block412	019Ch	0000h~00FFh	128KB	Block444	01BCh	0000h~00FFh	128KB
Block413	019Dh	0000h~00FFh	128KB	Block445	01BDh	0000h~00FFh	128KB
Block414	019Eh	0000h~00FFh	128KB	Block446	01BEh	0000h~00FFh	128KB
Block415	019Fh	0000h~00FFh	128KB	Block447	01BFh	0000h~00FFh	128KB

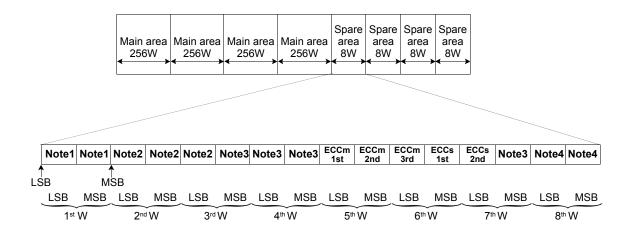


Block	Block Address	Page and Sector Address	Size	Block	Block Address	Page and Sector Address	Size
Block448	01C0h	0000h~00FFh	128KB	Block480	01E0h	0000h~00FFh	128KB
Block449	01C1h	0000h~00FFh	128KB	Block481	01E1h	0000h~00FFh	128KB
Block450	01C2h	0000h~00FFh	128KB	Block482	01E2h	0000h~00FFh	128KB
Block451	01C3h	0000h~00FFh	128KB	Block483	01E3h	0000h~00FFh	128KB
Block452	01C4h	0000h~00FFh	128KB	Block484	01E4h	0000h~00FFh	128KB
Block453	01C5h	0000h~00FFh	128KB	Block485	01E5h	0000h~00FFh	128KB
Block454	01C6h	0000h~00FFh	128KB	Block486	01E6h	0000h~00FFh	128KB
Block455	01C7h	0000h~00FFh	128KB	Block487	01E7h	0000h~00FFh	128KB
Block456	01C8h	0000h~00FFh	128KB	Block488	01E8h	0000h~00FFh	128KB
Block457	01C9h	0000h~00FFh	128KB	Block489	01E9h	0000h~00FFh	128KB
Block458	01CAh	0000h~00FFh	128KB	Block490	01EAh	0000h~00FFh	128KB
Block459	01CBh	0000h~00FFh	128KB	Block491	01EBh	0000h~00FFh	128KB
Block460	01CCh	0000h~00FFh	128KB	Block492	01ECh	0000h~00FFh	128KB
Block461	01CDh	0000h~00FFh	128KB	Block493 01EDh		0000h~00FFh	128KB
Block462	01CEh	0000h~00FFh	128KB	Block494	01EEh	0000h~00FFh	128KB
Block463	01CFh	0000h~00FFh	128KB	Block495	01EFh	0000h~00FFh	128KB
Block464	01D0h	0000h~00FFh	128KB	Block496 01F0h		0000h~00FFh	128KB
Block465	01D1h	0000h~00FFh	128KB	Block497	01F1h	0000h~00FFh	128KB
Block466	01D2h	0000h~00FFh	128KB	Block498	01F2h	0000h~00FFh	128KB
Block467	01D3h	0000h~00FFh	128KB	Block499	01F3h	0000h~00FFh	128KB
Block468	01D4h	0000h~00FFh	128KB	Block500	01F4h	0000h~00FFh	128KB
Block469	01D5h	0000h~00FFh	128KB	Block501	01F5h	0000h~00FFh	128KB
Block470	01D6h	0000h~00FFh	128KB	Block502	01F6h	0000h~00FFh	128KB
Block471	01D7h	0000h~00FFh	128KB	Block503	01F7h	0000h~00FFh	128KB
Block472	01D8h	0000h~00FFh	128KB	Block504	01F8h	0000h~00FFh	128KB
Block473	01D9h	0000h~00FFh	128KB	Block505	01F9h	0000h~00FFh	128KB
Block474	01DAh	0000h~00FFh	128KB	Block506	01FAh	0000h~00FFh	128KB
Block475	01DBh	0000h~00FFh	128KB	Block507	01FBh	0000h~00FFh	128KB
Block476	01DCh	0000h~00FFh	128KB	Block508	01FCh	0000h~00FFh	128KB
Block477	01DDh	0000h~00FFh	128KB	Block509	01FDh	0000h~00FFh	128KB
Block478	01DEh	0000h~00FFh	128KB	Block510	01FEh	0000h~00FFh	128KB
Block479	01DFh	0000h~00FFh	128KB	Block511	01FFh	0000h~00FFh	128KB



### 2.7.2 Internal Memory Spare Area Assignment

The figure below shows the assignment of the spare area in the Internal Memory NAND Array.



### Spare Area Assignment in the Internal Memory NAND Array Information

Word	Byte	Note	Description
1	LSB	1	Invalid Block information in 1st and 2nd page of an invalid block
'	MSB	'	Invalid Block information in 1st and 2nd page of an invalid block
2	LSB		
	MSB	2	Managed by internal ECC logic for Logical Sector Number data
3	LSB		
3	MSB		
4	LSB	3	Reserved for future use
4	MSB		
_	LSB		Dedicated to internal ECC logic. Read Only.  ECCm 1st for main area data
5	MSB		Dedicated to internal ECC logic. Read Only. ECCm 2nd for main area data
6	LSB		Dedicated to internal ECC logic. Read Only. ECCm 3rd for main area data
O	MSB		Dedicated to internal ECC logic. Read Only. ECCs 1st for 2nd word of spare area data
7	LSB		Dedicated to internal ECC logic. Read Only. ECCs 2nd for 3rd word of spare area data
	MSB	3	Reserved for future use
8	LSB	4	Available to the user
0	MSB	7	Available to the user

## 2.7.3 External Memory (BufferRAM) Address Map

The following table shows the External Memory address map in Word and Byte Order. Note that the data output is unknown while host reads a register bit of reserved area.

Division	Address (word order)	Address (byte order)	(1	Size total 128KI	3)	Usage	Description
Main area	0000h~00FFh	00000h~001FEh	512B	1KB	R	BootM 0	BootRAM Main sector0
(64KB)	0100h~01FFh	00200h~003FEh	512B	IND	K	BootM 1	BootRAM Main sector1
	0200h~02FFh	00400h~005FEh	512B			DataM 0_0	DataRAM Main page0/sector0
	0300h~03FFh	00600h~007FEh	512B			DataM 0_1	DataRAM Main page0/sector1
	0400h~04FFh	00800h~009FEh	512B			DataM 0_2	DataRAM Main page0/sector2
	0500h~05FFh	00A00h~00BFEh	512B	4KB	R/W	DataM 0_3	DataRAM Main page0/sector3
	0600h~06FFh	00C00h~00DFEh	512B	468	R/W	DataM 1_0	DataRAM Main page1/sector0
	0700h~07FFh	00E00h~00FFEh	512B			DataM 1_1	DataRAM Main page1/sector1
	0800h~08FFh	01000h~011FEh	512B			DataM 1_2	DataRAM Main page1/sector2
	0900h~09FFh	01200h~013FEh	512B			DataM 1_3	DataRAM Main page1/sector3
	0A00h~7FFFh	01400h~0FFFEh	59K	59K	-	Reserved	Reserved
Spare area	8000h~8007h	10000h~1000Eh	16B	32B	R	BootS 0	BootRAM Spare sector0
(8KB)	8008h~800Fh 10010h~1001	10010h~1001Eh	16B	325	K	BootS 1	BootRAM Spare sector1
	8010h~8017h	10020h~1002Eh	16B			DataS 0_0	DataRAM Spare page0/sector0
	8018h~801Fh	10030h~1003Eh	16B			DataS 0_1	DataRAM Spare page0/sector1
	8020h~8027h	10040h~1004Eh	16B		R/W	DataS 0_2	DataRAM Spare page0/sector2
	8028h~802Fh	10050h~1005Eh	16B	128B		DataS 0_3	DataRAM Spare page0/sector3
	8030h~8037h	10060h~1006Eh	16B	1200	FC/VV	DataS 1_0	DataRAM Spare page1/sector0
	8038h~803Fh	10070h~1007Eh	16B			DataS 1_1	DataRAM Spare page1/sector1
	8040h~8047h	10080h~1008Eh	16B			DataS 1_2	DataRAM Spare page1/sector2
	8048h~804Fh	10090h~1009Eh	16B			DataS 1_3	DataRAM Spare page1/sector3
	8050h~8FFFh	100A0h~11FFEh	8032B	8032B	-	Reserved	Reserved
Reserved (24KB)	9000h~BFFFh	12000h~17FFEh	24KB	24KB	-	Reserved	Reserved
Reserved (8KB)	C000h~CFFFh	18000h~19FFEh	8KB	8KB	-	Reserved	Reserved
Reserved (16KB)	D000h~EFFFh	1A000h~1DFFEh	16KB	16KB	-	Reserved	Reserved
Registers (8KB)	F000h~FFFFh	1E000h~1FFFEh	8KB	8KB	R or R/W	Registers	Registers



### 2.7.4 External Memory Detail Map Information

The tables below show Word Order Address Map information for the BootRAM and DataRAM main and spare areas.

#### • BootRAM(Main area)

-0000h~01FFh: 2(sector) x 512byte(NAND main area) = 1KB

0000h~00FFh(512B)	0100h~01FFh(512B)
BootM 0	BootM 1
(sector 0 of page 0)	(sector 1 of page 0)

### • DataRAM(Main area)

-0200h~09FFh: 8(sector) x 512byte(NAND main area) = 4KB

0200h~02FFh(512B)	0300h~03FFh(512B)	0400h~04FFh(512B)	0500h~05FFh(512B)
DataM 0_0	DataM 0_1	DataM 0_2	DataM 0_3
(sector 0 of page 0)	(sector 1 of page 0)	(sector 2 of page 0)	(sector 3 of page 0)
0600h~06FFh(512B) DataM 1_0 (sector 0 of page 1)	0700h~07FFh(512B) DataM 1_1 (sector 1 of page 1)	0800h~08FFh(512B) DataM 1_2 (sector 2 of page 1)	0900h~09FFh(512B) DataM 1_3 (sector 3 of page 1)

### • BootRAM(Spare area)

-8000h~800Fh: 2(sector) x 16byte(NAND spare area) = 32B

8000h~8007h(16B)	8008h~800Fh(16B)
BootS 0	BootS 1
(sector 0 of page 0)	(sector 1 of page 0)

### • DataRAM(Spare area)

-8010h~804Fh: 8(sector) x 16byte(NAND spare area) = 128B

8010h~8017h(16B)	8018h~801Fh(16B)	8020h~8027h(16B)	8028h~802Fh(16B)
DataS 0_0	DataS 0_1	DataS 0_2	DataS 0_3
(sector 0 of page 0)	(sector 1 of page 0)	(sector 2 of page 0)	(sector 3 of page 0)
8030h~8037h(16B)	8038h~803Fh(16B)	8040h~8047h(16B)	8048h~804Fh(16B)
DataS 1_0	DataS 1_1	DataS 1_2	DataS 1_3
(sector 0 of page 1)	(sector 1 of page 1)	(sector 2 of page 1)	(sector 3 of page 1)

<sup>\*</sup>NAND Flash array consists of 2KB page size and 128KB block size.



# 2.7.5 External Memory Spare Area Assignment

Equivalent to 1word of NAND Flash

			<b>-</b>					Equiv	alent to	1word	of NAN	ID Flas	sh					
Buf.	Word Address	Byte Address	F	E	D	С	В	Α	9	8	7	6	5	4	3	2	1	0
BootS 0	8000h	10000h		1						Е	31							
	8001h	10002h						Ма	naged	by In	ternal	ECC I	ogic					
	8002h	10004h		F	Reserv	ed for	the fu	iture u	se			Ма	naged	by In	ternal	ECC I	ogic	
	8003h	10006h					ı	Reserv	ed for	the cu	irrent a	and fut	ture us	se				
	8004h	10008h		ECC	Code	for M	lain ar	ea data	a (2 <sup>nd</sup> )			ECC	Code	e for M	ain are	ea data	a (1st)	
	8005h	1000Ah		ECC Code for Spare area data (1st)  FFh(Reserved for the future use)  ECC Code for Main area data  ECC Code for Spare area data												a (3 <sup>rd</sup> )		
	8006h	1000Ch														ea dat	a (2 <sup>nd</sup> )	
	8007h	1000Eh			Free Usage													
BootS 1	8008h	10010h			BI													
	8009h	10012h			Managed by Internal ECC logic													
	800Ah	10014h		F	Reserved for the future use Managed by Internal ECC logic											ogic		
	800Bh	10016h			Reserved for the current and future use													
	800Ch	10018h		ECC	Code	for M	lain ar	ea data	a (2 <sup>nd</sup> )			ECC	Code	for M	ain are	ea dat	a (1 <sup>st</sup> )	
	800Dh	1001Ah		ECC	Code	for Sp	pare a	rea da	a (1 <sup>st</sup> )			ECC	Code	for M	ain are	ea data	a (3 <sup>rd</sup> )	
	800Eh	1001Ch		FFh	(Rese	rved f	for the	future	use)			ECC	Code	for Sp	are ar	ea dat	a (2 <sup>nd</sup> )	
	800Fh	1001Eh								Free	Usage							
DataS	8010h	10020h								E	31							
0_0	8011h	10022h						Ma	naged	by In	ternal	ECC I	ogic					
	8012h	10024h		F	Reserv	ed for	r the fu	iture u	se			Ма	naged	by In	ternal	ECC I	ogic	
	8013h	10026h					ı	Reserv	ed for	the cu	irrent a	and fut	ture us	se				
	8014h	10028h		ECC	Code	for M	lain ar	ea data	a (2 <sup>nd</sup> )			ECC	Code	e for M	ain are	ea data	a (1 <sup>st</sup> )	
	8015h	1002Ah		ECC	Code	for Sp	pare a	rea da	a (1 <sup>st</sup> )			ECC	Code	e for M	ain are	ea data	a (3 <sup>rd</sup> )	
	8016h	1002Ch		FFh	(Rese	rved f	for the	future	use)			ECC	Code	for Sp	are ar	ea dat	a (2 <sup>nd</sup> )	
	8017h	1002Eh								Free	Usage	!						
DataS	8018h	10030h								E	31							
0_1	8019h	10032h						Ma	naged	by In	ternal	ECC I	ogic					
	801Ah	10034h		Reserved for the future use Managed by Internal ECC logic										ogic				
	801Bh	10036h					ı	Reserv	ed for	the cu	irrent a	and fut	ture us	se				
	801Ch	10038h		ECC	Code	for M	lain ar	ea data	a (2 <sup>nd</sup> )			ECC	Code	e for M	ain are	ea data	a (1 <sup>st</sup> )	
	801Dh	1003Ah		ECC	Code	for Sp	pare a	rea da	a (1st)			ECC	Code	for M	ain are	ea data	a (3 <sup>rd</sup> )	
	801Eh	1003Ch		FFh	(Rese	rved f	for the	future	use)			ECC	Code	for Sp	are ar	ea dat	a (2 <sup>nd</sup> )	
	801Fh	1003Eh								Free	Usage							

Buf.	Word Address	Byte Address	F	E	D	С	В	Α	9	8	7	6	5	4	3	2	1	0		
DataS 0_2	8020h	10040h						1		E	3I						1			
	8021h	10042h						Ма	naged	by Int	ternal	ECC I	ogic							
	8022h	10044h		F	Reserv	ed for	the fu	ture us	se			Ма	naged	by Int	ernal	ECC I	ogic			
	8023h	10046h					F	Reserv	ed for	the cu	rrent a	and fut	ture us	se						
	8024h	10048h		ECC	Code	for M	lain are	ea data	a (2 <sup>nd</sup> )			ECC	CC Code for Main area data (1 <sup>st</sup> )							
	8025h	1004Ah		ECC	Code	for Sp	oare ar	ea dat	a (1st)			ECC Code for Main area data (3 <sup>rd</sup> )								
	8026h	1004Ch		F	Reserv	ed for	the fu	ture us	se			ECC	Code	for Sp	are ar	e area data (2 <sup>nd</sup> )				
	8027h	1004Eh		Free Usage																
DataS 0_3	8028h	10050h		ВІ																
	8029h	10052h			Managed by Internal ECC logic															
	802Ah	10054h		F	Reserved for the future use Managed by Internal									ECC I	ogic					
	802Bh	10056h		Reserved for the current and future use																
	802Ch	10058h		ECC	ECC Code for Main area data (2 <sup>nd</sup> ) ECC Code for Main area data										a (1 <sup>st</sup> )					
	802Dh	1005Ah		ECC	ECC Code for Spare area data (1st) ECC Code for Main area da										ea dat	a (3 <sup>rd</sup> )				
	802Eh	1005Ch		F	Reserved for the future use ECC Code for Spa										are ar	ea dat	a (2 <sup>nd</sup> )			
	802Fh	1005Eh		Free Usage																
DataS 1_0	8030h	10060h								E	31									
	8031h	10062h						Ма	naged	by Int	ternal	ECC I	ogic							
	8032h	10064h		F	Reserv	ed for	the fu	ture us	se			Ма	naged	by Int	ernal	ECC I	ogic			
	8033h	10066h		Reserved for the current and future use																
	8034h	10068h		ECC	Code	for M	lain are	ea data	a (2 <sup>nd</sup> )		ECC Code for Main area data (1st)									
	8035h	1006Ah		ECC	Code	for Sp	oare ar	ea dat	a (1st)		ECC Code for Main area data (3rd)									
	8036h	1006Ch		F	Reserv	ed for	the fu	ture us	se			ECC	Code	for Sp	are ar	ea dat	a (2 <sup>nd</sup> )			
	8037h	1006Eh								Free	Usage									
DataS 1_1	8038h	10070h								Е	31									
	8039h	10072h						Ma	naged	by Int	ternal	ECC I	ogic							
	803Ah	10074h		F	Reserv	ed for	the fu	ture us	se			Ма	naged	by Int	ernal	ECC I	ogic			
	803Bh	10076h					F	Reserv	ed for	the cu	irrent a	and fut	ture us	se						
	803Ch	10078h		ECC	Code	for M	lain are	ea data	a (2 <sup>nd</sup> )			ECC	Code	for M	ain are	ea dat	a (1 <sup>st</sup> )			
	803Dh	1007Ah		ECC	Code	for S	oare a	ea dat	a (1st)			ECC	Code	for Ma	ain are	ea dat	a (3 <sup>rd</sup> )			
	803Eh	1007Ch		F	Reserv	ed for	the fu	ture us	se			ECC	Code	for Sp	are ar	ea dat	a (2 <sup>nd</sup> )			
	803Fh	1007Eh								Free	Usage									
DataS 1_2	8040h	10080h								Е	31									
	8041h	10082h		Managed by Int							ternal	ECC I	ogic							
	8042h	10084h		Reserved for the future use								Ма	naged	by Int	ernal	ECC I	ogic			
	8043h	10086h					F	Reserv	ed for	the cu	irrent a	and fut	ture us	se						
	8044h	10088h	0088h ECC				CC Code for Main area data (2 <sup>nd</sup> )					ECC Code for Main area data (1st)								
	8045h	1008Ah		ECC	Code	for S	oare a	ea dat	a (1st)			ECC	Code	for Ma	ain are	ea dat	a (3 <sup>rd</sup> )			
	8046h	1008Ch		F	Reserv	ed for	the fu	ture us	se			ECC	Code	for Sp	are ar	ea dat	a (2 <sup>nd</sup> )			
	8047h	1008Eh								Free	Usage									



			<b>«</b> —					Equiva	lent to	1word	of NAN	D Flas	h					<b>~</b>
Buf.	Word Address	Byte Address	F	E	D	С	В	Α	9	8	7	6	5	4	3	2	1	0
DataS 1_3	8048h	10090h								E	31							
	8049h	10092h		Managed by Internal ECC logic														
	804Ah	10094h		F	Reserv	ed for	the fu	ture us	e			Ма	naged	by Int	ernal	ECC I	ogic	
	804Bh	10096h					F	Reserv	ed for	the cu	irrent a	and fut	ure us	se				
	804Ch	10098h		ECC	Code	for M	ain are	ea data	(2 <sup>nd</sup> )			ECC	Code	for Ma	ain are	ea data	a (1 <sup>st</sup> )	
	804Dh	1009Ah		ECC	Code	for Sp	are ar	ea dat	a (1st)			ECC	Code	for Ma	ain are	ea data	a (3 <sup>rd</sup> )	
	804Eh	1009Ch		Reserved for the future use ECC Code for Spare area data (2 <sup>nd</sup> )														
	804Fh	1009Eh		Free Usage														

#### NOTE:

<sup>-</sup> BI: Bad block Information

<sup>&</sup>gt;Host can use complete spare area except BI and ECC code area. For example,

Host can write data to Spare area buffer except for the area controlled by ECC logic at program operation.

<sup>&</sup>gt;In case of 'with ECC' mode, OneNAND automatically generates ECC code for both main and spare data of memory during program operation, but does not update ECC code to spare bufferRAM during load operation.

<sup>&</sup>gt;When loading/programming spare area, spare area BufferRAM address(BSA) and BufferRAM sector count(BSC) is chosen via Start buffer register as it is.

# 2.8 Registers

Section 2.8 of this specification provides information about the OneNAND512 registers.

## 2.8.1 Register Address Map

This map describes the register addresses, register name, register description, and host accessibility.

Address (word order)	Address (byte order)	Name	Host Access	Description
F000h	1E000h	Manufacturer ID	R	Manufacturer identification
F001h	1E002h	Device ID	R	Device identification
F002h	1E004h	Version ID	R	N/A
F003h	1E006h	Data Buffer size	R	Data buffer size
F004h	1E008h	Boot Buffer size	R	Boot buffer size
F005h	1E00Ah	Amount of buffers	R	Amount of data/boot buffers
F006h	1E00Ch	Technology	R	Info about technology
F007h~F0FFh	1E00Eh~1E1FEh	Reserved	-	Reserved for user
F100h	1E200h	Start address 1	R/W	NAND Flash Block Address
F101h	1E202h	Start address 2	R/W	Reserved
F102h	1E204h	Start address 3	R/W	Destination Block address for Copy back program
F103h	1E206h	Start address 4	R/W	Destination Page & Sector address for Copy back program
F104h	1E208h	Start address 5	-	N/A
F105h	1E20Ah	Start address 6	-	N/A
F106h	1E20Ch	Start address 7	-	N/A
F107h	1E20Eh	Start address 8	R/W	NAND Flash Page & Sector address
F108h~F1FFh	1E210h~1E3FEh	Reserved	-	Reserved for user
F200h	1E400h	Start Buffer	R/W	Buffer Number for the page data transfer to/from the memory and the start Buffer Address The meaning is with which buffer to start and how many buffers to use for the data transfer
F201h~F207h	1E402h~1E40Eh	Reserved	-	Reserved for user
F208h~F21Fh	1E410h~1E43Eh	Reserved	-	Reserved for vendor specific purposes
F220h	1E440h	Command	R/W	Host control and memory operation commands
F221h	1E442h	System Configuration 1	R, R/W	memory and Host Interface Configuration
F222h	1E444h	System Configuration 2	-	N/A
F223h~F22Fh	1E446h~1E45Eh	Reserved	-	Reserved for user
F230h~F23Fh	1E460h~1E47Eh	Reserved	-	Reserved for vendor specific purposes
F240h	1E480h	Controller Status	R	Controller Status and result of memory operation
F241h	1E482h	Interrupt	R/W	Memory Command Completion Interrupt Status
F242h~F24Bh	1E484h~1E496h	Reserved	-	Reserved for user
F24Ch	1E498h	Start Block Address	R/W	Start memory block address in Write Protection mode



Address (word order)	Address (byte order)	Name	Host Access	Description
F24Dh	1E49Ah	Reserved	R/W	Reserved for user
F24Eh	1E49Ch	Write Protection Status	R	Current memory Write Protection status (unlocked/locked/tight-locked)
F24Fh~FEFFh	1E49Eh~1FDFEh	Reserved	-	Reserved for user
FF00h	1FE00h	ECC Status Register	R	ECC status of sector
FF01h	1FE02h	ECC Result of main area data	R	ECC error position of Main area data error for first selected Sector
FF02h	1FE04h	ECC Result of spare area data	R	ECC error position of Spare area data error for first selected Sector
FF03h	1FE06h	ECC Result of main area data	R	ECC error position of Main area data error for second selected Sector
FF04h	1FE08h	ECC Result of spare area data	R	ECC error position of Spare area data error for second selected Sector
FF05h	1FE0Ah	ECC Result of main area data	R	ECC error position of Main area data error for third selected Sector
FF06h	1FE0Ch	ECC Result of spare area data	R	ECC error position of Spare area data error for third selected Sector
FF07h	1FE0Eh	ECC Result of main area data	R	ECC error position of Main area data error for fourth selected Sector
FF08h	1FE10h	ECC Result of spare area data	R	ECC error position of Spare area data error for fourth selected Sector
FF09h~FFFFh	1FE12h~1FFFEh	Reserved	-	Reserved for vendor specific purposes

## 2.8.2 Manufacturer ID Register F000h (R)

This Read register describes the manufacturer's identification. Samsung Electronics Company manufacturer's ID is 00ECh.

F000h, default = 00ECh

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								nufID							

## 2.8.3 Device ID Register F001h (R)

This Read register describes the device.

F001h, see table for default.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	DeviceID														

#### **Device Identification**

Device Identification	Description
DeviceID [1:0] Vcc	00 = 1.8V, 01 = 2.65V/3.3V, 10/11 = reserved
DeviceID [2] Muxed/Demuxed	0 = Muxed, 1 = Demuxed
DeviceID [3] Single/DPP	0 = Single, 1 = DDP
DeviceID [6:4] Density	000 = 128Mb, 001 = 256Mb, 010 = 512Mb, 011 = 1Gb, 100 = 2Gb

#### **Device ID Default**

Device	DeviceID[15:0]
KFG1216Q2A	0024h
KFG1216D2A	0025h
KFG1216U2A	0025h



# 2.8.4 Version ID Register F002h

This register is reserved for manufacturer

## 2.8.5 Data Buffer Size Register F003h (R)

This Read register describes the size of the Data Buffer.

### F003h, default = 0800h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		DataBufSize													

#### **Data Buffer Size Information**

Version Identification	Description
DataBufSize	Total data buffer size in Words equal to 2 buffers of 1024 Words each $(2 \times 1024 = 2^{11})$ in the memory interface

## 2.8.6 Boot Buffer Size Register F004h (R)

This Read register describes the size of the Boot Buffer.

#### F004h, default = 0200h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							BootB	ufSize							

Register Information	Description
BootBufSize	Total boot buffer size in Words equal to 1 buffer of 512 Words $(1 \times 512 = 2^9)$ in the memory interface



## 2.8.7 Number of Buffers Register F005h (R)

This Read register describes the number of each Buffer.

F005h, default = 0201h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	DataBufAmount										BootBu	fAmount			

#### **Number of Buffers Information**

Register Information	Description
DataBufAmount	The number of data buffers = 2 (2 <sup>N</sup> , N=1)
BootBufAmount	The number of boot buffers = 1 (2 <sup>N</sup> , N=0)

## 2.8.8 Technology Register F006h (R)

This Read register describes the internal NAND array technology.

F006h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Te	ch							

### **Technology Information**

Technology	Register Setting
NAND SLC	0000h
NAND MLC	0001h
Reserved	0002h ~ FFFFh



### 2.8.9 Start Address1 Register F100h (R/W)

This Read/Write register describes the start address of the NAND Flash block address which will be loaded, programmed, or erased.

### F100h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved(00000)										FBA					

Device	Number of Block	FBA			
512Mb	512	FBA[8:0]			

#### Start Address1 Information

Register Information	Description
FBA	NAND Flash Block Address

## 2.8.10 Start Address2 Register F101h (R/W)

This register is reserved for future use.

# 2.8.11 Start Address3 Register F102h (R/W)

This Read/Write register describes the NAND Flash destination block address which will be copy back programmed.

### F102h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Rese	rved(000	0000)							FCBA				

Device	Number of Block	FBA
512Mb	512	FCBA[8:0]

### **Start Address3 Information**

Register Information	Description
FCBA	NAND Flash Copy Back Block Address

# 2.8.12 Start Address4 Register F103h (R/W)

This Read/Write register describes the NAND Flash destination page address in a block and the NAND Flash destination sector address in a page for copy back programming.

### F103h, default = 0000h

Ī	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Ī			R	eserved(	0000000	0)				FCPA					FC	SA

### **Start Address4 Information**

Item	Description	Default Value	Range
FCPA	NAND Flash Copy Back Page Address	000000	000000 ~ 111111, 6 bits for 64 pages
FCSA	NAND Flash Copy Back Sector Address	00	00 ~ 11, 2 bits for 4 sectors



# 2.8.13 Start Address5 Register F104h

This register is reserved for future use.

# 2.8.14 Start Address6 Register F105h

This register is reserved for future use.

# 2.8.15 Start Address7 Register F106h

This register is reserved for future use.

# 2.8.16 Start Address8 Register F107h (R/W)

This Read/Write register describes the NAND Flash start page address in a block for a page load, copy back program, or program operation and the NAND Flash start sector address in a page for a load, copy back program, or program operation.

### F107h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	·	Re	served (	(000000	00)					FI	PA			FSA	

### Start Address8 Information

Item	Description	Default Value	Range
FPA	NAND Flash Page Address	000000	000000 ~ 111111, 6 bits for 64 pages
FSA	NAND Flash Sector Address	00	00 ~ 11, 2 bits for 4 sectors



# 2.8.17 Start Buffer Register F200h (R/W)

This Read/Write register describes the BufferRAM Sector Count (BSC) and BufferRAM Sector Address (BSA).

The BufferRAM Sector Count (BSC) field specifies the number of sectors to be loaded, programmed, or copy back programmed. At 00 value (the default value), the number of sector is "4". If the internal RAM buffer reaches its maximum value of 11, it will count up to 0 value to meet the BSC value. For example, if BSA = 1101, BSC = 00, then the selected BufferRAM will count up from '1101  $\rightarrow$  1110  $\rightarrow$  1111  $\rightarrow$  1100'.

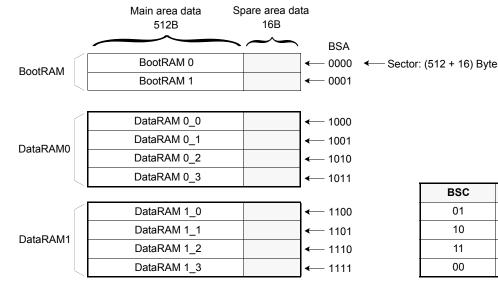
The BufferRAM Sector Address (BSA) is the sector 0~3 address in the internal BootRAM and DataRAM where data is placed.

### F200h, default = 0000h

Ī	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved(0000)				BSA				F	Reserved	00000)b	0)		BS		

#### Start Address8 Information

Item	Description
BSA[3]	Selection bit between BootRAM and DataRAM
BSA[2]	Selection bit between DataRAM0 and DataRAM1
BSA[1:0]	Selection bit between Sector0 and Sector1 in the internal BootRAM Selection bit between Sector0 to Sector3 in the internal DataRAM



BSC	Number of Sectors
01	1 sector
10	2 sector
11	3 sector
00	4 sector

# 2.8.18 Command Register F220h (R/W)

This Read/Write register describes the operation of the OneNAND interface.

Note that all commands should be issued right after INT is turned from ready state to busy state. (i.e. right after 0 is written to INT register.) After any command is issued and the corresponding operation is completed, INT goes back to ready state. (00F0h and 00F3h may be accepted during busy state of some operations. Refer to the rightmost column of the command register table below.)

### F220h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Com	mand							

CMD	Operation	Acceptable command during busy
0000h	Load single/multiple sector data unit into buffer	00F0h, 00F3h
0013h	Load single/multiple spare sector into buffer	00F0h, 00F3h
0080h	Program single/multiple sector data unit from buffer <sup>1)</sup>	00F0h, 00F3h
001Ah	Program single/multiple spare data unit from buffer	00F0h, 00F3h
001Bh	Copy back Program operation	00F0h, 00F3h
0023h	Unlock NAND array a block	-
002Ah	Lock NAND array a block	-
002Ch	Lock-tight NAND array a block	
0071h	Erase Verify Read	00F0h, 00F3h
0094h	Block Erase	00F0h, 00F3h
0095h	Multi-Block Erase	00F0h, 00F3h
00B0h	Erase Suspend	-
0030h	Erase Resume	00F0h, 00F3h
00F0h	Reset NAND Flash Core	-
00F3h	Reset OneNAND <sup>2)</sup>	-
0065h	OTP Access	00F0h, 00F3h

### NOTE:



<sup>1) 0080</sup>h programs both main and spare area, while 001Ah programs only spare area. Refer to chapter 5.8 for NOP limits in issuing these commands. When using 0080h and 001Ah command, Read-only part in spare area must be masked by FF. (Refer to chapter 2.7.2)

<sup>2) &#</sup>x27;Reset OneNAND' (=Hot reset) command makes the registers and NAND Flash core into default state as the warm reset(=reset by RP pin).

# 2.8.19 System Configuration 1 Register F221h (R, R/W)

This Read/Write register describes the system configuration.

### F221h, default = 40C0h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
R/W		R/W			R/W		R/W	R/W	R/W	R/W	R/W		R		R
RM		BRL			BL		ECC	RDY pol	INT pol	IOBE	RDY Conf	Res	served(	000)	BWPS

## Read Mode (RM)

RM	Read Mode	
0	Asynchronous read(default)	
1	Synchronous read	

## Read Mode Information[15]

Item	Definition	Description
RM	Read Mode	Selects between asynchronous read mode and synchronous read mode

## **Burst Read Latency (BRL)**

BRL	Latency Cycles	
000	8(N/A)	
001	9(N/A)	
010	10(N/A)	
011	3(up to 40MHz)	
100	4(default, min.)	
101	5	
110	6	
111	7	

## Burst Read Latency (BRL) Information[14:12]

Item	Definition	Description
BRL	Burst Read Latency	Specifies the access latency in the burst read transfer after the initial access



## Burst Length (BL)

BL	Burst Length(Main)	Burst Length(Spare)	
000	Continuous(default)		
001	001 4 words		
010	8 words		
011	16 words		
100	32 words	N/A	
101~111	Reserved		

## Burst Length (BL) Information[11:9]

Item	Definition	Description
BL	Burst Length	Specifies the size of the burst length during a synchronous read, wrap around and linear burst read

## Error Correction Code (ECC) Information[8]

Item	Definition	Description
ECC	Error Correction Code Operation	0 = with correction (default) 1 = without correction (bypassed)

## RDY Polarity (RDYpol) Information[7]

Item	Definition	Description
RDYpol	RDY signal polarity	1 = high for ready (default) 0 = low for ready

## INT Polarity (INTpol) Information[6]

INTpol	INT bit of Interrupt Status Register	INT Pin output
0	0 (busy)	High
ľ	1 (ready)	Low
1 (default)	0 (busy)	Low
i (delauit)	1 (ready)	High



## I/O Buffer Enable (IOBE)

IOBE is the I/O Buffer Enable for the INT and RDY signals. At startup, INT and RDY outputs are High-Z. Bits 6 and 7 become valid after IOBE is set to "1". IOBE can be reset by a Cold Reset or by writing "0" to bit 5 of System Configuration1 Register.

## I/O Buffer Enable Information[5]

Item	Definition	Description
IOBE	I/O Buffer Enable for INT and RDY signals	0 = disable (default) 1 = enable

## **RDY Configuration (RDY conf)**

## **RDY Configuration Information[4]**

ltem	Definition	Description
RDY conf	RDY configuration	0=active one clock before valid data(default) 1=active with valid data

### **Boot Buffer Write Protect Status (BWPS)**

### Boot Buffer Write Protect Status Information[0]

Item	Definition	Description
BWPS	Boot Buffer Write Protect Status	0 = locked (fixed)

# 2.8.20 System Configuration 2 Register F222h

This register is reserved for future use.

# 2.8.21 Controller Status Register F240h (R)

This Read register shows the overall internal status of the OneNAND and the controller.

#### F240h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OnG	Lock	Load	Prog	Erase	Error	Sus	Reserv ed(0)	RSTB	OTPL		Rese	erved(00	0000)		TO (0)

### OnGo

This bit shows the overall internal status of the OneNAND device.

### OnGo Information[15]

ltem	Definition	Description
OnGo	Internal Device Status	0 = ready 1 = busy

### Lock

This bit shows whether the host is loading data from the NAND Flash array into the locked BootRAM or whether the host is performing a program/erase of a locked block of the NAND Flash array.

## Lock Information[14]

Lock	Locked/Unlocked Check Result
0	Unlocked
1	Locked

### Load

This bit shows the Load Operation status.

### Load Information[13]

Item	Definition	Description
Load	Load Operation status	0 = ready (default) 1 = busy or error (see controller status output modes)



## **Program**

This bit shows the Program Operation status.

## **Program Information[12]**

Item	Definition	Description
Prog	Program Operation status	0 = ready (default) 1 = busy or error (see controller status output modes)

### **Erase**

This bit shows the Erase Operation status.

## Erase Information[11]

Item	Definition	Description
Erase	Erase Operation status	0 = ready (default) 1 = busy or error (see controller status output modes)

### Error

This bit shows the overall Error status, including Load Reset, Program Reset, and Erase Reset status.

## Error Information[10]

Error	Current Sector/Page Load/Program/CopyBack. Program/ Erase Result and Invalid Command Input
0	Pass
1	Fail

## Erase Suspend (Sus)

This bit shows the Erase Suspend status.

## Sus Information[9]

Sus	Erase Suspend Status
0	Erase Resume(Default)
1	Erase Suspend, Program Ongoing(Susp.), Load Ongoing(Susp.), Program Fail(Susp.), Load Fail(Susp.), Invalid Command(Susp.)



## Reset / Busy (RSTB)

This bit shows the Reset Operation status.

## RSTB Information[7]

Item	Definition	Description
RSTB	Reset Operation Status	0 = ready (default) 1 = busy (see controller status output modes)

## OTP Lock Status (OTP<sub>L</sub>)

This bit shows whether the OTP block is locked or unlocked. Locking the OTP has the effect of a 'write-protect' to guard against accidental re-programming of data stored in the OTP block.

The  $\mathsf{OTP}_{\mathsf{L}}$  status bit is automatically updated at power-on.

## OTP Lock Information[6]

OTPL	OTP Locked/Unlocked Status
0	OTP Block Unlock Status(Default)
1	OTP Block Lock Status(Disable OTP Program/Erase)

### Time Out (TO)

This bit determines if there is a time out for load, program, copy back program, and erase operations. It is fixed at 'no time out'.

## TO Information[0]

Item	Definition	Description
ТО	Time Out	0 = no time out

## **Controller Status Register Output Modes**

					Con	troller S	Status F	Register [15:0]				
Mode	OnGo	Lock	Load	Prog	Erase	Error	Sus	Reserved(0)	RSTB	OTPL	Reserved(0)	то
Load Ongoing	1	0	1	0	0	0	0	0	0	0/1	00000	0
Program Ongoing	1	0	0	1	0	0	0	0	0	0/1	00000	0
Erase Ongoing	1	0	0	0	1	0	0	0	0	0/1	00000	0
Reset Ongoing	1	0	0	0	0	0	0	0	1	0/1	00000	0
Multi-Block Erase Ongoing	1	0	0	0	1	0	0	0	0	0/1	00000	0
Erase Verify Read Ongoing	1	0	0	0	0	0	0	0	0	0/1	00000	0
Load OK	0	0	0	0	0	0	0	0	0	0/1	00000	0
Program OK	0	0	0	0	0	0	0	0	0	0/1	00000	0
Erase OK	0	0	0	0	0	0	0	0	0	0/1	00000	0
Erase Verify Read OK <sup>3)</sup>	0	0	0	0	0	0	0	0	0	0/1	00000	0
Load Fail <sup>1)</sup>	0	0	1	0	0	1	0	0	0	0/1	00000	0
Program Fail	0	0	0	1	0	1	0	0	0	0/1	00000	0
Erase Fail	0	0	0	0	1	1	0	0	0	0/1	00000	0
Erase Verify Read Fail <sup>3)</sup>	0	0	0	0	1	1	0	0	0	0/1	00000	0
Load Reset <sup>2)</sup>	0	0	1	0	0	1	0	0	1	0/1	00000	0
Program Reset	0	0	0	1	0	1	0	0	1	0/1	00000	0
Erase Reset	0	0	0	0	1	1	0	0	1	0/1	00000	0
Erase Suspend	0	0	0	0	1	0	1	0	0	0/1	00000	0
Program Lock	0	1	0	1	0	1	0	0	0	0/1	00000	0
Erase Lock	0	1	0	0	1	1	0	0	0	0/1	00000	0
Load Lock(Buffer Lock)	0	1	1	0	0	1	0	0	0	0/1	00000	0
OTP Program Fail(Lock)	0	1	0	1	0	1	0	0	0	1	00000	0
OTP Program Fail	0	0	0	1	0	1	0	0	0	0	00000	0
OTP Erase Fail	0	1	0	0	1	1	0	0	0	0/1	00000	0
Program Ongo- ing(Susp.)	1	0	0	1	1	0	1	0	0	0/1	00000	0
Load Ongoing(Susp.)	1	0	1	0	1	0	1	0	0	0/1	00000	0
Program Fail(Susp.)	0	0	0	1	1	1	1	0	0	0/1	00000	0
Load Fail(Susp.)	0	0	1	0	1	1	1	0	0	0/1	00000	0
Invalid Command	0	0	0	0	0	1	0	0	0	0/1	00000	0
Invalid Com- mand(Susp.)	0	0	0	0	1	1	1	0	0	0/1	00000	0

# NOTE:

- 1. ERm and/or ERs bits in ECC status register at Load Fail case is 10. (2bits error uncorrectable)
- 2. ERm and ERs bits in ECC status register at Load Reset case are 00. (No error)
- 3. Multi Block Erase status should be checked by Erase Verify Read operation.



# 2.8.22 Interrupt Status Register F241h (R/W)

This Read/Write register shows status of the OneNAND interrupts.

F241h, defaults = 8080h after Cold Reset; 8010h after Warm/Hot Reset

1	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	INT	Reserved(0000000)								WI	EI	RSTI		Reserve	ed(0000)	)

## Interrupt (INT)

This is the master interrupt bit. The INT bit is wired directly to the INT pin on the chip. Upon writing '0' to the INT bit, the INT pin goes low if INTpol is high and goes high if INTpol is low.

### INT Interrupt [15]

Status	Conditions	Defau	It State	Valid	Interrupt
Status	Conditions	Cold	Warm/hot	State	Function
		1	1	0	off
sets itself to '1'	One or more of RI, WI, RSTI and EI is set to '1', or 0065h, 0023h, 0071h, 002A and 002C commands are completed			0→1	Pending
clears to '0'	'0' is written to this bit, or Cold/Warm/Hot reset is being performed			1→0	off

## Read Interrupt (RI)

This is the Read interrupt bit.

## RI Interrupt [7]

Status	Conditions	Defau	It State	Valid	Interrupt
Status	Conditions	Cold	Warm/hot	State	Function
		1	0	0	off
sets itself to '1'	At the completion of an Load Operation (0000h, 0013h, Load Data into Buffer, or boot is done)			0→1	Pending
clears to '0'	'0' is written to this bit, or Cold/Warm/Hot reset is being performed			1→0	off

## Write Interrupt (WI)

This is the Write interrupt bit.

## WI Interrupt [6]

Status	Conditions	Defau	It State	Valid	Interrupt
Status	Conditions	Cold	Warm/hot	State	Function
		0	0	0	off
sets itself to '1'	At the completion of an Program Operation (0080h, 001Ah, 001Bh)			0→1	Pending
clears to '0'	'0' is written to this bit, or Cold/Warm/Hot reset is being performed			1→0	off



## Erase Interrupt (EI)

This is the Erase interrupt bit.

## El Interrupt [5]

Status	Conditions	Defau	It State	Valid	Interrupt
Status	Conditions	Cold	Warm/hot	State	Function
		0	0	0	off
sets itself to '1'	At the completion of an Erase Operation (0094h, 0095h, 0030h)			0→1	Pending
clears to '0'	'0' is written to this bit, or Cold/Warm/Hot reset is being performed			1→0	off

### Reset Interrupt (RSTI)

This is the Reset interrupt bit.

### RSTI Interrupt [4]

Status	Conditions	Defau	It State	Valid	Interrupt
Status	Conditions	Cold	Warm/hot	State	Function
		0	1	0	off
sets itself to '1'	At the completion of an Reset Operation (00B0h, 00F0h, 00F3h or warm reset is released)			0→1	Pending
clears to '0'	'0' is written to this bit			1→0	off

# 2.8.23 Start Block Address Register F24Ch (R/W)

This Read/Write register shows the NAND Flash block address in the Write Protection mode. Setting this register precedes a 'Lock Block' command, 'Unlock Block' command, or 'Lock-Tight' Command.

## F24Ch, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved(0000000)										SBA				

Device	Number of Block	SBA
512Mb	512	[8:0]

## SBA Information[9:0]

Item	Definition	Description
SBA	Start Block Address	Precedes Lock Block, Unlock Block, or Lock-Tight commands



# 2.8.24 End Block Address Register F24Dh

This register is reserved for future use.

# 2.8.25 NAND Flash Write Protection Status Register F24Eh (R)

This Read register shows the Write Protection Status of the NAND Flash memory array. To read the write protection status, FBA has to be set before reading the register.

### F24Eh, default = 0002h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved(000000000000)								US	LS	LTS				

#### Write Protection Status Information[2:0]

Item	Bit	Definition	Description
US	2	Unlocked Status	1 = current NAND Flash block is unlocked
LS	1	Locked Status	1 = current NAND Flash block is locked
LTS	0	Locked-Tight Status	1 = current NAND Flash block is locked-tight

# 2.8.26 ECC Status Register FF00h (R)

This Read register shows the Error Correction Status. The OneNAND can detect 1- or 2-bit errors and correct 1-bit errors. 3-bit or more error detection and correction is not supported.

ECC can be performed on the NAND Flash main and spare memory areas. The ECC status register can also show the number of errors in a sector as a result of an ECC check in during a load operation. ECC status bits are also updated during a boot loading operation.

### FF00h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
E	Rm3	EF	Rs3	ER	m2	EF	Rs2	ER	lm1	EF	Rs1	ER	m0	EF	Rs0

## **Error Status**

ERm, ERs	ECC Status
00	No Error
01	1-bit error(correctable)
10	2 bits error (uncorrectable)
11	Reserved



### ECC Information[15:0]

Item	Definition	Description
ERm0	1st selected sector of the main BufferRAM	Status of errors in the 1st selected sector of the main BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERm1	2nd selected sector of the main BufferRAM	Status of errors in the 2nd selected sector of the main BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERm2	3rd selected sector of the main BufferRAM	Status of errors in the 3rd selected sector of the main BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERm3	4th selected sector of the main BufferRAM	Status of errors in the 4th selected sector of the main BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERs0	1st selected sector of the spare BufferRAM	Status of errors in the 1st selected sector of the spare BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERs1	2nd selected sector of the spare BufferRAM	Status of errors in the 2nd selected sector of the spare BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERs2	3rd selected sector of the spare BufferRAM	Status of errors in the 3rd selected sector of the spare BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.
ERs3	4th selected sector of the spare BufferRAM	Status of errors in the 4th selected sector of the spare BufferRAM as a result of an ECC check during a load operation.  Also updated during a Bootload operation.

# 2.8.27 ECC Result of 1<sup>st</sup> Selected Sector, Main Area Data Register FF01h (R)

This Read register shows the Error Correction result for the 1st selected sector of the main area data. ECCposWord0 is the error position address in the Main Area data of 256 words. ECCposIO0 is the error position address which selects 1 of 16 DQs. ECCposWord0 and ECCposIO0 are also updated at boot loading.

### FF01h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserve	ed(0000)	)				ECCpo	sWord0					ECCp	ooslO0	

# 2.8.28 ECC Result of 1<sup>st</sup> Selected Sector, Spare Area Data Register FF02h (R)

This Read register shows the Error Correction result for the 1st selected sector of the spare area data. ECClogSector0 is the error position address for 1.5 words of 2nd and 3rd words in the spare area. ECCposIO0 is the error position address which selects 1 of 16 DQs. ECClogSector0 and ECCposIO0 are also updated at boot loading.

### FF02h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved(000000000)										Sector0		ECCp	ooslO0	



# 2.8.29 ECC Result of 2<sup>nd</sup> Selected Sector, Main Area Data Register FF03h (R)

This Read register shows the Error Correction result for the 2nd selected sector of the main area data. ECCposWord1 is the error position address in the Main Area data of 256 words. ECCposIO1 is the error position address which selects 1 of 16 DQs. ECCposWord1 and ECCposIO1 are also updated at boot loading.

### FF03h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserve	d(0000)		ECCposWord1 ECCposIO1									oslO1		

# 2.8.30 ECC Result of 2<sup>nd</sup> Selected Sector, Spare Area Data Register FF04h (R)

This Read register shows the Error Correction result for the 2nd selected sector of the spare area data. ECClogSector1 is the error position address for 1.5 words of 2nd and 3rd words in the spare area. ECCposIO1 is the error position address which selects 1 of 16 DQs. ECClogSector1 and ECCposIO1 are also updated at boot loading.

### FF04h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved(000000000)									ECClog			ECCp	oslO1	

# 2.8.31 ECC Result of 3<sup>rd</sup> Selected Sector, Main Area Data Register FF05h (R)

This Read register shows the Error Correction result for the 3rd selected sector of the main area data. ECCposWord2 is the error position address in the Main Area data of 256 words. ECCposIO2 is the error position address which selects 1 of 16 DQs. ECCposWord2 and ECCposIO2 are also updated at boot loading.

## FF05h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserve	ed(0000)	)	ECCposWord2 ECCposIO2											

# 2.8.32 ECC Result of 3<sup>rd</sup> Selected Sector, Spare Area Data Register FF06h (R)

This Read register shows the Error Correction result for the 3rd selected sector of the spare area data. ECClogSector2 is the error position address for 1.5 words of 2nd and 3rd words in the spare area. ECCposIO2 is the error position address which selects 1 of 16 DQs. ECClogSector2 and ECCposIO2 are also updated at boot loading.

### FF06h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Res	erved(0	000000	000)				ECClog	Sector2		ECCp	oslO2	



# 2.8.33 ECC Result of 4<sup>th</sup> Selected Sector, Main Area Data Register FF07h (R)

This Read register shows the Error Correction result for the 4th selected sector of the main area data. ECCposWord3 is the error position address in the Main Area data of 256 words. ECCposIO3 is the error position address which selects 1 of 16 DQs. ECCposWord3 and ECCposIO3 are also updated at boot loading.

### FF07h, default = 0000h

ĺ	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ĺ		Reserve	d(0000)		ECCposWord3 ECCposle									oslO3		

# 2.8.34 ECC Result of 4<sup>th</sup> Selected Sector, Spare Area Data Register FF08h (R)

This Read register shows the Error Correction result for the 4th selected sector of the spare area data. ECClogSector3 is the error position address for 1.5 words of 2nd and 3rd words in the spare area. ECCposIO3 is the error position address which selects 1 of 16 DQs. ECClogSector3 and ECCposIO3 are also updated at boot loading.

#### FF08h, default = 0000h

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved(000000000)									ECClog	Sector3		ECCp	osIO3	

## **ECC Log Sector**

ECClogSector0~ECClogSector3 indicates the error position in the 2nd word and LSB of 3rd word in the spare area. Refer to note 2 in chapter 2.7.2

### **ECClogSector Information [5:4]**

ECClogSector	Error Position
00	2nd word
01	3rd word
10, 11	Reserved



# 3.0 DEVICE OPERATION

This section of the datasheet discusses the operation of the OneNAND device. It is followed by AC/DC Characteristics and Timing Diagrams which may be consulted for further information.

The OneNAND supports either a command-based or a register-based interface for performing operations on the device including reading device ID, writing data to buffer etc.

The command-based interface is active in the boot partition. Commands can only be written with a boot area address. Boot area data is only returned if no command has been issued prior to the read.

When reading and programming the OneNAND, data handling is data unit-based. One data is read at the time from NAND Flash array.

# 3.1 Command Based Operation

The command-based interface is active in the boot partition. Commands can only be written with a boot area address. Boot area data is only returned if no command has been issued prior to the read.

The entire address range, except for the boot area, can be used for the data buffer. All commands are written to the boot partition. Writes outside the boot partition are treated as normal writes to the buffers or registers.

The command consists of one or more cycles depending on the command. After completion of the command the device starts its execution. Writing incorrect information including address and data to the boot partition or writing an improper command will terminate the previous command sequence and make the device enter the ready status.

The defined valid command sequences are stated in Command Sequences Table.

### **Command Sequences**

Command Definition		Cycles	1st cycle	2nd cycle
Read Data from Buffer	Add	1	DP <sup>1)</sup>	
Read Data ITOITI Bullet	Data	'	Data	
Write Data to Buffer	A/rite Date to Duffer Add		DP	
White Data to Bullet	Data	ı	Data	
Reset OneNAND	Add	1	BP <sup>2)</sup>	
	Data	'	00F0h	
Load Data into Buffer <sup>3)</sup>	Add	2	BP	BP
Load Data Into Bullet	Data	2	00E0h	0000h <sup>4)</sup>
Read Identification Data <sup>6)</sup>	Add	2	BP	XXXXh <sup>5)</sup>
Read Identification Data 97	Data	2	0090h	Data

#### NOTE

- 1) DP(Data Partition) : DataRAM Area
- 2) BP(Boot Partition) : BootRAM Area [0000h ~ 01FFh, 8000h ~ 800Fh]
- 3) Load Data into Buffer operation is available within a block(128KB)
- 4) Load 2KB unit into DataRAM0. Current Start address(FPA) is automatically incremented by 2KB unit after the load.
- 5) 0000h -> Data is Manufacturer ID
- 0001h -> Data is Device ID
- 0002h -> Current Block Write Protection Status
- 6) WE toggling can terminate 'Read Identification Data' operation.



# 3.1.1 Reading Data From Buffer

The buffer memory can be read by addressing a Read to the desired buffer area.

# 3.1.2 Writing Data to Buffer

The buffer memory can be written to by addressing a Write to a desired buffer area.

# 3.1.3 Reset OneNAND Command

The Reset command is given by writing 00F0h to the boot partition address. Reset will return all default values into the device.

## 3.1.4 Load Data Into Buffer Command

Load Data into Buffer command is a two-cycle command. Two sequential designated command activates this operation. Sequentially writing 00E0h and 0000h to the boot partition [0000h~01FFh, 8000h~800Fh] will load one page to DataRAM0. This operation refers to FBA and FPA. FSA, BSA, and BSC are not considered.

At the end of this operation, FPA will be automatically increased by 1. So continuous issue of this command will sequentially load data in next page to DataRAM0. This page address increment is restricted within a block.

The default value of FBA and FPA is 0. Therefore, initial issue of this command after power on will load the first page of memory, which is usually boot code.

# 3.1.5 Read Identification Data Command

The Read Identification Data command consists of two cycles. It gives out the devices identification data according to the given address. The first cycle is 0090h to the boot partition address and second cycle is read from the addresses specified in Identification Data Description Table.

## **Identification Data Description**

Address	Data Out	
0000h	Manufacturer ID : 00ECh	
0001h	Device ID : refer to chapter 2.8.3	
0002h	Current Block Write Protection Status 1)	

Note 1) To read the write protection status, FBA has to be set before issuing this command.



# 3.2 Device Bus Operation

The device bus operations are shown in the table below.

Operation	CE	OE	WE	ADD0~15	DQ0~15	RP	CLK	AVD
Standby	Н	Х	X	Х	High-Z	Н	Х	Х
Warm Reset	X	Х	Х	Х	High-Z	L	Х	Х
Asynchronous Write	L	Н	L	Add. In	Data In	Н	L	Х
Asynchronous Read	L	L	Н	Add. In	Data Out	Н	L	or L
Load Initial Burst Address	L	Н	Н	Add. In	X	Н		
Burst Read	L	L	Н	х	Burst Data Out	Н		х
Terminate Burst Read Cycle	Н	х	Н	Х	High-Z	Н	х	Х
Terminate Burst Read Cycle via RP	Х	Х	Х	Х	High-Z	L	Х	Х
Terminate Current Burst Read Cycle and Start New Burst Read Cycle		Н	Н	Add In	High-Z	Н		

Note : L=V<sub>IL</sub> (Low), H=V<sub>IH</sub> (High), X=Don't Care.



# 3.3 Reset Mode Operation

The One NAND has 4 reset modes: Cold/Warm/Hot Reset, and NAND Flash Array Reset. Section 3.3 discusses the operation of these reset modes.

The Register Reset Table shows the which registers are affected by the various types or Reset operations.

### **Internal Register Reset Table**

	Internal Registers	Default	Cold Reset	Warm Reset (RP)	Hot Reset (00F3h)	Hot Reset (BP-F0)	NAND Flash Reset(00F0h)
F000h	Manufacturer ID Register (R)	00ECh	N/A	N/A	N/A		N/A
F001h	Device ID Register (R): OneNAND	(Note 3)	N/A	N/A	N/	/A	N/A
F002h	Version ID Register (R): N/A	N/A	N/A	N/A	N/	/A	N/A
F003h	Data Buffer size Register (R)	0800h	N/A	N/A	N/	/A	N/A
F004h	Boot Buffer size Register (R)	0200h	N/A	N/A	N/	/A	N/A
F005h	Amount of Buffers Register (R)	0201h	N/A	N/A	N/	/A	N/A
F006h	Technology Register (R)	0000h	N/A	N/A	N/	/A	N/A
F100h	Start Address1 Register (R/W): FBA	0000h	0000h	0000h	000	00h	N/A
F101h	Start Address2 Register (R/W): Reserved	0000h	0000h	0000h	000	00h	N/A
F102h	Start Address3 Register (R/W): FCBA	0000h	0000h	0000h	000	00h	N/A
F103h	Start Address4 Register (R/W): FCPA, FCSA	0000h	0000h	0000h	0000h		N/A
F107h	Start Address8 Register (R/W): FPA, FSA	0000h	0000h	0000h	0000h		N/A
F200h	Start Buffer Register (R/W): BSA, BSC	0000h	0000h	0000h	000	00h	N/A
F220h	Command Register (R/W)	0000h	0000h	0000h	000	00h	N/A
F221h	System Configuration 1 Register (R/W)	40C0h	40C0h	O (Note1)	O (N	ote1)	N/A
F240h	Controller Status Register (R)	0000h	0000h	0000h	000	00h	N/A
F241h	Interrupt Status Register (R/W)	-	8080h	8010h	801	l0h	N/A
F24Ch	Start Block Address (R/W)	0000h	0000h	0000h	N/	/A	N/A
F24Dh	End Block Address: N/A	N/A	N/A	N/A	N/	/A	N/A
F24Eh	NAND Flash Write Protection Status (R)	0002h	0002h	0002h	N/	/A	N/A
FF00h	ECC Status Register (R) (Note2)	0000h	0000h	0000h	000	)0h	N/A
FF01h	ECC Result of Sector 0 Main area data Register(R)	0000h	0000h	0000h	000	)0h	N/A
FF02h	ECC Result of Sector 0 Spare area data Register (R)	0000h	0000h	0000h	000	00h	N/A
FF03h	ECC Result of Sector 1 Main area data Register(R)	0000h	0000h	0000h	000	00h	N/A
FF04h	ECC Result of Sector 1 Spare area data Register (R)	0000h	0000h	0000h	000	)0h	N/A
FF05h	ECC Result of Sector 2 Main area data Register(R)	0000h	0000h	0000h	000	)0h	N/A
FF06h	ECC Result of Sector 2 Spare area data Register (R)	0000h	0000h	0000h	000	)0h	N/A
FF07h	ECC Result of Sector 3 Main area data Register(R)	0000h	0000h	0000h	000	)0h	N/A
FF08h	ECC Result of Sector 3 Spare area data Register (R)	0000h	0000h	0000h	000	)0h	N/A

NOTE: 1) RDYpol, INTpol, IOBE is reset by Cold reset. The other bits are reset by Cold/Warm/Hot reset.



<sup>2)</sup> ECC Status Register & ECC Result Registers are reset when any command is issued.

<sup>3)</sup> Refer to Device ID Register F001h.

# 3.3.1 Cold Reset Mode Operation

## See Timing Diagram 6.11

At system power-up, the voltage detector in the device detects the rising edge of Vcc and releases an internal power-up reset signal. This triggers bootcode loading. Bootcode loading means that the boot loader in the device copies designated sized data (1KB) from the beginning of memory into the BootRAM. This sequence is the Cold Reset of OneNAND.

The POR(Power On Reset) triggering level is typically 1.5V. Boot code copy operation activates 400us after POR. Therefore, the system power should reach 1.7V within 400us from the POR triggering level for bootcode data to be valid. For 2.65V and 3.3V device, POR level is typically 1.8V and system power should reach 2.2V within 400us.

It takes approximately 70us to copy 1KB of bootcode. Upon completion of loading into the BootRAM, it is available to be read by the host. The INT pin is not available until after IOBE = 1 and IOBE bit can be changed by host.

# 3.3.2 Warm Reset Mode Operation

## See Timing Diagrams 6.12

A Warm Reset means that the host resets the device by using the /RP pin. When the a /RP low is issued, the device logic stops all current operations and executes internal reset operation and resets current NAND Flash core operation synchronized with the falling edge of /RP.

During an Internal Reset Operation, the device initializes internal registers and makes output signals go to default status. The BufferRAM data is kept unchanged after Warm/Hot reset operations.

The device guarantees the logic reset operation in case /RP pulse is longer than tRP min. (200ns). The device may reset if tRP < tRP min(200ns), but this is not guaranteed.

Warm reset will abort the current NAND Flash core operation. During a warm reset, the content of memory cells being altered is no longer valid as the data will be partially programmed or erased.

Warm reset has no effect on contents of BootRAM and DataRAM.

# 3.3.3 Hot Reset Mode Operation

## See Timing Diagram 6.13

A Hot Reset means that the host resets the device by Reset command. The reset command can be either Command based or Register Based. Upon receiving the Reset command, the device logic stops all current operation and executes an internal reset operation and resets the current NAND Flash core operation.

During an Internal Reset Operation, the device initializes internal registers and makes output signals go to default status. The BufferRAM data is kept unchanged after Warm/Hot reset operations.

Hot reset will abort the current NAND Flash core operation. During a Hot reset, the content of memory cells being altered is no longer valid as the data will be partially programmed or erased.

Hot reset has no effect on contents of BootRAM and DataRAM.

# 3.3.4 NAND Flash Core Reset Mode Operation

### See Timing Diagram 6.14

The Host can reset the NAND Flash Core operation by issuing a NAND Flash Core reset command. NAND Flash core reset will abort the current NAND Flash core operation. During a NAND Flash core reset, the content of memory cells being altered is no longer valid as the data will be partially programmed or erased.

Hot reset has no effect on contents of BootRAM and DataRAM, as well as register values.



# 3.4 Write Protection Operation

The OneNAND can be write-protected to prevent re-programming or erasure of data. The areas of write-protection are the BootRAM, and the NAND Flash Array.

# 3.4.1 BootRAM Write Protection Operation

At system power-up, voltage detector in the device detects the rising edge of Vcc and releases the internal power-up reset signal which triggers bootcode loading. And the designated size data(1KB) is copied from the first page of the first block in the NAND flash array to the BootRAM.

After the bootcode loading is completed, the BootRAM is always locked to protect the boot code from the accidental write.

# 3.4.2 NAND Flash Array Write Protection Operation

The device has both hardware and software write protection of the NAND Flash array.

### **Hardware Write Protection Operation**

The hardware write protection operation is implemented by executing a Cold or Warm Reset. On power up, the NAND Flash Array is in its default, locked state. The entire NAND Flash array goes to a locked state after a Cold or Warm Reset.

#### **Software Write Protection Operation**

The software write protection operation is implemented by writing a Lock command (002Ah) or a Lock-tight command (002Ch) to command register (F220h).

Lock (002Ah) and Lock-tight (002Ch) commands write protects the block defined in the Start Block Address Register F24Ch.

# 3.4.3 NAND Array Write Protection States

There are three lock states in the NAND Array: unlocked, locked, and locked-tight.

OneNAND512 supports lock/unlock/lock-tight by **one block**, so each block should be locked/unlocked/locked-tight individually.

### Write Protection Status

The current block Write Protection status can be read in NAND Flash Write Protection Status Register(F24Eh). There are three bits - US, LS, LTS -, which are not cleared by hot reset. These Write Protection status registers are updated when FBA is set, and when Write Protection command is entered.

The followings summarize locking status.

example)

In default, [2:0] values are 010.

- -> If host executes unlock block operation, then [2:0] values turn to 100.
- -> If host executes lock-tight block operation, then [2:0] values turn to 001.



# 3.4.3.1 Unlocked NAND Array Write Protection State

An Unlocked block can be programmed or erased. The status of an unlocked block can be changed to locked or locked-tight using the appropriate software command. (locked-tight state can be achieved via lock-tight command which follows lock command)

Only one block can be released from lock state to unlock state with Unlock command and addresses. The unlocked block can be changed with new lock command. Therefore, each block has its own lock/unlock/lock-tight state.

Unlocked	
	Unlock Command Sequence: Start block address+Unlock block command (0023h)
	Guit block ddaress - Gillock block communa (60261)

# 3.4.3.2 Locked NAND Array Write Protection State

A Locked block cannot be programmed or erased. All blocks default to a locked state following a Cold or Warm Reset. Unlocked blocks can be changed to locked using the Lock block command. The status of a locked block can be changed to unlocked or locked-tight using the appropriate software command.

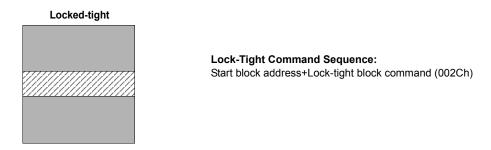
Locked	
	Lock Command Sequence: Start block address+Lock block command (002Ah)



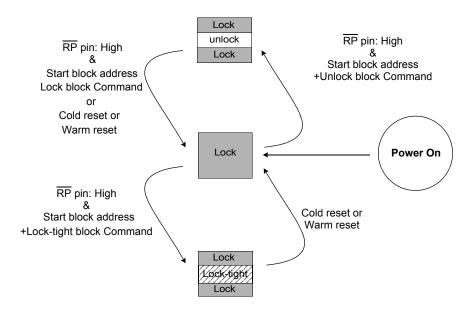
# 3.4.3.3 Locked-tight NAND Array Write Protection State

A block that is in a locked-tight state can only be changed to lock state after a Cold or Warm Reset. Unlock and Lock command sequences will not affect its state. This is an added level of write protection security.

A block must first be set to a locked state before it can be changed to locked-tight using the Lock-tight command. locked-tight blocks will revert to a locked state following a Cold or Warm Reset.

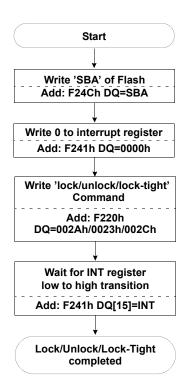


# 3.4.4 NAND Flash Array Write Protection State Diagram





**Data Protection Operation Flow Diagram** 



Note) Samsung strongly recommend to follow the above flow chart

# 3.5 Data Protection During Power Down Operation

See Timing Diagram 6.15

The device is designed to offer protection from any involuntary program/erase during power-transitions.

An internal voltage detector disables all functions whenever Vcc is below POR level, about 1.3V. It is recommended that the /RP pin, which provides hardware protection, should be kept at VIL before power-down.

# 3.6 Load Operation

### See Timing Diagram 6.8

The Load operation is initiated by setting up the start address from which the data is to be loaded. The Load command is issued in order to initiate the load.

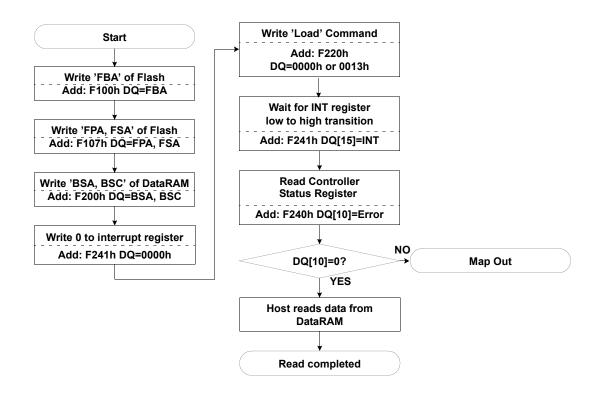
During a Load operation, the device:

- -Transfers the data from NAND Flash array into the BufferRAM
- -ECC is checked and any detected and corrected error is reported in the status response as well as any unrecoverable error.

Once the BufferRAM has been filled, an interrupt is issued to the host so that the contents of the BufferRAM can be read. The read from the BufferRAM can be an asynchronous read mode or synchronous read mode. The status information related to load operation can be checked by the host if required.

The device has a dual data buffer memory architecture (DataRAM0, DataRAM1), each 2KB in size. Each DataRAM buffer has 4 Sectors. The device is capable of independent and simultaneous data-read operation from one data buffer and data-load operation to the other data buffer. Refer to the information for more details in section 3.12.1, "Read-While-Load Operation".

## **Load Operation Flow Chart Diagram**





# 3.7 Read Operation

See Timing Diagrams 6.1, 6.2, 6.3, 6.4, 6.5 and 6.6

The device has two read modes; Asynchronous Read and Synchronous Burst Read.

The initial state machine automatically sets the device into the Asynchronous Read Mode (RM=0) to prevent the spurious altering of memory content upon device power up or after a Hardware reset. No commands are required to retrieve data in Asynchronous Read Mode.

The Synchronous Read Mode is enabled by setting RM bit of System Configuration1 Register (F221h) to Synchronous Read Mode (RM=1). See Section 2.8.19 for more information about System Configuration1 Register.

# 3.7.1 Asynchronous Read Mode Operation (RM=0)

See Timing Diagrams 6.3, 6.4, 6.5 and 6.6

In an Asynchronous Read Mode, data is output with respect to a logic input, /AVD.

Output data will appear on DQ15-DQ0 in when a valid address is asserted on A15-A0 while driving /AVD and /CE to VIL. / WE is held at VIH. The function of the /AVD signal is to latch the valid address.

Address access time from /AVD low (tAA) is equal to the delay from valid addresses to valid output data.

The Chip Enable access time (tCE) is equal to the delay from the falling edge of /CE to valid data at the outputs.

The Output Enable access time (tOE) is the delay from the falling edge of /OE to valid data at the output.

# 3.7.2 Synchronous Read Mode Operation (RM=1)

See Timing Diagrams 6.1 and 6.2

In a Synchronous Read Mode, data is output with respect to a clock input.

The device is capable of a continuous linear burst operation and a fixed-length linear burst operation of a preset length. Burst address sequences for continuous and fixed-length burst operations are shown in the table below.

## **Burst Address Sequences**

	Start	Burst Address Sequence(Decimal)				
	Addr.	Continuous Burst	4-word Burst	8-word Burst	16-word Burst	32-word Burst
	0	0-1-2-3-4-5-6	0-1-2-3-0	0-1-2-3-4-5-6-7-0	0-1-2-3-413-14-15-0	0-1-2-3-429-30-31-0
Wrap	1	1-2-3-4-5-6-7	1-2-3-0-1	1-2-3-4-5-6-7-0-1	1-2-3-4-514-15-0-1	1-2-3-4-530-31-0-1
around	2	2-3-4-5-6-7-8	2-3-0-1-2	2-3-4-5-6-7-0-1-2	2-3-4-5-615-0-1-2	2-3-4-5-631-0-1-2
		•	•	•		

In the burst mode, the initial word will be output asynchronously, regardless of BRL. While the following words will be determined by BRL value.

The latency is determined by the host based on the BRL bit setting in the System Configuration 1 Register. The default BRL is 4 latency cycles. At clock frequencies of 40MHz or lower, latency cycles can be reduced to 3. BRL can be set up to 7 latency cycles.

The BRL registers can be read during a burst read mode by using the /AVD signal with an address.



# 3.7.2.1 Continuous Linear Burst Read Operation

See Timing Diagram 6.2

## First Clock Cycle

The initial word is output at tIAA after the rising edge of the first CLK cycle. The RDY output indicates the initial word is ready to the system by pulsing high. If the device is accessed synchronously while it is set to Asynchronous Read Mode, the first data can still be read out.

### **Subsequent Clock Cycles**

Subsequent words are output (Burst Access Time from Valid Clock to Output) tBA after the rising edge of each successive clock cycle, which automatically increments the internal address counter.

### **Terminating Burst Read**

The device will continue to output sequential burst data until the system asserts /CE high, or /RP low, wrapping around until it reaches the designated address (see Section 2.7.3 for address map information). Alternately, a Cold/Warm/Hot Reset, asserting /CE high, or a /WE low pulse will terminate the burst read operation.

### Synchronous Read Boundary

Division	Add.map(word order)	
BootRAM Main(0.5Kw)	0000h~01FFh	Not Support
BrfferRAM0 Main(1Kw)	0200h~05FFh	
BufferRAM1 Main(1Kw)	0600h~09FFh	Not Support
Reserved Main	0A00h~7FFFh	
BootRAM Spare(16w)	8000H~800Fh	Not Support
BufferRAM0 Spare(32w)	8010h~802Fh	Not Support
BufferRAM1 Spare(32w)	8030h~804Fh	Not Support
Reserved Spare	8050h~8FFFh	
Reserved Register	9000h~EFFFh	
Register(4Kw)	F000h~FFFFh	

<sup>\*</sup> Reserved area is not available on Synchronous read

# 3.7.2.2 4-, 8-, 16-, 32-Word Linear Burst Read Operation

See Timing Diagram 6.1

An alternate Burst Read Mode enables a fixed number of words to be read from consecutive address.

The device supports a burst read from consecutive addresses of 4-, 8-, 16-, and 32-words with a linear-wrap around. When the last word in the burst has been reached, /CE and /OE are asserted high.

In this mode, the start address for the burst read can be any address of the address map with one exception. The device does not support a 32-word linear burst read on the spare area of the BufferRAM.



# 3.7.2.3 Programmable Burst Read Latency Operation

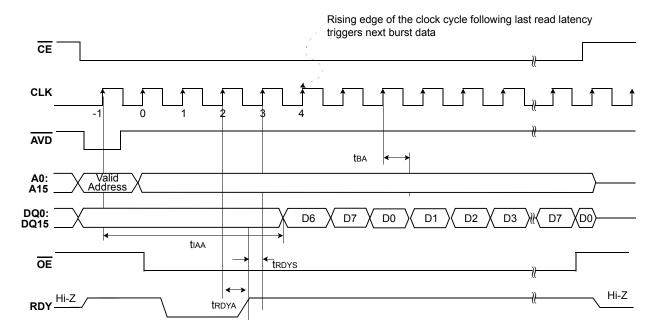
See Timing Diagrams 6.1 and 6.2

Upon power up, the number of initial clock cycles from Valid Address (/AVD) to initial data defaults to four clocks.

The number of clock cycles (n) which are inserted after the clock which is latching the address. The host can read the first data with the (n+1)th rising edge.

The number of total initial access cycles is programmable from three to seven cycles. After the number of programmed burst clock cycles is reached, the rising edge of the next clock cycle triggers the next burst data.

### Four Clock Burst Read Latency (default condition)



# 3.7.3 Handshaking Operation

The handshaking feature allows the host system to simply monitor the RDY signal from the device to determine when the initial word of burst data is ready to be read.

To set the number of initial cycles for optimal burst mode, the host should use the programmable burst read latency configuration (see Section 2.8.19, "System Configuration1 Register").

The rising edge of RDY which is derived from one cycles ahead of data fetch clock indicates the initial word of valid burst data.



# 3.7.4 Output Disable Mode Operation

When the CE or OE input is at  $V_{\text{IH}}$ , output from the device is disabled. The outputs are placed in the high impedance state.

# 3.8 Program Operation

### See Timing Diagram 6.9

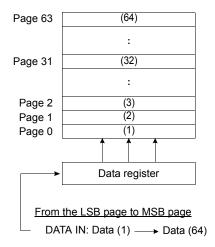
The Program operation is used to program data from the on-chip BufferRAMs into the NAND FLASH memory array.

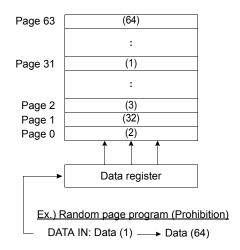
The device has two 2KB data buffers, each 1 Page (2KB + 64B) in size. Each page has 4 sectors of 512B each main area and 16B spare area. The device can be programmed in units of 1~4 sectors.

The architecture of the DataRAMs permits a simultaneous data-write operation from the Host to one of data buffers and a program operation from the other data buffer to the NAND Flash Array memory. Refer to Section 3.12.2, "Write While Program Operation", for more information.

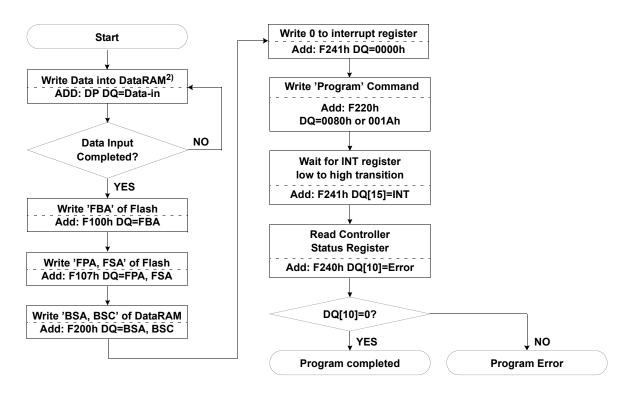
### Addressing for program operation

Within a block, the pages must be programmed consecutively from the LSB (least significant bit) page of the block to MSB (most significant bit) pages of the block. Random page address programming is prohibited.





### **Program Operation Flow Diagram**



\* : If program operation results in an error, map out the block including the page in error and copy the target data to another block.

Note 1) This must happen before data input

2) Data input could be done anywhere between "Start" and "Write Program Command".

During the execution of the Internal Program Routine, the host is not required to provide any further controls or timings. Furthermore, all commands, except a Reset command, will be ignored. A reset during a program operation will cause data corruption at the corresponding location.

If a program error is detected at the completion of the Internal Program Routine, map out the block, including the page in error, and copy the target data to another block. An error is signaled if DQ10 = "1" of Controller Status Register(F240h).

Data input from the Host to the DataRAM can be done at any time during the Internal Program Routine after "Start" but before the "Write Program Command" is written.



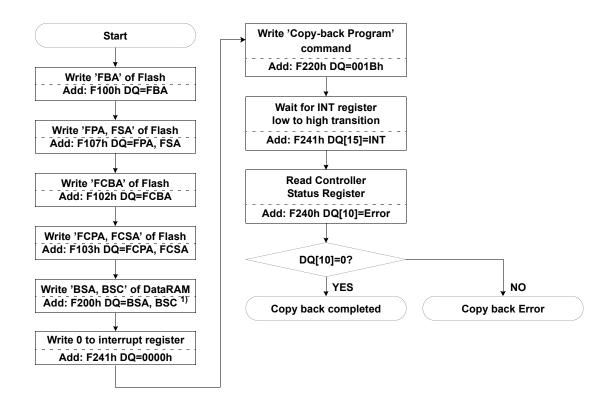
# 3.9 Copy-Back Program Operation

The Copy-Back program is configured to quickly rewrite data stored in one page without utilizing memory other than OneNAND. Since the time-consuming cycles of serial access and re-loading cycles are removed, the system performance is improved. The benefit is especially obvious when a portion of block is updated and the rest of the block also need to be copied to the newly assigned free block.

Data from the source page is saved in one of the on-chip DataRAM buffers and then programmed directly into the destination page. The DataRAM overwrites the previous data using the Buffer Sector Address (BSA) and Buffer Sector Count (BSC).

The Copy-Back Program Operation does this by performing sequential page-reads without a serial access and executing a copy-program using the address of the destination page.

### **Copy-Back Program Operation Flow Chart**



\* : If program operation results in an error, map out the block including the page in error and copy the target data to another block.

Note 1) Selected DataRAM by BSA & BSC is used for Copy back operation, so previous data is overwritten.

2) FBA, FPA and FSA should be input prior to FCBA, FCPA and FCSA.



The Copy-Back steps shown in the flow chart are:

- Data is read from the NAND Array using Flash Block Address (FBA), Flash Page Address (FPA) and Flash Sector Address (FSA). FBA, FPA, and FSA identify the source address to read data from NAND Flash array.
- The BufferRAM Sector Count (BSC) and BufferRAM Sector Address (BSA) identifies how many sectors and the location of the sectors in DataRAM that are used.
- The destination address in the NAND Array is written using the Flash Copy-Back Block Address (FCBA), Flash Copy-Back Page Address (FCPA), and Flash Copy-Back Sector Address (FCSA).
- The Copy-Back Program command is issued to start programming.
- Upon completion of copy-back programming to the destination page address, the Host checks the status to see if the operation was successfully completed. If there was an error, map out the block including the page in error and copy the target data to another block.



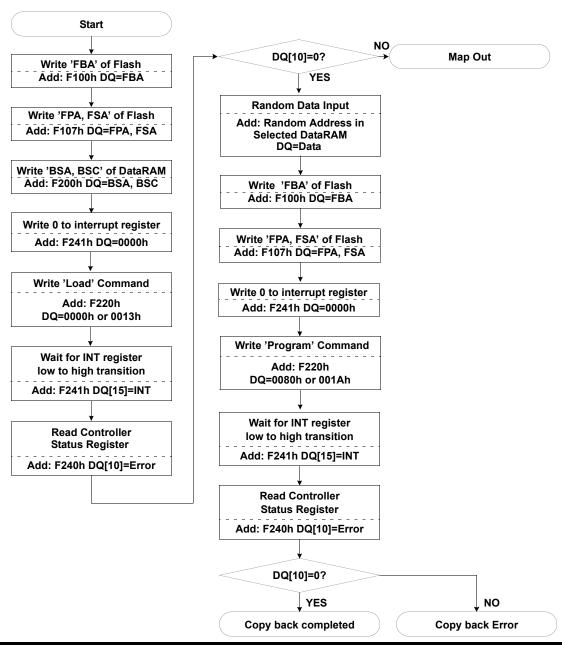
# 3.9.1 Copy-Back Program Operation with Random Data Input

The Copy-Back Program Operation with Random Data Input in OneNAND consists of 2 phase, Load data into DataRAM, Modify data and program into designated page. Data from the source page is saved in one of the on-chip DataRAM buffers and modified by the host, then programmed into the destination page.

As shown in the flow chart, data modification is possible upon completion of load operation. ECC is also available at the end of load operation. Therefore, using hardware ECC of OneNAND, accumulation of 1 bit error can be avoided.

Copy-Back Program Operation with Random Data Input will be effectively utilized at modifying certain bit, byte, word, or sector of source page to destination page while it is being copied.

## Copy-Back Program Operation with Random Data Input Flow Chart





# 3.10 Erase Operation

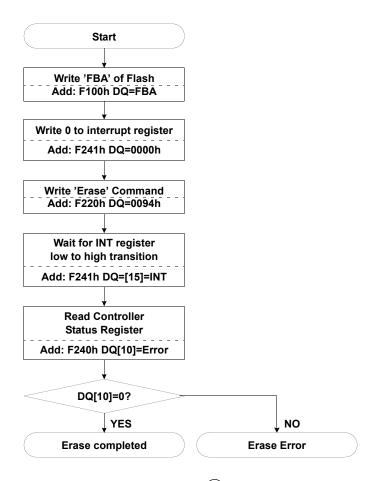
There are multiple methods for erasing data in the device including Block Erase and Multi-Block Erase.

# 3.10.1 Block Erase Operation

See Timing Diagram 6.10

The device can be erased one block at a time. To erase a block is to write all 1's into the desired memory block by executing the Internal Erase Routine. All previous data is lost.

## **Block Erase Operation Flow Chart**



(\*): If erase operation results in an error, map out the failing block and replace it with another block.

In order to perform the Internal Erase Routine, the following command sequence is necessary.

- The Host sets the block address of the memory location.
- The Erase Command initiates the Internal Erase Routine. During the execution of the Routine, the host is not required to provide further controls or timings. During the Internal erase routine, all commands, except the Reset command and Erase Suspend Command, written to the device will be ignored.

A reset during an erase operation will cause data corruption at the corresponding location.

# 3.10.2 Multi-Block Erase Operation

See Timing Diagram 6.10

Using Multi-Block Erase, the device can erase up to 64 multiple blocks simultaneously.

Multiple blocks can be erased by issuing a Multi-Block Erase command and writing the block address of the memory location to be erased. The final Flash Block Address (FBA) and Block Erase command initiate the internal multi block erase routine. During a Multi-Block Erase, the OnGo bit of the Controller Status Register is set to '1'(busy) from the time first block address to be latched is written until the actual erase has finished.

During block address latch sequence, issuing of other commands except Block Erase, Multi Block Erase, and Erase suspend at INT=High will abort the current operation. So to speak, It will cancel the previously latched addresses of Multi Block Erase Operation. On the other hand, Other command issue at INT=low will be ignored.

A reset during an erase operation will cause data corruption at the address location being operated on during the reset.

Despite a failed block during Multi-Block Erase operation, the device will continue the erase operation until all other specified blocks are erased.

Erase Suspend Command issue during Multi Block Erase Address latch sequence is prohibited.

### **Locked Blocks**

If there are locked blocks in the specified range, the Multi-Block Erase operation works as the follows.

Case 1: All specified blocks except BA(2) will be erased.

[BA(1)+0095h] + [BA((2), locked))+0095h] + ... + [BA(N-1)+0095h] + [BA(N)+0094h]

Case 2: Multi-Block Erase Operation is suspended and fails to start if the last Block Erase command is put together with the locked block address until right command and address input are issued.

[BA(1)+0095h] + [BA(2)+0095h] + ... + [BA(N-1)+0095h] + [BA((N), locked))+0094h]

Case 3: All specified blocks except BA(N) are erased.

[BA(1)+0095h] + [BA(2)+0095h] + ... + [BA(N-1)+0095h] + [BA((N, locked))+0094h] + [BA(N+1)+0094h]

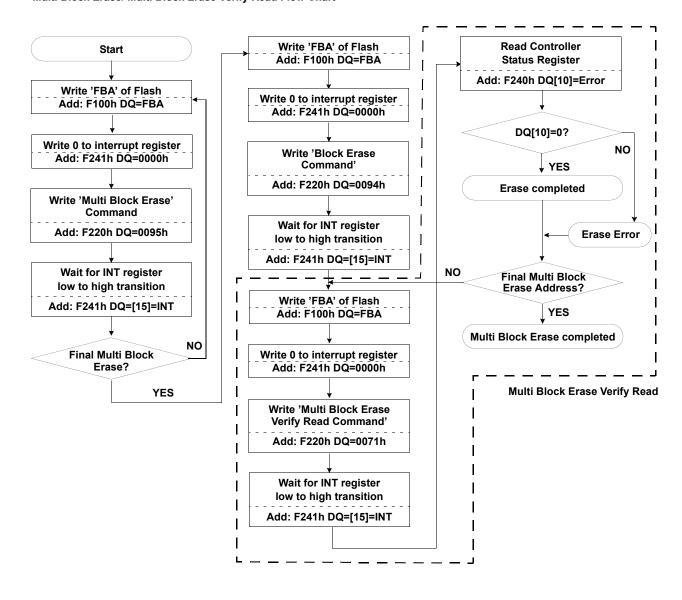


# 3.10.3 Multi-Block Erase Verify Read Operation

After a Multi-Block Erase Operation, verify Erase Operation result of each block with Multi-Block Erase Verify Command combined with address of each block.

If a failed address is identified, it must be managed in firmware.

## Multi Block Erase/ Multi Block Erase Verify Read Flow Chart





## 3.10.4 Erase Suspend / Erase Resume Operation

The Erase Suspend/Erase Resume Commands interrupt and restart a Block Erase or Multi-Block Erase operation so that user may perform another urgent operation on the block that is not being designated by Erase/Multi-Block Erase Operation.

### **Erase Suspend During a Block Erase Operation**

When Erase Suspend command is written during a Block Erase or Multi-Block Erase operation, the device requires a maximum of 500us to suspend erase operation. Erase Suspend Command issue during Block Address latch sequence is prohibited.

After the erase operation has been suspended, the device is ready for the next operation including a load, program, copy-back program, Lock, Unlock, Lock-tight, Hot Reset, NAND Flash Core Reset, Command Based Reset, Multi-Block Erase Read Verity, or OTP Access.

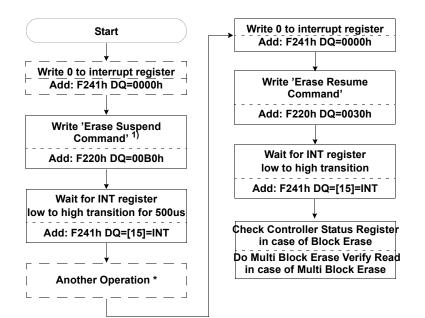
The subsequent operation can be to any block that was NOT being erased.

A special case arises pertaining Erase Suspend to the OTP. A Reset command is used to exit from the OTP Access mode. If the Reset-triggered exit from the OTP Access Mode happens during an Erase Suspend Operation, the erase routine could fail. Therefore to exit from the OTP Access Mode without suspending the erase operation stop, a 'NAND Flash Core Reset' command should be issued.

For the duration of the Erase Suspend period the following commands are not accepted:

• Block Erase/Multi-Block Erase/Erase Suspend

### **Erase Suspend and Erase Resume Operation Flow Chart**



\* Another Operation; Load, Program Copy-back Program, OTP Access<sup>2)</sup>, Hot Reset, Flash Reset, CMD Reset, Multi Block Erase Verify, Lock, Lock-tight, Unlock

Note 1) Erase Suspend command input is prohibited during Multi Block Erase address latch period.

2) If OTP access mode exit happens with Reset operation during Erase Suspend mode, Reset operation could hurt the erase operation. So if a user wants to exit from OTP access mode without the erase operation stop, Reset NAND Flash Core command should be used.



### **Erase Resume**

When the Erase Resume command is executed, the Block Erase will restart. The Erase Resume operation does not actually resume the erase, but starts it again from the beginning.

When an Erase Suspend or Erase Resume command is executed, the addresses are in Don't Care state.

# 3.11 OTP Operation

On Block of the NAND Flash Array memory is reserved as a One-Time Programmable Block memory area.

The OTP block can be read, programmed and locked using the same operations as any other NAND Flash Array memory block. OTP block cannot be erased.

OTP block is fully-guaranteed to be a valid block.

## **Entering the OTP Block**

The OTP block is separately accessible from the rest of the NAND Flash Array by using the OTP Access command instead of the Flash Block Address (FBA).

### **Exiting the OTP Block**

To exit the OTP Access Mode, a Cold-, Warm-, Hot-, or NAND Flash Core Reset operation is performed.

### Exiting the OTP Block during an Erase Operation

If the Reset-triggered exit from the OTP Access Mode happens during an Erase Suspend Operation, the erase routine could fail. Therefore to exit from the OTP Access Mode without suspending the erase operation stop, a 'NAND Flash Core Reset' command should be issued.

## The OTP Block Page Assignments

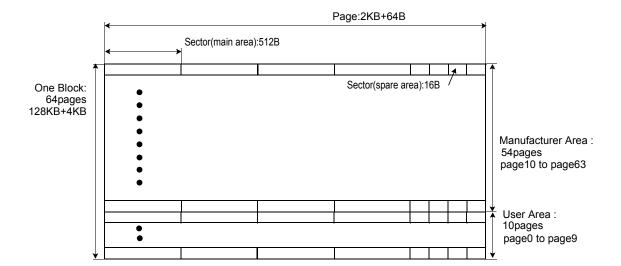
OTP area is one block size (128KB+4KB, 64 Pages) and is divided into two areas. The 10-page User Area is available as an OTP storage area. The 54-page Manufacturer Area is programmed by the manufacturer prior to shipping the device to the user.

## **OTP Block Page Allocation Information**

Area	Page Use			
User 0 ~ 9 (10 pages)		Designated as user area		
Manufacturer	10 ~ 63 (54 pages)	Used by the device manufacturer		



**OTP Area Structure** 



## 3.11.1 OTP Load Operation

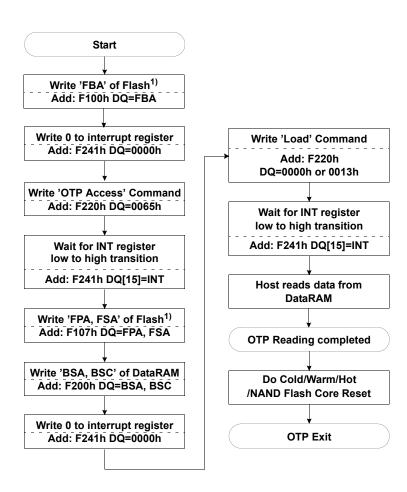
An OTP Load Operation accesses the OTP area and transfers identified content from the OTP to the DataRAM on-chip buffer, thus making the OTP contents available to the Host.

The OTP area is a separate part of the NAND Flash Array memory. It is accessed by issuing OTP Access command(65h) instead of a Flash Block Address (FBA) command.

After being accessed with the OTP Access Command, the contents of OTP memory area are loaded using the same operations as a normal load operation to the NAND Flash Array memory (see section 3.6 for more information).

To exit the OTP access mode following an OTP Load Operation, a Cold-, Warm-, Hot-, or NAND Flash Core Reset operation is performed.

## **OTP Read Operation Flow Chart**



Note 1) FBA(NAND Flash Block Address) could be any address.



## 3.11.2 OTP Program Operation

An OTP Program Operation accesses the OTP area and programs content from the DataRAM on-chip buffer to the designated page(s) of the OTP.

A memory location in the OTP area can be programmed only one time (no erase operation permitted).

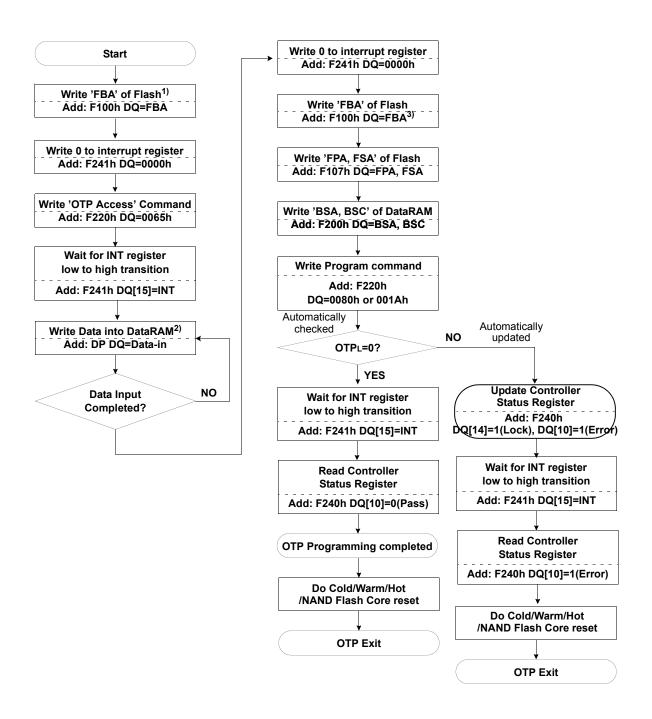
The OTP area is programmed using the same sequence as normal program operation after being accessed by the command (see section 3.8 for more information).

### **Programming the OTP Area**

- · Issue the OTP Access Command
- Write data into the DataRAM (data can be input at anytime between the "Start" and "Write Program" commands
- Issue a Flash Block Address (FBA) which is unlocked area address of NAND Flash Array address map.
- Issue a Write Program command to program the data from the DataRAM into the OTP
- When the OTP programming is complete, do a Cold-, Warm-, Hot-, NAND Flash Core Reset to exit the OTP Access mode.



**OTP Program Operation Flow Chart** 



Note 1) FBA(NAND Flash Block Address) could be any address.

2) Data input could be done anywhere between "Start" and "Write Program Command".

# 3.11.3 OTP Lock Operation

Even though the OTP area can only be programmed once without erase capability, it can be locked when the device starts up to prevent any changes from being made.

Unlike the main area of the NAND Flash Array memory, once the OTP block is locked, it cannot be unlocked.

## Locking the OTP

Programming to the OTP area can be prevented by locking the OTP area. Locking the OTP area is accomplished by programming XXXCh to 8th word of sector0 of page0 of the spare0 memory area.

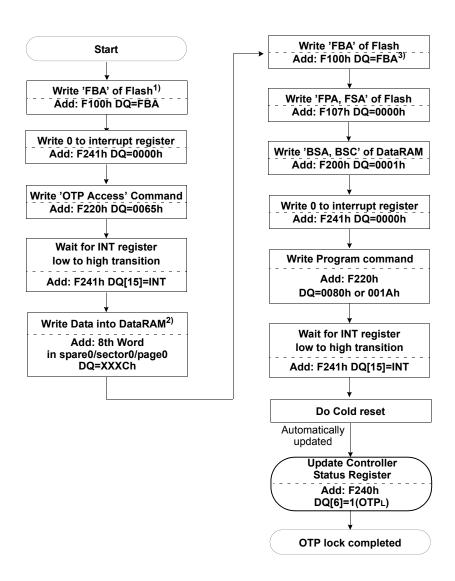
At device power-up, this word location is checked and if XXXCh is found, the OTPL bit of the Controller Status Register is set to "1", indicating the OTP is locked. When the Program Operation finds that the status of the OTP is locked, the device updates the Error Bit of the Controller Status Register as "1" (fail).

## **OTP Lock Operation Steps**

- Issue the OTP Access Command
- Fill data to be programmed into DataRAM (data can be input at anytime between the "Start" and "Write Program" commands)
- Write 'XXXCH' data into the 8th word of sector0 of page0 of the spare0 memory area of the DataRAM.
- Issue a Flash Block Address (FBA) to any address in the NAND Flash Array address map.
- Issue a Program command to program the data from the DataRAM into the OTP
- When the OTP lock is complete, do a Cold Reset to exit the OTP Access mode and update OTP lock bit[6].
- OTP lock bit[6] of the Controller Status Register will be set to "1" and the OTP will be locked.



**OTP Lock Operation Flow Chart** 



Note 1) FBA(NAND Flash Block Address) could be any address.

2) Data input could be done anywhere between "Start" and "Write Program Command".

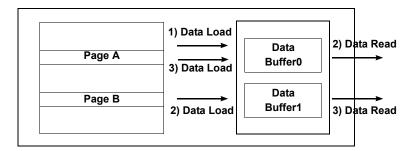
3) FBA should point the unlocked area address among NAND Flash Array address map.

# 3.12 **Dual Operations**

The device has independent dual data buffers on-chip (except during the Boot Load period) that enables higher performance read and program operation.

# 3.12.1 Read-While-Load Operation

This operation accelerates the read performance of the device by enabling data to be read out by the host from one DataRAM buffer while the other DataRAM buffer is being loaded with data from the NAND Flash Array memory.

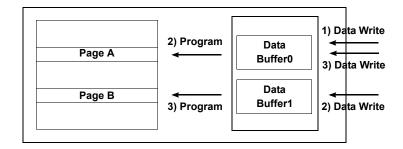


The dual data buffer architecture provides the capability of executing a data-read operation from one of DataRAM buffers during a simultaneous data-load operation from Flash to the other buffer. Simultaneous load and read operation to same data buffer is prohibited. See sections 3.6 and 3.7 for more information on Load and Read Operations.

If host sets FBA, FSA, or FPA while loading into designated page, it will fail the internal load operation. Address registers should not be updated until internal operation is completed.

## 3.12.2 Write-While-Program Operation

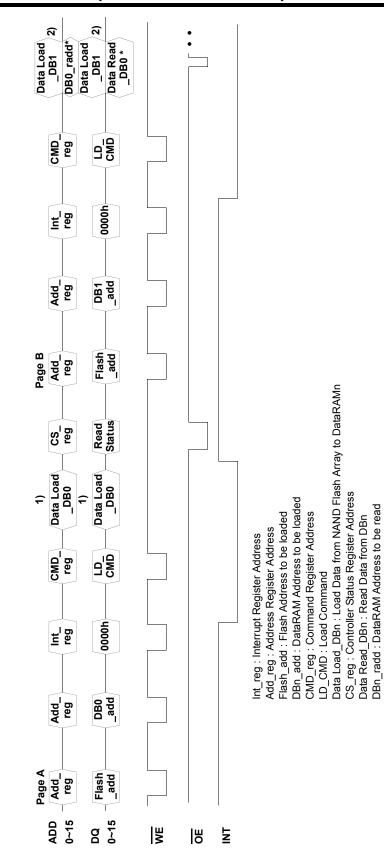
This operation accelerates the programming performance of the device by enabling data to be written by the host into one DataRAM buffer while the NADN Flash Array memory is being programmed with data from the other DataRAM buffer.



The dual data buffer architecture provides the capability of executing a data-write operation to one DataRAM buffers during simultaneous data-program operation to Flash from the other buffer. Simultaneous program and write operation to same data buffer is prohibited. See sections 3.8 for more information on Program Operation.

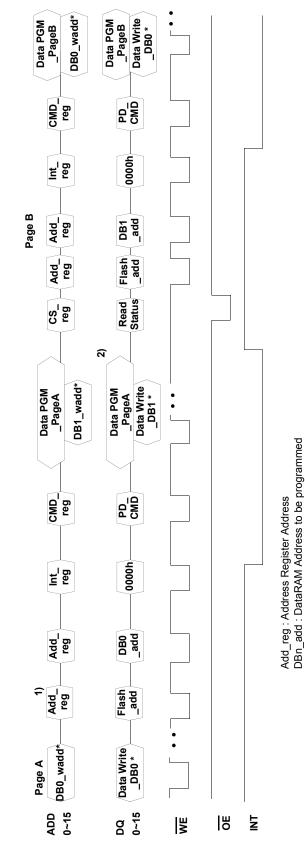
If host sets FBA, FSA, or FPA while programming into designated page, it will fail the internal program operation. Address registers should not be updated until internal operation is completed.







Read While Load Diagram



Flash\_add: Flash Address to be programmed DBn\_wadd: DataRAM Address to be written Data Write\_DBn: Write Data to DataRAMn CMD\_reg : Command Register Address PD\_CMD : Program Command Int\_reg : Interrupt Register Address

Data PGM\_PageA: Program Data from DataRAM to PageA

CS\_reg: Controller Status Register Address

Write While Program Diagram

# 3.13 ECC Operation

The OneNAND device has on-chip ECC with the capability of detecting 2 bit errors and correcting 1-bit errors in the NAND Flash Array memory main and spare areas.

As the device transfers data from a BufferRAM to the NAND Flash Array memory Page Buffer for Program Operation, the device initiates a background operation which generates an Error Correction Code (ECC) of 24bits for each sector main area data and 10bits for 2nd and 3rd word data of each sector spare area.

During a Load operation from the NAND Flash Array memory Page, the on-chip ECC engine generates a new ECC. The 'Load ECC result' is compared to the originally 'Program ECC' thus detecting the number and position of errors. Single-bit error is corrected.

ECC is updated by the device automatically. After a Load Operation, the Host can determine whether there was error by reading the 'ECC Status Register' (refer to section 2.8.26).

Error types are divided into 'no error', '1bit correctable error', and '2bit error uncorrectable error'.

OneNAND supports 2bit EDC even though 2bit error seldom or never occurs. Hence, it is not recommeded for Host to read 'ECC Status Register' for checking ECC error because the built-in Error Correction Logic of OneNAND automatically corrects ECC error.

When the device reads the NAND Flash Array memory main and spare area data with an ECC operation, the device doesn't place the newly generated ECC for main and spare area into the buffer. Instead it places the ECC which was generated and written during the program operation into the buffer.

An ECC operation is also done during the Boot Loading operation.

## 3.13.1 ECC Bypass Operation

In an ECC bypass operation, the device does not generate ECC as a background operation. The result does not indicate error position (refer to the ECC Result Table).

In a Program Operation the ECC code to NAND Flash Array memory spare area is not updated.

During a Load operation, the on-chip ECC engine does not generate a new ECC internally. Also the ECC Status & Result to Registers are invalid. The error is not corrected and detected by itself, so that ECC bypass operation is not recommended for host.

ECC bypass operation is set by the 9bit of System Configuration 1 Register (see section 2.8.19)

## ECC Code and ECC Result by ECC Operation

	Program operation	Load operation					
Operation ECC Code Update to NAND Flash Array Spare Area		ECC Code at BufferRAM Spare Area	ECC Status & Result Update to Registers	1bit Error			
ECC operation	Update	Pre-written ECC code <sup>(1)</sup> loaded	Update	Correct			
ECC bypass	Not update	Pre-written code loaded	Invalid	Not correct			

## NOTE:

1. Pre-written ECC code: ECC code which is previously written to NAND Flash Spare Area in program operation.



# 3.14 Invalid Block Operation

Invalid blocks are defined as blocks in the device's NAND Flash Array memory that contain one or more invalid bits whose reliability is not guaranteed by Samsung.

The information regarding the invalid block(s) is called the Invalid Block Information. Devices with invalid block(s) have the same quality level as devices with all valid blocks and have the same AC and DC characteristics.

An invalid block(s) does not affect the performance of valid block(s) because it is isolated from the bit line and the common source line by a select transistor.

The system design must be able to mask out the invalid block(s) via address mapping. The 1st block, which is placed on 00h block address, is always fully guaranteed to be a valid block.

Due to invalid marking, during load operation for indentifying invalid block, a load error may occur.

## 3.14.1 Invalid Block Identification Table Operation

A system must be able to recognize invalid block(s) based on the original invalid block information and create an invalid block table.

Invalid blocks are identified by erasing all address locations in the NAND Flash Array memory except locations where the invalid block(s) information is written prior to shipping.

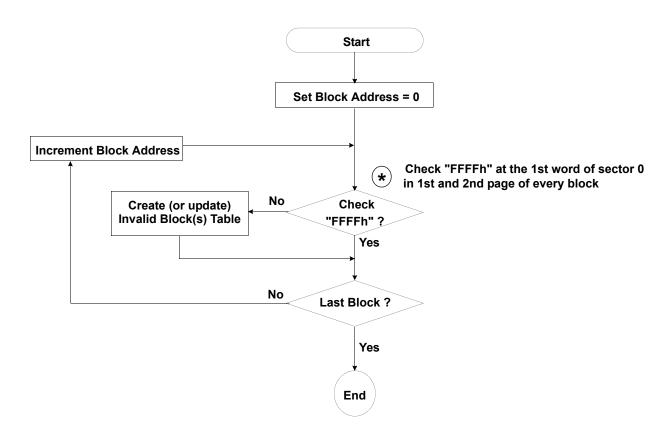
An invalid block(s) status is defined by the 1st word in the spare area. Samsung makes sure that either the 1st or 2nd page of every invalid block has non-FFFFh data at the 1st word of sector0.

Since the invalid block information is also erasable in most cases, it is impossible to recover the information once it has been erased. Any intentional erase of the original invalid block information is prohibited.

The following suggested flow chart can be used to create an Invalid Block Table.



**Invalid Block Table Creation Flow Chart** 



# 3.14.2 Invalid Block Replacement Operation

Within its life time, additional invalid blocks may develop with NAND Flash Array memory. Refer to the device's qualification report for the actual data.

The following possible failure modes should be considered to implement a highly reliable system.

In the case of a status read failure after erase or program, a block replacement should be done. Because program status failure during a page program does not affect the data of the other pages in the same block, a block replacement can be executed with a page-sized buffer by finding an erased empty block and reprogramming the current target data and copying the rest of the replaced block.

## **Block Failure Modes and Countermeasures**

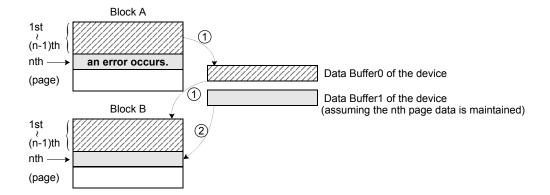
Failure Mode	Detection and Countermeasure sequence
Erase Failure	Status Read after Erase> Block Replacement
Program Failure	Status Read after Program> Block Replacement
Single Bit Failure in Load Operation	Error Correction by ECC mode of the device



Referring to the diagram for further illustration, when an error happens in the nth page of block 'A' during program operation, copy the data in the 1st ~ (n-1)th page to the same location of block 'B' via data buffer0.

Then copy the nth page data of block 'A' in the data buffer1 to the nth page of block 'B' or any free block. Do not further erase or program block 'A' but instead complete the operation by creating an 'Invalid Block Table' or other appropriate scheme.

## **Block Replacement Operation Sequence**





# 4.0 DC CHARACTERISTICS

# 4.1 Absolute Maximum Ratings

Parameter		Symbol	Rating	Unit
	Vcc	Vcc (for 1.8V)	-0.5 to + 2.45	
Voltage on any pin relative to Vss	VCC	Vcc (for 2.65V/3.3V)	-0.5 to + 4.6	V
	All Pins	V <sub>IN</sub> (for 1.8V)	-0.5 to + 2.45	V
	All Pills	V <sub>IN</sub> (for 2.65V/3.3V)	-0.5 to + 4.6	
T	Extended	Thias	-30 to +125	°C
Temperature Under Bias	Industrial	I Dias	-40 to +125	-0
Storage Temperature		Tstg	-65 to +150	°C
Short Circuit Output Current		los	5	mA
Recommended Operating Temperature		TA (Extended Temp.)	-30 to +85	°C
		TA (Industrial Temp.)	-40 to +85	٠

### NOTES

- 1. Minimum DC voltage is -0.5V on Input/ Output pins. During transitions, this level should not fall to POR level(typ. 1.5V@1.8V device, 1.8V@2.65V/ 3.3V device). Maximum DC voltage is Vcc+0.6V on input / output pins which, during transitions, may overshoot to Vcc+2.0V for periods <20ns.
- Permanent device damage may occur if ABSOLUTE MAXIMUM RATINGS are exceeded. Functional operation should be restricted to the conditions detailed in the operational sections of this data sheet. Exposure to absolute maximum rating conditions for extended periods may affect reliability.

# 4.2 Operating Conditions

Voltage reference to GND

Parameter	Symbol	1.8V Device		2.65V Device		3.3V Device			Unit				
	Symbol	Min	Тур.	Max	Min	Тур.	Max	Min	Тур.	Max	Onit		
	Vcc-core / Vcc	1.7	1.7	1.7	1.8	1.95	2.4	2.65	2.9	2.7	3.3	3.6	V
Supply Voltage	Vcc-IO / Vccq	1.7	1.0	1.95	2.4	2.03	2.9	2.1	3.3	3.0	v l		
	Vss	0	0	0	0	0	0	0	0	0	V		

### NOTES:

- 1. The system power should reach 1.7V after POR triggering level(typ. 1.5V) within 400us.
- For 2.65V and 3.3V device, the system power should reach 2.2V after POR triggering level(typ. 1.8V) within 400us.
- 2. Vcc-Core (or Vcc) should reach the operating voltage level prior to or at the same time as Vcc-IO (or Vccq).



# 4.3 DC Characteristics

Barramatar Sumbal 3		T4 O	Took Conditions		V dev	ice	2.6	5V dev	/ice	3.3V device			1114
Parameter	Symbol	Test Conditions	rest Conditions		Тур	Max	Min	Тур	Max	Min	Тур	Max	Unit
Input Leakage Current	ILI	VIN=Vss to Vcc, Vcc=V	CCmax	- 1.0	-	+ 1.0	- 1.0	-	+ 1.0	- 1.0	1	+ 1.0	μА
Output Leakage Cur- rent	llo	VOUT=Vss to Vcc, Vcc= CE or OE=ViH(Note 1)	VCCmax,	- 1.0	-	+ 1.0	- 1.0	-	+ 1.0	- 1.0	-	+ 1.0	μА
Active Asynchronous Read Current (Note 2)	Icc1	CE=VIL, OE=VIH		-	8	15	-	18	22	-	18	22	mA
Active Burst Read Cur-	Icc2	CE=VIL. OE=VIH	54Mhz	-	12	20	-	22	30	-	22	30	mA
rent (Note 2)	1002	CE-VIL, OE-VIH	1MHz	-	3	4	-	7	9	-	7	9	mA
Active Write Current (Note 2)	Іссз	CE=VIL, OE=VIH		-	8	15	-	17	22	-	17	22	mA
Active Load Current (Note 3)	ICC4	CE=VIL, OE=VIH, WE=\	/IH	-	30	40	-	30	40	-	30	40	mA
Active Program Cur- rent (Note 3)	ICC5	CE=VIL, OE=VIH, WE=\	/IH	-	25	30	-	28	35	-	28	35	mA
Active Erase Current (Note 3)	Icc6	CE=VIL, OE=VIH, WE=\	/IH	-	20	25	-	23	30	-	23	30	mA
Multi Block Erase Cur- rent (Note 3)	ICC7	CE=VIL, OE=VIH, WE=\ 64blocks	/IH,	-	20	25	-	23	30	-	23	30	mA
Standby Current	Isb	CE= RP=Vcc ± 0.2V		-	10	50	-	35	80	-	35	80	μА
Input Low Voltage	VIL	-		-0.5	-	0.4	-0.5	-	0.4	0	1	8.0	V
Input High Voltage	VIH	-		Vccq- 0.4	-	VCCq +0.4	Vccq- 0.4	-	VCCq +0.4	0.7*V CCq	-	Vccq	٧
Output Low Voltage	Vol	IOL = 100 μA , VCC=VCC VCCq=VCCqmin	Cmin ,	-	-	0.2	-	-	0.2	-	1	0.22* Vccq	٧
Output High Voltage	Vон	IOH = -100 μA , VCC=VC	CCmin ,	Vccq- 0.1	-	-	VCCq- 0.4	-	-	0.8*V CCq	1	-	٧

Note 1.  $\overline{\text{CE}}$  should be VIH for RDY. IOBE should be '0' for INT.

Note 2. Icc active for Host access

Note 3. ICC active for Internal operation. (without host access)

# 5.0 AC CHARACTERISTICS

# 5.1 AC Test Conditions

Parameter		Value
Input Pulse Levels		0V to Vcc
Innut Discound Fall Times	CLK	3ns
Input Rise and Fall Times	other inputs	5ns
Input and Output Timing Levels		Vcc/2
Output Load		CL = 30pF



# 5.2 Device Capacitance

**CAPACITANCE**(TA = 25 °C, VCC = 1.8V/2.65V/3.3V, f = 1.0MHz)

ltem	Symbol Test Condition		KFG12	Unit	
	Symbol	rest Condition	Min	Max	Offic
Input Capacitance	CIN1	VIN=0V	-	10	pF
Control Pin Capacitance	CIN2	V <sub>IN</sub> =0V	-	10	pF
Output Capacitance	Соит	Vout=0V	-	10	pF
INT Capacitance	CINT	Vout=0V	-	15	pF

NOTE: Capacitance is periodically sampled and not 100% tested.

# **5.3** Valid Block Characteristics

Parameter	Symbol	Min	Тур.	Max	Unit
Valid Block Number	N∨B	502	-	512	Blocks

### NOTES:

2. The 1st block, which is placed on 00h block address, is fully guaranteed to be a valid block.



<sup>1.</sup> The device may include invalid blocks when first shipped. Additional invalid blocks may develop while being used. The number of valid blocks is presented with both cases of invalid blocks considered. Invalid blocks are defined as blocks that contain one or more bad bits. Do not erase or program factory-marked bad blocks.

## **AC Characteristics for Synchronous Burst Read** 5.4

See Timing Diagrams 6.1 and 6.2

Parameter	Symbol	KFG12	- Unit	
Parameter	Symbol	Min	Max	Unit
Clock	CLK	1	54	MHz
Clock Cycle	tclk	18.5	-	ns
Initial Access Time	tiaa	-	76	ns
Burst Access Time Valid Clock to Output Delay	tва	-	14.5	ns
AVD Setup Time to CLK	tavds	7	-	ns
AVD Hold Time from CLK	tavdh	7	-	ns
Address Setup Time to CLK	tacs	7	-	ns
Address Hold Time from CLK	tach	7	-	ns
Data Hold Time from Next Clock Cycle	tврн	4	-	ns
Output Enable to Data	toe	-	20	ns
CE Disable to Output High Z	tcez <sup>1)</sup>	-	20	ns
OE Disable to Output High Z	toez <sup>1)</sup>	-	17	ns
CE Setup Time to CLK	tces	7	-	ns
CLK High or Low Time	tclkh/L	tclk/3	-	ns
CLK <sup>2)</sup> to RDY valid	trdyo	-	14.5	ns
CLK to RDY Setup Time	trdya	-	14.5	ns
RDY Setup Time to CLK	trdys	4	-	ns
CE low to RDY valid	tcer	-	15	ns

Note\_\_\_\_\_\_\_

1. If  $\overline{OE}$  is disabled at the same time or before  $\overline{CE}$  is disabled, the output will go to high-z by toez. If  $\overline{CE}$  is disabled at the same time or before  $\overline{OE}$  is disabled, the output will go to high-z by toez. If  $\overline{CE}$  and  $\overline{OE}$  are disabled at the same time, the output will go to high-z by toez.

<sup>2.</sup> It is the following clock of address fetch clock.

### 5.5 **AC Characteristics for Asynchronous Read**

See Timing Diagrams 6.3, 6.4, 6.5 and 6.6

Parameter	Symbol	KFG12	KFG1216x2A			
Faranietei	Syllibol	Min	Max	Unit		
Access Time from CE Low	tce	-	76	ns		
Asynchronous Access Time from AVD Low	taa	-	76	ns		
Asynchronous Access Time from address valid	tacc	-	76	ns		
Read Cycle Time	trc	76	-	ns		
AVD Low Time	tavdp	12	-	ns		
Address Setup to rising edge of AVD	taavds	7	-	ns		
Address Hold from rising edge of AVD	taavdh	7	-	ns		
Output Enable to Output Valid	toE	-	20	ns		
WE disable to OE enable	toeh	0	-	ns		
CE Setup to AVD falling edge	tca	0	-	ns		
CE Disable to Output & RDY High Z <sup>1)</sup>	tcez	-	20	ns		
OE Disable to Output High Z <sup>1)</sup>	toez	-	17	ns		
WE Disable to AVD Enable	twea	15	-	ns		

### NOTE:

These parameters are not 100% tested.

## AC Characteristics for Warm Reset (RP), Hot Reset 5.6 and NAND Flash Core Reset

See Timing Diagrams 6.12, 6.13 and 6.14

Parameter	Symbol	Min	Max	Unit
RP & Reset Command Latch to BootRAM Access	tReady1 (BufferRAM)	-	5	μS
RP & Reset Command Latch(During Load Routines) to INT High (Note1)	tReady2 (NAND Flash Array)	-	10	μS
RP & Reset Command Latch(During Program Routines) to INT High (Note1)	tReady2 (NAND Flash Array)	-	20	μS
RP & Reset Command Latch(During Erase Routines) to INT High (Note1)	tReady2 (NAND Flash Array)	-	500	μS
RP & Reset Command Latch(NOT During Internal Routines) to INT High (Note1)	tReady2 (NAND Flash Array)	-	10	μS
RP Pulse Width (Note2)	tRP	200	-	ns



In If  $\overline{\underline{CE}}$  is disabled at the same time or before  $\overline{\underline{CE}}$  is disabled, the output will go to high-z by toez. If  $\overline{\underline{CE}}$  is disabled at the same time or before  $\overline{\underline{OE}}$  is disabled, the output will go to high-z by toez. If  $\overline{\underline{CE}}$  and  $\overline{\underline{OE}}$  are disabled at the same time, the output will go to high-z by toez.

<sup>1.</sup> These parameters are tested based on INT bit of interrupt register. Because the time on INT pin is related to the pull-up and pull-down resistor value.

<sup>2.</sup> The device may reset if tRP < tRP min(200ns), but this is not guaranteed.

# 5.7 AC Characteristics for Asynchronous Write/Load/ Program/Erase Operation

See Timing Diagrams 6.7, 6.8 and 6.9

Parameter	Symbol	Min	Тур	Max	Unit
WE Cycle Time	twc	70	-	-	ns
AVD low pulse width	tavdp	12	-	-	ns
Address Setup Time	tawes	0	-	-	ns
Address Hold Time	tан	30	-	-	ns
Data Setup Time	tos	25	-	-	ns
Data Hold Time	tон	0	-	-	ns
CE Setup Time	tcs	0	-	-	ns
CE Hold Time	tсн	10	-	-	ns
WE Pulse Width	twpL	40	-	-	ns
WE Pulse Width High	twpн	30	-	-	ns
AVD Disable to WE Disable	t∨LWH	15	-	-	ns
WE Disable to AVD Enable	twea	15	-	-	ns

# 5.8 AC Characteristics for Load/Program/Erase Performance

See Timing Diagrams 6.8, 6.9, and 6.10

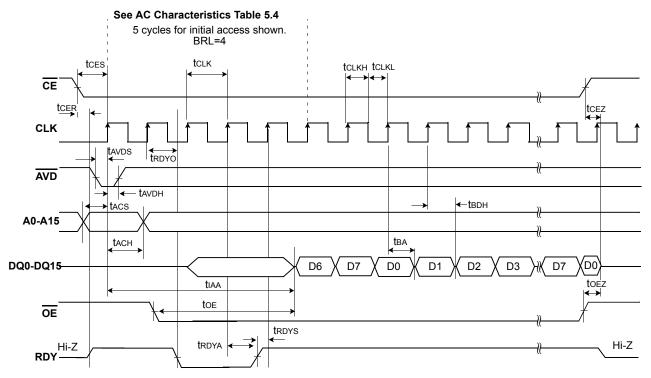
Parameter		Symbol	Min	Тур	Max	Unit
Sector Read time(Note 1)		tRD1	-	23	35	μS
Page Read time(Note 1)		trd2	-	30	45	μS
Sector Program time(Note 1)		tPGM1	-	205	720	μS
Page Program time(Note 1)		tPGM2	-	220	750	μS
OTP Access Time(Note 1)		tотр	-	500	700	ns
Lock/Unlock/Lock-tight Time(Note 1)		tLOCK	-	500	700	ns
Erase Suspend Time(Note 1)		tesp	-	400	500	μS
Erase Resume Time(Note 1)	1 Block	ters1	-	2	3	ms
	2~64 Blocks	ters2		4	6	ms
Number of Partial Program Cycles in the sector (Including main and spare area)		NOP	-	-	2	cycles
Block Erase time (Note 1)	1 Block	tBERS1	-	2	3	ms
	2~64 Blocks	tBERS2	-	4	6	ms
Multi BlocK Erase Verify Read time(Note 1)		trd3	-	70	100	μS

Note 1)These parameters are tested based on INT bit of interrupt register. Because the time on INT pin is related to the pull-up and pull-down resistor value.

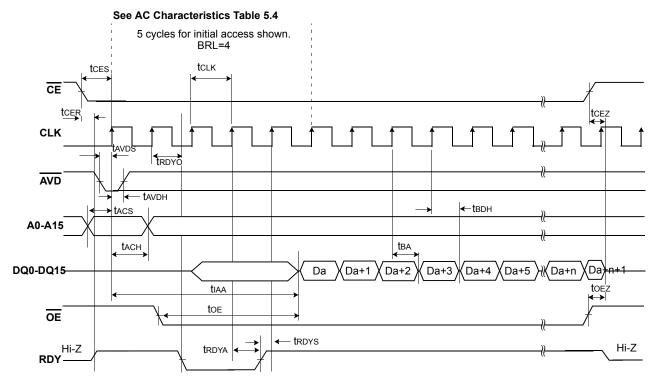


# 6.0 TIMING DIAGRAMS

# 6.1 8-Word Linear Burst Mode with Wrap Around



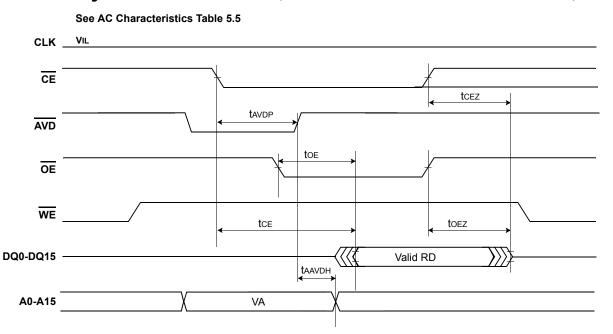
# 6.2 Continuous Linear Burst Mode with Wrap Around





Hi-Z

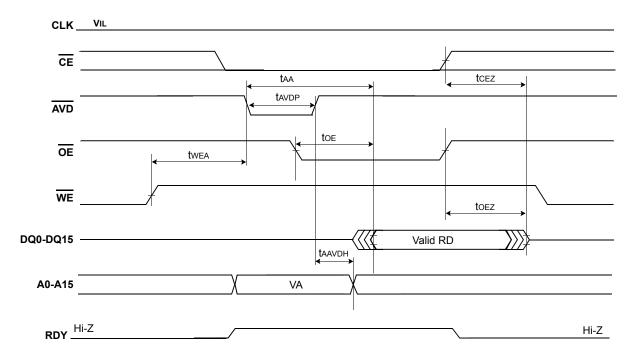
## Asynchronous Read (VA and AVD Transition Before CE Low) 6.3



NOTE: VA=Valid Read Address, RD=Read Data.

## Asynchronous Read (VA and AVD Transition After CE Low) 6.4

See AC Characteristics Table 5.5



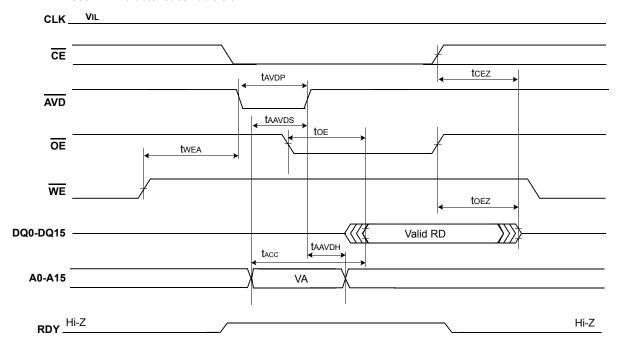
NOTE: VA=Valid Read Address, RD=Read Data.



RDY Hi-Z

# 6.5 Asynchronous Read (VA and AVD Transition After CE Low)

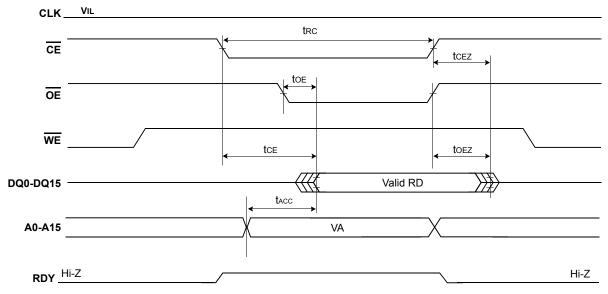
See AC Characteristics Table 5.5



NOTE: VA=Valid Read Address, RD=Read Data.

# 6.6 Asynchronous Read (AVD is tied to CE)

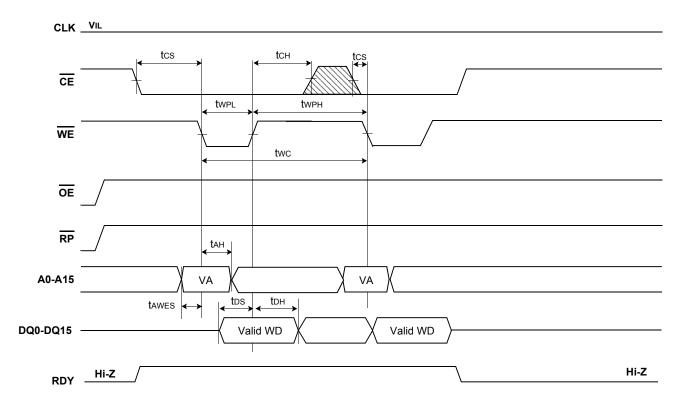
See AC Characteristics Table 5.5



NOTE: VA=Valid Read Address, RD=Read Data.

# 6.7 Asynchronous Write

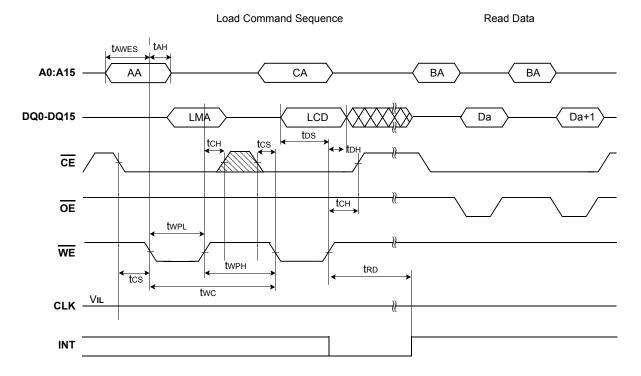
See AC Characteristics Table 5.7



NOTE: VA=Valid Read Address, WD=Write Data.

# 6.8 Load Operation Timing

See AC Characteristics Tables 5.7 and 5.8



### NOTES:

- 1. AA = Address of address register
  - CA = Address of command register
  - LCD = Load Command
  - LMA = Address of memory to be loaded
  - BA = Address of BufferRAM to load the data
  - BD = Program Data
  - SA = Address of status register
- 2. "In progress" and "complete" refer to status register
- 3. Status reads in this figure is asynchronous read, but status read in synchronous mode is also supported.

# 6.9 Program Operation Timing

See AC Characteristics Tables 5.7 and 5.8

Program Command Sequence (last two cycles) Read Status Data tawes tан A0:A15 ВА CA SA SA DQ0-DQ15 BD Complete tos CE tсн OE twpl WE twph tcs **t**PGM twc CLK VIL INT

## NOTES:

1. AA = Address of address register

CA = Address of command register

PCD = Program Command

PMA = Address of memory to be programmed

BA = Address of BufferRAM to load the data

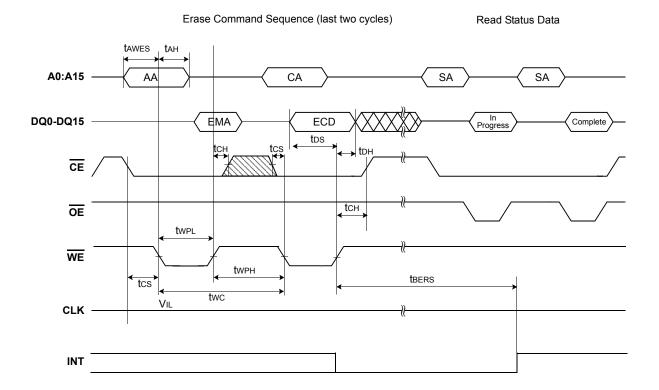
BD = Program Data

SA = Address of status register

- 2. "In progress" and "complete" refer to status register
- 3. Status reads in this figure is asynchronous read, but status read in synchronous mode is also supported.

# 6.10 Block Erase Operation Timing

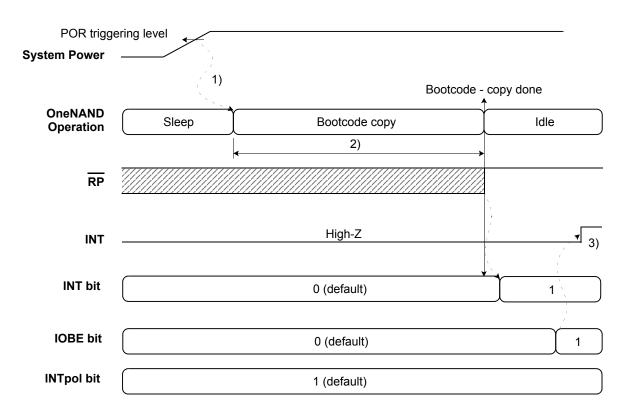
See AC Characteristics Tables 5.7 and 5.8



### NOTES:

- 1. AA = Address of address register
  - CA = Address of command register
  - ECD = Erase Command
  - EMA = Address of memory to be erased
  - SA = Address of status register
- 2. "In progress" and "complete" refer to status register
- 3. Status reads in this figure is asynchronous read, but status read in synchronous mode is also supported.

# 6.11 Cold Reset Timing



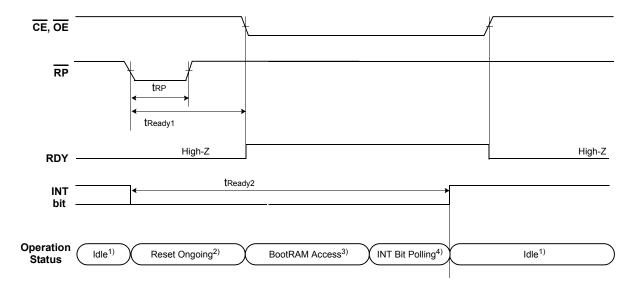
Note: 1) Bootcode copy operation starts 400us later than POR activation.

The system power should reach Vcc after POR triggering level(typ. 1.5V) within 400us for valid boot code data.

- 2) 1K bytes Bootcode copy takes 70us(estimated) from sector0 and sector1/page0/block0 of NAND Flash array to BootRAM. Host can read Bootcode in BootRAM(1K bytes) after Bootcode copy completion.
- 3) INT register goes 'Low' to 'High' on the condition of 'Bootcode-copy done' and  $\overline{RP}$  rising edge.
- If RP goes 'Low' to 'High' before 'Bootcode-copy done', INT register goes to 'Low' to 'High' as soon as 'Bootcode-copy done'

# 6.12 Warm Reset Timing

See AC Characteristics Tables 5.6

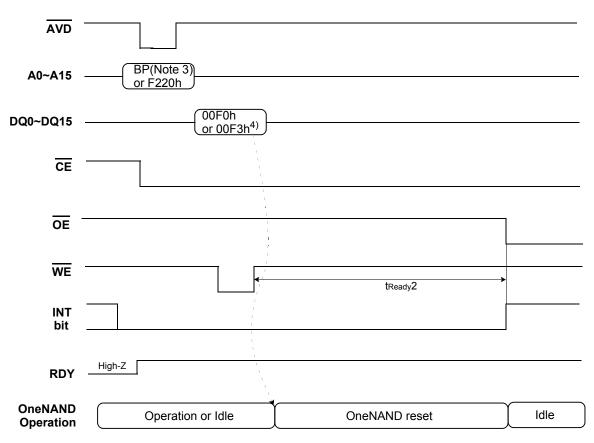


### NOTES:

- 1. The status which can accept any register based operation(Load, Program, Erase command, etc).
- 2. The status where reset is ongoing.
- 3. The status allows only BootRAM(BL1) read operation for Boot Sequence.(refer to 7.2.2 Boot Sequence)
- 4. To read BL2 of Boot Sequence, Host should wait INT until becomes ready. and then, Host can issue load command. (refer to 7.2.2 Boot Sequence, 7.1 Methods of Determing Interrupt status)



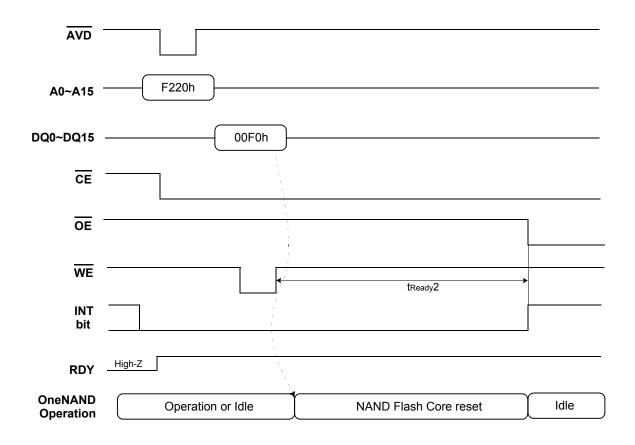
# 6.13 Hot Reset Timing



### NOTE:

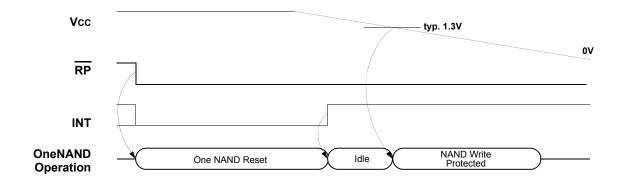
- 1. Internal reset operation means that the device initializes internal registers and makes output signals go to default status and bufferRAM data are kept unchanged after Warm/Hot reset operations.
- 2. Reset command : Command based reset or Register based reset
- 3. BP(Boot Partition): BootRAM area [0000h~01FFh, 8000h~800Fh]
- 4. 00F0h for BP, and 00F3h for F220h

# 6.14 NAND Flash Core Reset Timing



# 6.15 Data Protection Timing During Power Down

The device is designed to offer protection from any involuntary program/erase during power-transitions. An internal voltage detector disables all functions whenever Vcc is below about 1.3V.  $\overline{\text{RP}}$  pin provides hardware protection and is recommended to be kept at VIL before power-down.





# 7.0 TECHNICAL AND APPLICATION NOTES

From time-to-time supplemental technical information and application notes pertaining to the design and operation of the device in a system are included in this section. Contact your Samsung Representative to determine if additional notes are available.

# 7.1 Methods of Determining Interrupt Status

There are two methods of determining Interrupt Status on the OneNAND. Using the INT pin or monitoring the Interrupt Status Register Bit.

The OneNAND INT pin is an output pin function used to notify the Host when a command has been completed. This provides a hardware method of signaling the completion of a program, erase, or load operation.

In its normal state, the INT pin is high if the INT polarity bit is default. Before a command is written to the command register, the INT bit must be written to '0' so the INT pin transitions to a low state indicating start of the operation. Upon completion of the command operation by the OneNAND's internal controller, INT returns to a high state.

INT is an open drain output allowing multiple INT outputs to be Or-tied together. INT does not float to a hi-Z condition when the chip is deselected or when outputs are disabled. Refer to section 2.8 for additional information about INT.

INT can be implemented by tying INT to a host GPIO or by continuous polling of the Interrupt status register.

# 7.1.1 The INT Pin to a Host General Purpose I/O

INT can be tied to a Host GPIO to detect the rising edge of INT, signaling the end of a command operation.

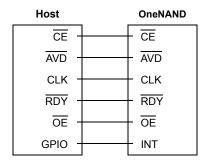


This can be configured to operate either synchronously or asynchronously as shown in the diagrams below.



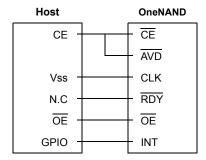
## Synchronous Mode Using the INT Pin

When operating synchronously, INT is tied directly to a Host GPIO.



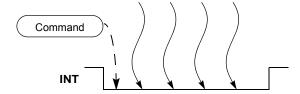
### **Asynchronous Mode Using the INT Pin**

When configured to operate in an asynchronous mode, /CE and /AVD of the OneNAND are tied to /CE of the Host. CLK is tied to the Host Vss (Ground). /RDY is tied to a no-connect. /OE of the OneNAND and Host are tied together and INT is tied to a GPIO.



# 7.1.2 Polling the Interrupt Register Status Bit

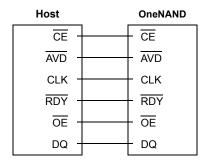
An alternate method of determining the end of an operation is to continuously monitor the Interrupt Status Register Bit instead of using the INT pin.



This can be configured in either a synchronous mode or an asynchronous mode.

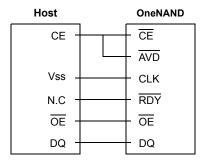
## Synchronous Mode Using Interrupt Status Register Bit Polling

When operating synchronously, /CE, /AVD, CLK, /RDY, /OE, and DQ pins on the host and OneNAND are tied together.



## Asynchronous Mode Using Interrupt Status Register Bit Polling

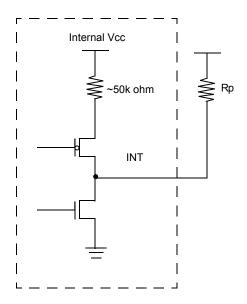
When configured to operate in an asynchronous mode, /CE and /AVD of the OneNAND are tied to /CE of the Host. CLK is tied to the Host Vss (Ground). /RDY is tied to a no-connect. /OE and DQ of the OneNAND and Host are tied together.

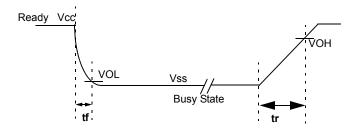


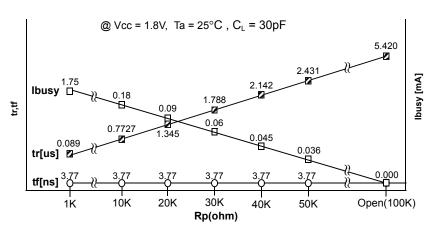
# 7.1.3 Determining Rp Value

Because the pull-up resistor value is related to tr(INT) an appropriate value can obtained with the following reference charts.

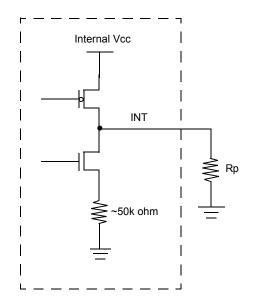
# INT pol = 'High'

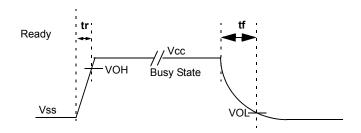


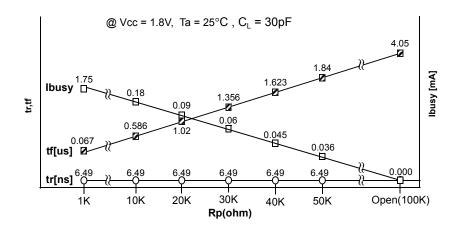




INT pol = 'Low'







## 7.2 Boot Sequence

One of the best features OneNAND has is that it can be a booting device itself since it contains an internally built-in boot loader despite the fact that its core architecture is based on NAND Flash. Thus, OneNAND does not make any additional booting device necessary for a system, which imposes extra cost or area overhead on the overall system.

As the system power is turned on, the boot code originally stored in NAND Flash Arrary is moved to BootRAM automatically and then fetched by CPU through the same interface as SRAM's or NOR Flash's if the size of the boot code is less than 1KB. If its size is larger than 1KB and less than or equal to 3KB, only 1KB of it can be moved to BootRAM automatically and fetched by CPU, and the rest of it can be loaded into one of the DataRAMs whose size is 2KB by Load Command and CPU can take it from the DataRAM after finishing the code-fetching job for BootRAM. If its size is larger than 3KB, the 1KB portion of it can be moved to BootRAM automatically and fetched by CPU, and its remaining part can be moved to DRAM through two DataRAMs using dual buffering and taken by CPU to reduce CPU fetch time.

A typical boot scheme usually used to boot the system with OneNAND is explained at Patition of NAND Flash Array and OneNAND Boot Sequence. In this boot scheme, boot code is comprised of BL1, where BL stands for Boot Loader, BL2, and BL3. Moreover, the size of the boot code is larger than 3KB (the 3rd case above). BL1 is called primary boot loader in other words. Here is the table of detailed explanations about the function of each boot loader in this specific boot scheme.

## 7.2.1 Boot Loaders in OneNAND

### **Boot Loaders in OneNAND**

Boot Loader	Description		
BL1	Moves BL2 from NAND Flash Array to DRAM through two DataRAMs using dual buffering		
BL2	Moves OS image (or BL3 optionally) from NAND Flash Array to DRAM through two DataRams using dual buffering		
BL3 (Optional)	Moves or writes the image through USB interface		

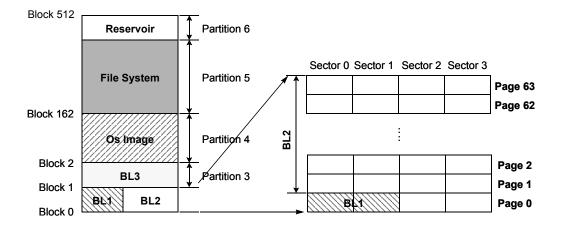
NAND Flash Array of OneNAND is divided into the partitions as described at Partition of NAND Flash array to show where each component of code is located and how much portion of the overall NAND Flash Array each one occupies. In addition, the boot sequence is listed below and depicted at Boot Sequence.

## 7.2.2 Boot Sequence

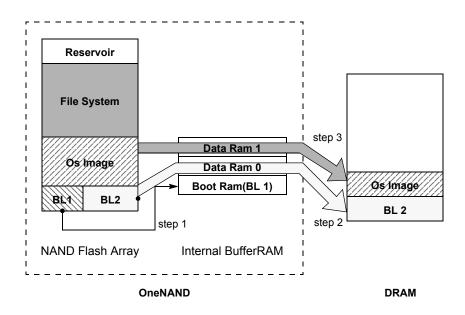
## Boot Sequence :

- Power is on BL1 is loaded into BootRAM
- BL1 is executed in BootRAM
   BL2 is loaded into DRAM through two DataRams using dual buffering by BL1
- BL2 is executed in DRAM
   OS image is loaded into DRAM through two DataRams using dual buffering by BL2
- 4. OS is running





Partition of NAND Flash array



NOTE: Step 2 and Step 3 can be copied into DRAM through two DataRAMs using dual buffering

## **OneNAND Boot Sequence**

# 8.0 PACKAGE DIMENSIONS

